

1 **Optimal configuration of a power grid system with a dynamic performance sharing**
2 **mechanism¹**

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11 **Abstract:** Performance sharing is an effective policy for a power grid system to satisfy the
12 power demand of different districts to greatest extent. Through transmission lines, the districts
13 with sufficient power can share the redundant power with the districts with power deficit. The
14 existing research has incorporated the performance sharing mechanism into systems with
15 simple structures such as parallel systems and series-parallel systems. However, little
16 concentration has been spent on more complex structures. This necessitates the need of this
17 paper that models a power distribution with a more complex reliability structure. We assume
18 that the system is composed of generators and nodes. Both the performance of each generator
19 and the demand of each node in the network are assumed to be random variables. This paper
20 first proposes a dynamic performance sharing policy to minimize the unsupplied demand for a
21 given system with fixed capacity and demand. The optimal allocation of generators, which
22 minimizes the expected system unsupplied demand, is then studied. Numerical examples are
23 proposed to illustrate the applications of the proposed procedures.

24 **Keywords:** power grid system; performance sharing mechanism; optimal generator allocation
25 policy; Hybridized Particle Swarm Optimization (HPSO) algorithm

26 **1. Introduction**

27 Power grid systems are essential for people's daily life and industrial production processes (Li
28 and Zio, 2012; Rodgers et al., 2018; Figueroa-Candia et al., 2018). Therefore, it is important to
29 ensure that the power demand is met (Zimmerman et al., 2011; Zhai et al., 2018). One of the
30 solutions is to optimally allocate electricity in real time according to the supply and demand
31 within a power grid system (Wang et al., 2018). Frank et al. (2016) and Lavaei et al. (2012)
32 studied the optimization problems in the electric power systems, known collectively as Optimal
33 Power Flow (OPF). For the future smart grids, another solution to meet power demand is
34 Demand Response (DR), which is achieved by active customer participation in real time to
35 maintain balance between generation and demand with two-way communication (Pourmousavi
36 and Nehrir, 2014). DR solutions can be deployed to encourage electricity consumers in

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37 scheduling their controllable demands to off-peak hours (Bahrami et al., 2017). A power grid
38 system may contain many generators, which are used to satisfy the power demand in different
39 districts (Xiang et al., 2017). As power production may be influenced by the state of generators
40 and the demand may also vary, the power produced in a district may not be able to satisfy its
41 power demand (Faza, 2017). In light of this, the performance sharing mechanism is essential
42 such that the districts with sufficient power supply can share the redundant power with the
43 districts with deficient power supply (Tolani & Sensarma, 2017). The power grid system can
44 be regarded reliable as long as the demands of all the districts can be satisfied after proper
45 power redistribution.

46 Some works studied the reliability of systems with performance sharing mechanism. Levitin
47 (2011) studied a system consisting of some unreparable units, each of which has a random
48 performance and a random demand to be satisfied. The units can share performance with each
49 other but the total amount of shared performance is limited by the bandwidth of the performance
50 sharing mechanism. The system works if the demand of all the units can be satisfied with proper
51 performance sharing. A Universal Generating Function (UGF) technique is proposed to
52 evaluate the reliability of the system. Xiao and Peng (2014) analyzed the case where the system
53 has several subsystems to share performance with each other and each subsystem contains
54 several units structured in parallel. The optimal allocation of units and the optimal preventive
55 maintenance interval of the units are investigated. Xiao and Peng (2014) studied the optimal
56 maintenance and protection of the system units, which are subject to both internal failures and
57 external impacts. Yu et al. (2014) extended Levitin (2011) to the case of repairable systems.
58 The availability of each system unit is evaluated first by considering the state transitions, which
59 are due to degradation or maintenance, and then the UGF technique is used to calculate the
60 system availability. Peng et al. (2017) extended Levitin (2011) to the case where the
61 performance sharing group has limited size and studied the optimal choice of units to include
62 in the performance sharing group. Yu et al. (2017) proposed a model of a system with multiple
63 phases, where the units can share performance with each other in each phase. Zhai et al. (2017)
64 studied the optimal defense and attack of a system with a performance sharing mechanism.
65 Zhao et al. (2018) studied a k -out-of- n system with performance sharing mechanism. Wu et al.
66 (2019) studied the reliability of capacitated systems with a performance sharing mechanism.

67 However, all the above works are restricted to systems with single performance sharing
68 group. Peng et al. (2016) considered a different scenario, where a system has two performance
69 sharing groups. Later, Peng (2019) studied a system with hierarchical performance sharing
70 groups. However, these models are still not able to accommodate more distinct system
71 structures and more arbitrary performance sharing among the system units. In practical systems,
72 what is constrained may not be the total amount of performance that can be shared in the whole
73 system or a subsystem. Instead, the performance sharing between any two nodes may be
74 constrained by the performance sharing mechanism between them. For example, a power grid
75 system may be of a networked structure with many nodes representing different districts, and
76 the power sharing between each two nodes may be constrained by the bandwidth of the
77 transmission lines between them. In this paper, a general network model is proposed to model
78 a power grid system with the performance sharing mechanism. Each node in the network has a
79 random performance and a random demand to be satisfied. The link between any two nodes

80 serves a channel for the performance sharing between the two nodes, and has a bandwidth that
 81 restricts the amount of performance that can be shared. For the case where the allocation of
 82 generators, the realization of the generators' capacity and nodes' demand is given, a dynamic
 83 performance sharing mechanism is proposed to minimize the unsupplied demand. The expected
 84 system unsupplied demand can be obtained by taking into account all the possible realizations
 85 of the nodes' performance and demand. The allocation of generators, which minimizes the
 86 expected system unsupplied demand, can be solved.

87 The remainder of this paper is as follows. Section 2 describes the system and assumptions.
 88 Section 3 focuses on the optimal power sharing problem for given realization of the nodes'
 89 demand and the generators' capacity and given generators' allocation. The case where the power
 90 transmission loss can be neglected and the case where the power transmission loss cannot be
 91 neglected are considered. Section 4 formulates the optimal allocation of the generators that
 92 minimizes the expected unsupplied demand of the systems. Section 5 presents the optimization
 93 techniques adopted. Section 6 presents numerical examples to illustrate the application of the
 94 proposed model. Section 7 presents a case study based on practical background. Section 8
 95 concludes the paper.

96 2. System description

97 Consider a system consisting of N nodes, from node 1 to node N , representing N districts in
 98 need of power supply. The number of available generators is M . M generators are assigned
 99 to N nodes, and the system structure is shown in the Fig 1.

100 Fig 1. A typical power grid system

101 Where, a box represents a node, a line represents a power transmission line, and M_{ij_r} refers
 102 to the j_r th generator in the i th node. $i=1,2,\dots$ $j_1 + j_2 + \dots$ M .

103 At each node, several power generators can be deployed. It is assumed that both the power
 104 to be generated by each generator and the power demand of each district are random variables.
 105 It is also assumed that the power generated by different generators is statistically independent,
 106 so are the power demands on different nodes. In fact, the demand and generating capacity are
 107 changing from time to time (Lisnianski & Ding, 2016). However, if one focuses on the
 108 demand and the generated power for a specific period of concern, then both the demand and
 109 capacity are random variables. In practical applications, the probability distribution of a
 110 generator's capacity can be estimated from historical data. Similarly, the power demand at each
 111 node may also change from time to time under the influence of various factors and its
 112 probability distribution can be estimated from historical data of power consumption. To
 113 facilitate discussion, the distribution for the generators' capacity is also called a capacity
 114 distribution, and the distribution for the demand at each node is also called a demand
 115 distribution.

116 Suppose that the capacity of generator k ($k=1,2,\dots,M$) has realizations

117 $a_k(1) < a_k(2) < \dots < a_k(L_k)$, where L_k is the number of possible realizations for the

118 capacity of generator k . The corresponding probabilities for $a_k(1), a_k(2), \dots, a_k(L_k)$
119 are $p_k(1), p_k(2), \dots, p_k(L_k)$, respectively, where $p_k(1) + p_k(2) + \dots + p_k(L_k) = 1$. The
120 demand of node $i (i = 1, 2, \dots, N)$ has realizations $d_i(1) < d_i(2) < \dots < d_i(W_i)$, where W_i
121 is the number of possible realizations for the demand of node i . The corresponding
122 probabilities for $d_i(1), d_i(2), \dots, d_i(W_i)$ are $q_i(1), q_i(2), \dots, q_i(W_i)$, where
123 $q_i(1) + q_i(2) + \dots + q_i(W_i) = 1$. Between each pair of districts, there may be a power
124 distribution line facilitating the power transmission between them. The amount of power
125 transmission is constrained by the bandwidth of the transmission line between them. Typically,
126 the maximum amount of the power transmission between nodes i and j is denoted as c_{ij} .
127 The case of no transmission line between nodes i and j corresponds to $c_{ij} = 0$. If the
128 allocation of generators, the realization of the generators' capacity and the nodes' demand are
129 given, the power is dynamically transmitted between the system nodes so that the unsupplied
130 demand of the system is minimized. The optimal allocation of the generators that minimizes
131 the expected unsupplied demand of the system needs to be solved.

132 The notations used in this paper are defined as follows.

Notations	Definitions
N	The number of nodes
M	The number of generators
c_{ij}	The transmission bandwidth between nodes i and j
d_i	The power demand at node i
g_i	The amount of power produced by generators from node i
X_{ij}	The performance transmitted from any node i to any other node j

$K_{ij}, (0 \leq K_{ij} \leq 1)$	The loss rate for power transmission from node i to node j
f_{before}	The unsupplied power demand before performance sharing
f_{after}	The unsupplied power demand after performance sharing
L_k	The number of possible realizations for the capacity of generator k
$u_k(z) = \sum_{l=1}^{L_k} p_k(l) z^{a_k(l)},$	The UGF used to represent the capacity distribution of generator k
$a_k(l)$	The capacity realization of generator
$p_k(l)$	The corresponding probability of the capacity realization of the generator
$d_i(l)$	The demand realization of the node i
$q_i(l)$	The corresponding probability of the demand realization of the node i .
$n(i)$	The number of generators allocated into node i
$g_i(l)$	The capacity realization of the node i
$p_i(l)$	The corresponding probability of the capacity realization of the node i
$H = \{h(k), 1 \leq k \leq M\}$	The generator allocation strategy where each generators k is assigned to the $h(k)$ th node
W_i	The number of demand realizations for the node i
$V = \prod_{i=1}^N L_{ig} W_i$	The number of different realizations of the generators' capacity and the nodes' demand
f_t	The minimized unsupplied demand for each realization t

$P_{f_i}(t=1,2,\dots)$	The probability for the realization
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133 **3. Optimal power sharing policy**

134 For a given generator allocation policy and given realizations of the generators' performances
 135 and the nodes' demand, the power sharing policy that minimizes the unsupplied demand of the
 136 power grid needs to be solved. To facilitate discussion, the following assumptions are made:

137 (1) Suppose the total amount of power produced by generators from node i is g_i , the power
 138 demand at node i is d_i . The excess supply at node i is therefore given by $\max(g_i - d_i, 0)$,
 139 the deficit at node is $\max(d_i - g_i, 0)$.

140 (2) Between each pair of districts, there may be a power distribution line facilitating the power
 141 transmission. The amount of power that can be transmitted between each pair of nodes is
 142 constrained by the bandwidth of the transmission line between them. We assume that the
 143 maximum amount of power that can be transmitted from node i to node j is constrained by
 144 the bandwidth c_{ij} .

145 (3) Let X_{ij} be the power exchange from any node i to any other node j . In practice, the
 146 power value is usually rounded to a finite number of digits. After choosing a proper unit, the
 147 power value can be integer. So, it is assumed that both X_{ij} and X_{ji} are integer variables. In
 148 addition, it is assumed $X_{ij} = 0$ for $i=j$, and $\min(X_{ij}, X_{ji}) = 0$ for $i \neq j$.

149 **3.1. The case where power transmission loss is neglected**

150 Based on the above assumptions, if the power transmission loss between the system nodes
 151 can be neglected, the optimal power sharing problem can be represented by the following
 152 integer programming model

153
$$\text{Min } f = \sum_{i=1}^N \left\{ \text{Max} \left(d_i - \left(g_i + \sum_{j=1}^N X_{ji} - \sum_{j=1}^N X_{ij} \right), 0 \right) \right\}; \quad (1)$$

154 Subject to:

155
$$0 \leq X_{ij} \leq c_{ij}, \quad i, j = 1, 2, \dots \quad (2)$$

156
$$X_{ij} = 0, \quad i = j, \quad (3)$$

157
$$\min(X_{ij}, X_{ji}) = 0, \quad i \neq j, \quad (4)$$

158
$$g_i + \sum_{j=1}^N X_{ji} - \sum_{j=1}^N X_{ij} \geq 0, \quad (i=1,2,\dots) \quad (5)$$

159
$$X_{ij} \in Z, \quad (i, j = 1, 2, \dots) \quad (6)$$

160 where (1) is the objective function to be minimized, which represents the minimum total amount

161 of unsupplied demand for all the system nodes; $g_i + \sum_{j=1}^N X_{ji} - \sum_{j=1}^N X_{ij}$ is regarded as the

162 amount of power in node i after performance sharing. $d_i - \left(g_i + \sum_{j=1}^N X_{ji} - \sum_{j=1}^N X_{ij} \right)$ is

163 regarded as the difference between node i 's demand and capacity.

164 $Max \left(d_i - \left(g_i + \sum_{j=1}^N X_{ji} - \sum_{j=1}^N X_{ij} \right), 0 \right)$ represents the deficiency of node i after

165 performance sharing. (2) ensures that the power sharing amount between any two nodes does

166 not exceed the bandwidth of the transmission line between them. (3) ensures that there no power

167 transmitted from a node to itself. (4) ensures that the power transmission between any two

168 nodes happens for at most one direction. (5) ensures that the power supply of each node after

169 power redistribution is not negative, and (6) is assumed for simplicity. The model was solved

170 by the integer programming toolbox of MATLAB software.

171 **3.2. The case where power transmission loss is incorporated**

172 In the actual transmission process, power loss may occur due to the line resistance. Taking

173 into account the transmission loss, the objective function of the optimal power sharing problem

174 can be formulated as:

175
$$Min f = \sum_{i=1}^N \left\{ Max \left(d_i - \left(g_i + \sum_{j=1}^N (1 - K_{ji}) X_{ji} - \sum_{j=1}^N X_{ij} \right), 0 \right) \right\}; \quad (7)$$

176 where K_{ij} ($0 \leq K_{ij} \leq 1$) is the loss rate for power transmission from node i to node j and

177 $\sum_{j=1}^N (1 - K_{ji}) X_{ji}$ is the actual amount of power that can be transmitted to node i from all

178 other nodes after deducting the power transmission loss. For simplicity, it is assumed that

179 $K_{ij} = K_{ji}$ ($i, j = 1, 2, \dots$) which implies that the power transmission loss rate does not

180 depend on the direction of the transmission. $g_i + \sum_{j=1}^N (1 - K_{ji}) X_{ji} - \sum_{j=1}^N X_{ij}$ is regarded as the
 181 amount of power in node i after performance sharing. Eq. (7) is the objective function that
 182 needs minimizing and it represents the total amount of unsupplied demand for all the system
 183 nodes after performance sharing. The constraints are still as the same as Section 3.1, as
 184 formulated by Eqs (2)-(6). The model is solved by the integer programming toolbox of
 185 MATLAB software.

186 **3.3. The power sharing rate**

187 In order to know the effect of performance sharing in reducing the unsupplied demand, we
 188 introduce an index called Power Sharing Rate (PSR). In particular, when the allocation of
 189 generators is given and both the generators' capacity and the nodes' demand are also given, the
 190 unsupplied power demand f can be calculated. We assume that the unsupplied power
 191 demand before power sharing is f_{before} , and the minimum achievable unsupplied power
 192 demand after power sharing is f_{after} . The power sharing rate corresponding to a fixed power
 193 sharing policy is given by:

$$194 \quad PSR = \frac{f_{\text{before}} - f_{\text{after}}}{f_{\text{before}}}.$$

195 where $f_{\text{before}} > 0$. When $f_{\text{before}} = 0$, $PSR = 0$.

196 **4. The optimal allocation of generators**

197 As the capacity of each generator and the demand of each node in the network are random, it is
 198 essential to allocate the generators in an optimal way to minimize the expected unsupplied
 199 demand of the power grid system, by taking into account all the possible realizations of
 200 generators' capacity and the nodes' demand. Section 4.1 presents a UGF based approach to
 201 represent the unsupplied demand of the power grid systems, and Section 4.2 models the optimal
 202 generators allocation problem.

203 **4.1. Generator capacity distribution and node demand distribution**

204 Based on the system description in Section 2, this paper constructs a generator capacity
 205 distribution function and a node demand distribution function using UGF (Lisnianski, 2007;
 206 Lisnianski & Ding, 2009; Meena & Vasanthi, 2016; Khorshidi et al., 2016; Liu et al., 2017).
 207 The UGF approach is based on the definition of a u-function of multistate elements, which are
 208 of discrete random variables and composition operators over u-functions. It is a polynomial
 209 used to represent the distribution of discrete random variables, where the exponent represents
 210 the realization of the random variables and the coefficient represents the corresponding

211 probability of the realization. In particular, the UGF used to represent the capacity distribution
 212 of generator k is defined as

$$213 \quad u_k(z) = \sum_{l=1}^{L_k} p_k(l) z^{a_k(l)}, \quad (8)$$

214 where $a_k(l)$ is the l -th capacity realization of generator, and $p_k(l)$ is the corresponding
 215 probability of the capacity realization of the generator.

216 The demand distribution of the node i is defined as below

$$217 \quad u_{id}(z) = \sum_{l=1}^{W_i} q_i(l) z^{d_i(l)}, \quad (9)$$

218 where $d_i(l)$ is the demand realization of the node i , and $q_i(l)$ is the corresponding
 219 probability of the demand realization of the node i .

220 For each node i , the total amount of capacity equals to the summation of the capacity for
 221 all the generators allocated into node i . Therefore, its capacity distribution can be obtained
 222 from the capacity distribution of all the generators belonging to the node i as

$$223 \quad u_{ig}(z) = \prod_{k=1}^{n(i)} u_{i(k)}(z) = \prod_{k=1}^{n(i)} \left(\sum_{l=1}^{L_{ig}} p_i(l) z^{g_i(l)} \right), \quad (10)$$

224 where $n(i)$ is the number of generators allocated into node i , $i(k)$ is the index for the k
 225 -th generator allocated to node i , $g_i(l)$ is the capacity realization of the node i , and
 226 $p_i(l)$ is the corresponding probability of the capacity realization of the node i . Note that, in
 227 case where node i does not contain any generators, the capacity distribution of node i can
 228 be expressed by the UGF $u_{ig}(z) = z^0$.

229 **4.2. Representation of the generators' allocation**

230 The generator allocation problem can be considered as a problem of partitioning M
 231 generators into N mutually disjoint subsets $F_i (1 \leq i \leq N)$ such that

$$232 \quad \bigcup_{i=1}^N \quad (11)$$

233 and

$$234 \quad F_i \cap F_j = \phi, \quad i \neq j. \quad (12)$$

235 The partition of the set F can be represented by the vector $H = \{h(k), 1 \leq k \leq M\}$,
236 which denotes that any generator k is assigned to the $h(k)$ th node. The cardinality of each
237 subset $F_i (1 \leq i \leq N)$ can be easily obtained as

$$238 \quad n(i) = |F_i| = \sum_{j=1}^M 1(h(j) = i), \quad (i = 1, 2, \dots, N). \quad (13)$$

239 Where, “1” is an indicator function, satisfying $1(\text{TRUE})=1$ and $1(\text{FALSE})=0$.

240 **4.3. Optimal allocation of generators**

241 For fixed allocation of generators, the expected unsupplied demand of the power grid system
242 can be obtained by considering all the possible realizations of the generators' capacity and the
243 nodes' demand. It can be seen from Section 3, Subsections 4.1 and 4.2 that the numbers of states
244 for the N nodes are L_{1g}, L_{2g}, \dots , respectively. The number of demand realizations for the

245 node i is W_i . Therefore, there are $V = \prod_{i=1}^N L_{ig} W_i$ different realizations of the generators'

246 capacity and the nodes' demand, where the unsupplied demand needs to be minimized under
247 each realization using the procedures proposed in Section 3. Denote the minimized unsupplied
248 demand for each realization t as f_t and the corresponding probability for the realization as

249 $P_{f_t} (t = 1, 2, \dots)$. Finally, the expected value of the minimum unsupplied power demand of
250 the system is obtained by

$$251 \quad E(f_t) = \sum_{t=1}^V P_{f_t} f_t. \quad (14)$$

252 In addition, the expected power sharing rate can be obtained by taking into account all the
253 possible realizations of the generators' capacity and the nodes' demand as:

$$254 \quad E(\text{PSR}) = \sum_{t=1}^V P_{f_t} \text{PSR}_t,$$

255 where PSR_t is the power sharing rate under the optimal sharing policy for the realization t .

256 Based on the above description, the optimization problem is to find vector
 257 $H = \{h(k), 1 \leq k \leq M\}$ that minimizes the expected unsupplied demand of the system. That
 258 is

$$259 \quad H_{optimal} = \arg \min E[f_t(H)];$$

260 (15)

261 Subject to

$$262 \quad H = \{h(k), 1 \leq k \leq M\}, \quad (16)$$

$$263 \quad h(k) = 1, \dots \quad \dots \quad (17)$$

264 5. Optimization technique

265 To solve the optimal generator allocation policy in order to minimize the objective function (15)
 266 is a difficult combinatorial optimization problem. Though the exact solution is possible to find
 267 with enumeration technique, it may be too time consuming. In practice, the quality of the
 268 solution is of concern, thus a typical way for such combinatorial optimization problem is to
 269 employ a heuristic to find the near optimal solution. In particular, this paper adopts a Hybridized
 270 Particle Swarm optimization (HPSO) algorithm that combines Particle Swarm optimization
 271 algorithm (PSO) with other algorithms (Chen et al., 2014; Jamrus et al., 2018; Zhang et al.,
 272 2016).

273 PSO is inspired from the behaviour characteristics of the biological population and is used
 274 to solve optimization problems in some research fields. In PSO, each particle represents a
 275 potential solution of the problem and each particle corresponds to a fitness value determined
 276 by fitness function (Wang & Tang, 2011). The velocity of the particle determines the direction
 277 and distance of the particle moving. The velocity is dynamically adjusted with the movement
 278 experience of the particle itself and other particles, so as to realize the optimization of the
 279 individual in the solvable space. To apply the described swarm optimization technique to a
 280 particular problem, the key actually lies in linking the solution to the fitness function of the
 281 solution. In our case, the fitness function of each solution H is defined in Eq. (14) in Section 4,
 282 which again needs to use the result from Eq. (7) in Section 3.

283 PSO initializes a group of particles in the feasible solution space. Each particle represents
 284 a potential optimal solution of the extremum optimization problem. The characteristics of the
 285 particle are represented by the position, velocity and fitness values. The fitness function
 286 represents the objective function, which indicates the advantages and disadvantages of the
 287 particles. Particles move in the solution space and update individual positions by tracking
 288 individual extremum and group extremum.

289 Although PSO is simple and can converge quickly, with the increasing number of
 290 iterations, the convergence of the population is concentrated, and the examples are more and
 291 more similar, which may not jump out around the local optimal solution. To overcome the

292 shortcoming, the HPSO introduces crossover and mutation operations from genetic algorithms,
 293 and searches for the optimal solution by crossing the individual extreme value with the
 294 population extreme value and mutation of the particle itself (Gong et al., 2016; Shi et al., 2003).

295 **6. Illustrative examples**

296 Consider the allocation of 5 generators into the 4 nodes in the following illustrative example.
 297 The topology of the system is as shown in Fig 2. The capacity distribution for each generator
 298 is as listed in Table 1, where the capacity distribution of each generator has three states.

299 Fig 2. The topology diagram of power grid with four nodes

300 Table 1. The capacity distribution for each generator

301 The demand distribution for each node is as listed in Table 2. There are three states for
 302 each node's demand distribution.

303 Table 2. The demand distribution for each node

304 **6.1. Fixed allocation**

305 Suppose the allocation vector $H = \{h(k), 1 \leq k \leq 5\} = \{2, 3, 4, 1, 1\}$. Through the definition
 306 of H , it is easy to know that the sets of generators allocated to the 4 nodes are $\{4,5\}$, $\{1\}$, $\{2\}$,
 307 and $\{3\}$, respectively. That is, it is a generator allocation policy.

308 According to the given generator allocation policy, the capacity distribution of the four
 309 nodes based on the formulas (8) and (10) are given respectively by

$$310 \quad u_{1g}(z) = (0.7z^4 + 0.2z^3 + 0.1z^0)(0.6z^4 + 0.2z^3 + 0.2z^0) \\
 = 0.42z^8 + 0.26z^7 + 0.04z^6 + 0.2z^4 + 0.06z^3 + 0.02z^0,$$

$$311 \quad u_{2g}(z) = 0.8z^4 + 0.1z^3 + 0.1z^0,$$

$$312 \quad u_{3g}(z) = 0.8z^4 + 0.1z^3 + 0.1z^0,$$

313 and

$$314 \quad u_{4g}(z) = 0.7z^4 + 0.2z^3 + 0.1z^0.$$

315 It is known from the above discussion that there are six states for the capacity distribution
 316 of the first node, and there are three states for the capacity distribution of the other three nodes.
 317 In addition, the demand for each node has three possible realizations. Thus, there are
 318 $3^7 \times 6 = 13122$ different realizations for the generators' capacity and the nodes' demand.

319 Take one of these cases as an example, as shown in Table 3. The probability of this case is
 320 $0.26 \times 0.1 \times 0.1 \times 0.7 \times 0.1 \times 0.1 \times 0.6 \times 0.15 = 1.638 \times 10^{-6}$. The optimal power sharing policy
 321 is discussed for the case where the power transmission loss can be neglected and the case where

322 the power transmission loss cannot be neglected.

323 Table 3. A specific case for illustration

324 6.1.1. The case where power transmission loss is neglected

325 (1) The case where the transmission bandwidth between nodes is the same

326 Assuming that the transmission bandwidth between nodes i and j is $c_{ij} = 0$. The
327 minimum value of the unsupplied power demand that can be achieved for the case presented in
328 Table 3 is solved to be 4. The contribution to the expected unsupplied power demand is
329 $1.638 \times 10^{-6} \times 4 = 6.552 \times 10^{-6}$. The PSR for this case is 0.

330 If the transmission bandwidth between nodes i and j is $c_{ij} = 1$, the optimal solution,
331 the minimum value of the unsupplied power demand that can be achieved for the case presented
332 in Table 3, the contribution to the expected unsupplied power demand, and the PSR are
333 respectively $x_{12} = x_{13} = x_{14} = x_{23} = x_{34} = x_{42} = 1$, $x_{21} = x_{24} = x_{31} = x_{32} = x_{41} = x_{43} = 0$, 1,
334 $1.638 \times 10^{-6} \times 1 = 1.638 \times 10^{-6}$, and $3/4 = 0.75$. If the transmission bandwidth between nodes i
335 and j is $c_{ij} = 2$, they become $x_{12} = 2$,

336 $x_{21} = x_{23} = x_{31} = x_{32} = x_{34} = x_{43} = x_{41} = x_{24} = x_{42} = 0$, $x_{13} = x_{14} = 1$, 0, $1.638 \times 10^{-6} \times 0 = 0$
337 and $4/4 = 1$, respectively.

338 It can be seen that with the increase of the transmission bandwidth between nodes i and
339 j , the minimum value of the unsupplied power demand becomes smaller and the PSR increases.

340 Considering all the realizations of the generators' capacity and the nodes' power demand,
341 the expected values of the minimum unsupplied power demand of the system $E(f_t)$ and the
342 PSR are obtained as below.

343 When the transmission bandwidth between nodes i and j is $c_{ij} = 0$, the expected
344 value of the minimum unsupplied power demand of the system $E(f_t)$ is 2.1050 and the
345 expected PSR is 0. When the transmission bandwidth between nodes i and j is $c_{ij} = 1$,
346 these values become 1.0156 and 0.5175, respectively. When the transmission bandwidth
347 between nodes i and j is $c_{ij} = 2$, these values become 0.8955 and 0.5746, respectively.

348 It can be seen that with the increase of the transmission bandwidth between nodes i and
349 j , the expected value of the minimum unsupplied power demand of the system $E(f_t)$

350 becomes smaller and the expected PSR increases.

351 (2) The case where the transmission bandwidth between nodes is different

352 In order to examine the importance of different transmission lines, we investigate the case
353 where the bandwidths of the transmission lines differ. Specifically, it is assumed that the
354 bandwidths of the transmission lines are as shown in Table 4. The corresponding expected value
355 of the minimum unsupplied demand together with the expected PSR are also shown for each
356 setting of the bandwidths. To be consistent with some other papers in the reliability field (Song
357 & Schmeiser, 2009), the expected value of minimum unsupplied demand is accurate to the last
358 four decimal points. For guaranteeing the quality of our obtained optimal solution, for each
359 setting, the optimization code is run 5 times and consistent results are observed.

360 Table 4. Results for conditions of different transmission bandwidths

361 From Table 4, we can see that the transmission line between nodes 1 and 4 and the
362 transmission line between nodes 1 and 3 are more important than other lines. Actually, two
363 generators instead of one are allocated to node 1 which makes it more likely to have excess. On
364 the other hand, the demand of node 3 and node 4 are relatively higher. Therefore, the power
365 transmission from node 1 to node 3 and node 4 is more important than other lines. To further
366 confirm the finding, the case presented in Table 5 is examined, which also implies the
367 importance of power transmission from node 1 to node 3 and node 4.

368 Table 5. Results for conditions of different transmission bandwidths

369 *6.1.2. The case where power transmission loss is incorporated*

370 (1) The case where the transmission bandwidth between nodes is the same

371 Assuming loss rate $K_{ij} = 0.05$, where the capacity and the demand of the nodes are still as

372 shown in Table 3. When the transmission bandwidth between nodes i and j is $c_{ij} = 0$, it

373 is not allowed to transmit power, so there is no power loss. The minimum value of the
374 unsupplied power demand is 4, and the contribution to the expected unsupplied power demand
375 is $1.638 \times 10^{-6} \times 4 = 6.552 \times 10^{-6}$. The PSR for this case is 0. When the transmission

376 bandwidth between nodes i and j is $c_{ij} = 1$, the optimal solution is obtained as

377 $x_{12} = x_{13} = x_{14} = 1$ and $x_{21} = x_{23} = x_{24} = x_{31} = x_{32} = x_{34} = x_{41} = x_{42} = x_{43} = 0$, respectively.

378 The minimum value of the unsupplied power demand is 1.15, and the contribution to the
379 expected unsupplied power demand is $1.6380 \times 10^{-6} \times 1.15 = 1.8837 \times 10^{-6}$.

380 The PSR for this case is $2.85/4 = 0.7125$. When the transmission bandwidth between nodes

381 i and j is $c_{ij} = 2$, the optimal solution is obtained as $x_{12} = 2$, $x_{13} = x_{14} = 1$, and

382 $x_{41} = x_{24} = x_{42} = x_{21} = x_{23} = x_{31} = x_{32} = x_{34} = x_{43} = 0$. The minimum value of the unsupplied

383 power demand is 0.2, and the contribution to the expected unsupplied power demand is

384 $1.6380 \times 10^{-6} \times 0.2 = 3.276 \times 10^{-7}$. The PSR for this case is $3.8/4=0.95$. It can be seen that with
385 the increase of the transmission bandwidth between nodes i and j , the minimum value of
386 the unsupplied power demand becomes smaller and the PSR increases.

387 Combined with all cases, the expected value of minimum unsupplied power demand of
388 the system $E(f_t)$ and the PSR can be obtained. When the transmission bandwidth between
389 nodes i and j is $c_{ij} = 0$, the expected value of minimum unsupplied power demand of the
390 system $E(f_t)$ is 2.1050 and the expected PSR is 0. When the transmission bandwidth
391 between nodes i and j is $c_{ij} = 1$, the expected value of the minimum unsupplied power
392 demand of the system $E(f_t)$ is 1.0774 and the expected PSR is 0.4882. When the
393 transmission bandwidth between nodes i and j is $c_{ij} = 2$, the expected value of the
394 minimum unsupplied power demand of the system $E(f_t)$ is 0.9736 and the expected PSR is
395 0.5375.

396 The results of the optimal solution with or without power loss are listed in Table 6. It is
397 obvious that the transmission bandwidth and the power transmission loss are positively
398 proportional. Actually, more power can be transmitted when the bandwidth is bigger, which
399 also causes more power to be lost through transmission, which indeed results in more
400 unsupplied demand.

401 Table 6. Comparison of solution with or without power loss

402 (2) The case where the transmission bandwidth between nodes is different

403 Similarly, the results for the case where transmission bandwidth between different nodes are
404 obtained, as shown in Table 7, which also shows the importance of power transmission from
405 node 1 to node 3 and node 4.

406 Table 7. Results for conditions of different transmission bandwidths

407 **6.2. Optimal allocation**

408 In order to minimize the power demand of the system, we analyze the optimal allocation
409 policy of generators below. The optimal allocation policy of generator is obtained by the
410 hybridized particle swarm optimization (HPSO) algorithm. The hybridized particle swarm
411 optimization algorithm in this paper combines particle swarm optimization algorithm and
412 genetic algorithm. The algorithm flow chart is as shown in Fig 3.

413 Fig 3. Flow chart of hybridized particle swarm optimization algorithm

414 The algorithm is realized by programming in MATLAB. The size of the selected
415 population is 200, the number of times of evolution is 100. The computer is a Windows 10
416 operating system, i7-8550u CPU, 4.0GHZ, Intel Quad-core processor, 16GB memory.

417 *6.2.1. The optimal allocation policy of generators without considering power transmission loss*

418 (1) The case where the transmission bandwidth between nodes is the same

419 When the transmission bandwidth between nodes i and j is $c_{ij} = 0$, the optimal allocation
420 policy of generator is $\{2\}, \{1\}, \{4,5\}, \{3\}$. The minimum expected unsupplied demand value
421 of the system is 2.0450 and the expected PSR is 0. When the transmission bandwidth between
422 nodes i and j is $c_{ij} = 1$, the optimal allocation policy of generator is $\{1\}, \{2\}, \{3,5\}, \{4\}$.
423 The minimum expected unsupplied demand value of the system is 1.0141 and the expected PSR
424 is 0.5041. When the transmission bandwidth between nodes i and j is $c_{ij} = 2$, the optimal
425 allocation policy of generator is $\{1\}, \{2\}, \{4,5\}, \{3\}$. The minimum expected unsupplied
426 demand value of the system is 0.8954 and the expected PSR is 0.5622. It can be seen that with
427 the increase of the transmission bandwidth between nodes i and j , the minimum expected
428 unsupplied demand value of the system becomes smaller and the expected PSR increases.

429 (2) The case where the transmission bandwidth between nodes is different

430 Table 8 presents the results for the case where the transmission bandwidths of different
431 transmission lines are different. It can be seen that the transmission line between nodes 3 and
432 node 4 is more important in this case. Actually, the demand of node 3 and node 4 are relatively
433 higher, which makes either of them more likely to experience deficit given the optimal
434 allocation. Therefore, the power transmission between them is essential.

435 Table 8. Results for conditions of different transmission bandwidths

436 *6.2.2. The optimal allocation policy of generators considering power transmission loss*

437 (1) The case where the transmission bandwidth between nodes is the same

438 Still consider the power transmission loss rate $K_{ij} = 0.05$. When the transmission bandwidth
439 between nodes i and j is $c_{ij} = 0$, the optimal allocation policy of generators, the
440 minimum expected unsupplied demand value of the system, and the expected PSR are
441 respectively $\{2\}, \{1\}, \{4,5\}, \{3\}, 2.0450$ and 0. When the transmission bandwidth between
442 nodes i and j is $c_{ij} = 1$, they become $\{2\}, \{1\}, \{3,5\}, \{4\}, 1.0774$, and 0.4732. When the
443 transmission bandwidth between nodes i and j is $c_{ij} = 2$, they become $\{1\}, \{2\}, \{4,5\},$
444 $\{3\}, 0.9724$, and 0.5245. It can be seen that with the increase of the transmission bandwidth
445 between nodes i and j , the minimum expected unsupplied demand value of the system

446 becomes smaller and the expected PSR increases. Under the same bandwidth, the minimum
 447 expected unsupplied demand value of the system for the case with power loss is larger than that
 448 of the system for the case without power loss, and the expected PSR for the case with power
 449 loss is smaller than that of the system for the case without power loss.

450 (2) The case where the transmission bandwidth between nodes is different

451 The results for the cases where the transmission widths are different are presented in Table 9.
 452 Similar as Table 8, it also shows that the transmission line between nodes 3 and 4 is more
 453 important.

454 Table 9. Results for conditions of different transmission bandwidths

455 7. Case study

456 Decentralized wind power generation projects are being undertaking in different areas of China,
 457 in particular, in western China. Imagine that there is an area with six cities, each of which is
 458 connected with a wind power station. Suppose that the power stations for different cities are
 459 connected with transmission lines, as shown in Fig. 4. The power transmission bandwidth
 460 between nodes is the same as the transmission lines of the same material and specification are
 461 used. Due to increasing power consumption of this area, it is intended to allocate eight wind
 462 power generators. The capacity distribution of the wind power generators and the distribution
 463 of demand for wind power for the six cities are as shown in Tables 10 and 11, respectively:

464 Fig 4. The topology diagram of power grid with six cities

465 Table 10. The capacity distribution for each generator

466 Table 11. The demand distribution for each node

467 7.1. Fixed allocation

468 Suppose the allocation is fixed and the allocation vector is
 469 $H = \{h(k), 1 \leq k \leq 8\} = \{3, 3, 6, 2, 1, 4, 5, 5\}$. Through the definition of H , it is easy to
 470 know that the sets of generators allocated to the 6 nodes are $\{5\}$, $\{4\}$, $\{1, 2\}$, $\{6\}$, $\{7, 8\}$ and
 471 $\{3\}$, respectively.

472 According to the given generator allocation policy, the capacity distribution of the six
 473 nodes based on the formulas (8) and (10) are given respectively by

$$474 \quad u_{1g}(z) = 0.7z^4 + 0.2z^3 + 0.1z^0,$$

$$475 \quad u_{2g}(z) = 0.8z^4 + 0.1z^3 + 0.1z^0,$$

$$476 \quad u_{3g}(z) = (0.8z^4 + 0.1z^3 + 0.1z^0)(0.8z^4 + 0.1z^3 + 0.1z^0) \\ = 0.64z^8 + 0.16z^7 + 0.01z^6 + 0.16z^4 + 0.02z^3 + 0.01z^0,$$

$$477 \quad u_{4g}(z) = 0.7z^4 + 0.2z^3 + 0.1z^0,$$

478
$$u_{5g}(z) = (0.7z^4 + 0.2z^3 + 0.1z^0)(0.6z^4 + 0.2z^3 + 0.2z^0)$$

$$= 0.42z^8 + 0.26z^7 + 0.04z^6 + 0.2z^4 + 0.06z^3 + 0.02z^0,$$

479 and

480
$$u_{6g}(z) = 0.8z^4 + 0.1z^3 + 0.1z^0.$$

481 It is known from the above discussion that there are 2125764 ($=3^{10} \times 6^2$) different
482 realizations for the generators' capacity and the nodes' demand.

483 Considering all the realizations of the generators' capacity and the nodes' power demand,
484 we discuss the optimal power sharing policy for the case where the power transmission loss can
485 be neglected and the case where the power transmission loss cannot be neglected.

486 *7.1.1. The case where power transmission loss is neglected*

487 When power transmission loss is neglected, we consider all the realizations of the generators'
488 capacity and the nodes' power demand, respectively. The expected values of the minimum
489 unsupplied power demand of the system $E(f_i)$ and the PSR are obtained as below.

490 When the transmission bandwidth between nodes i and j is $c_{ij} = 0$, the expected
491 value of the minimum unsupplied power demand of the system $E(f_i)$ is 2.5360 and the
492 expected PSR is 0. When the transmission bandwidth between nodes i and j is $c_{ij} = 1$,
493 these values become 1.1963 and 0.5283, respectively. When the transmission bandwidth
494 between nodes i and j is $c_{ij} = 2$, these values become 0.6964 and 0.7254, respectively .

495 It can be seen that with the increase of the transmission bandwidth between nodes i and
496 j , the expected value of the minimum unsupplied power demand of the system $E(f_i)$
497 becomes smaller and the expected PSR increases.

498 *7.1.2. The case where power transmission loss is incorporated*

499 Assume that loss rate is $K_{ij} = 0.05$. When the transmission bandwidth between nodes
500 i and j is $c_{ij} = 0$, the expected value of minimum unsupplied power demand of the system
501 $E(f_i)$ is 2.1657 and the expected PSR is 0. When the transmission bandwidth between nodes
502 i and j is $c_{ij} = 1$, the expected value of the minimum unsupplied power demand of the
503 system $E(f_i)$ is 1.2740 and the expected PSR is 0.4117. When the transmission bandwidth
504 between nodes i and j is $c_{ij} = 2$, the expected value of the minimum unsupplied power

505 demand of the system $E(f_i)$ is 0.8506 and the expected PSR is 0.6072.

506 It can be seen that with the increase of the transmission bandwidth between nodes i and
507 j , the minimum expected unsupplied demand value of the system becomes smaller and the
508 expected PSR increases. Under the same bandwidth, the minimum expected unsupplied
509 demand value of the system for the case with power loss is larger than that of the system for
510 the case without power loss, and the expected PSR for the case with power loss is smaller than
511 that of the system for the case without power loss.

512 **7.2. Optimal allocation**

513 In order to minimize the power demand of the system, we analyze the optimal allocation
514 policy of generators below.

515 *7.2.1. The optimal allocation policy of generators without considering power transmission loss*

516 If power transmission loss is neglected, the optimal allocation vector
517 $H = \{h(k), 1 \leq k \leq 8\} = \{5, 5, 6, 1, 2, 4, 3, 3\}$. That is, the optimal allocation policy of
518 generator is $\{4\}$, $\{5\}$, $\{7, 8\}$, $\{6\}$, $\{1, 2\}$, and $\{3\}$. If the transmission bandwidth between nodes
519 i and j is $c_{ij} = 0$, the minimum expected unsupplied demand value of the system is 2.4270
520 and the expected PSR is 0. If the transmission bandwidth between nodes i and j is $c_{ij} = 1$,
521 the minimum expected unsupplied demand value of the system is 1.0094 and the expected PSR
522 is 0.5841. If the transmission bandwidth between nodes i and j is $c_{ij} = 2$, the minimum
523 expected unsupplied demand value of the system is 0.5940 and the expected PSR is 0.7553. It
524 can be seen that with the increase of the transmission bandwidth between nodes i and j , the
525 minimum expected unsupplied demand value of the system decreases and the expected PSR
526 increases.

527 *7.2.2. The optimal allocation policy of generators considering power transmission loss*

528 Suppose that the power transmission loss rate $K_{ij} = 0.05$. The optimal allocation policy of
529 generators is $\{2\}$, $\{6\}$, $\{5, 8\}$, $\{7\}$, $\{3, 4\}$, $\{1\}$. If the transmission bandwidth between nodes
530 i and j is $c_{ij} = 0$, the minimum expected unsupplied demand value of the system is 2.4270
531 and the expected PSR is 0. If the transmission bandwidth between nodes i and j is $c_{ij} = 1$,
532 the minimum expected unsupplied demand value of the system, and the expected PSR are
533 respectively 1.0753 and 0.5569. If the transmission bandwidth between nodes i and j is
534 $c_{ij} = 2$, they become 0.7284, and 0.6999.

535 Similarly, the same conclusion can be obtained that with the increase of the transmission
536 bandwidth between nodes i and j , the minimum expected unsupplied demand value of the
537 system becomes smaller and the expected PSR increases. Under the same bandwidth, the
538 minimum expected unsupplied demand value of the system for the case with power loss is
539 larger than that of the system for the case without power loss, and the expected PSR for the
540 case with power loss is smaller than that of the system for the case without power loss.

541 **8. Conclusions**

542 Performance sharing is a method of improving the reliability of systems. However, the existing
543 works are restricted to systems of simple structures. Differently, this paper considers a complex
544 power grid system with performance sharing mechanisms, where each node of the power
545 network can be distributed with power generators. For fixed generators' allocation, procedures
546 are proposed to evaluate the expected system unsupplied demand. After that, a hybridized
547 particle swarm optimization (HPSO) algorithm is used to find out the optimal allocation policy
548 of generators that minimizes the expected unsupplied demand of the system. Both the case
549 where power transmission loss can and cannot be neglected and the case where the power
550 transmission loss cannot be neglected are considered. Examples are presented to illustrate
551 show the applications of the proposed procedures.

552 There are some limitations in this paper. First, this paper studies a static optimization
553 problem where the demand and the capacity are random variables. In practice, the demand and
554 the capacity are changing from time to time, and thus the dynamic system configuration and
555 dynamic performance sharing may be of interest. In addition, the demands of different nodes
556 are assumed to be independent, so as to the capacities of different capacities. The possible
557 dependence between demands, between capacities, and even between demand and capacities
558 can be investigated in the future.

559 Besides the above mentioned limitations, this work can also be extended in the following
560 ways. In this paper, only the generators' allocation is studied. In the future, policies of
561 maintenance and protection of the generators can be incorporated as well. For example, each
562 generator may be subject to internal failures and external attacks. The availability of each
563 generator depends on the preventive maintenance interval and the protection effort on each
564 generator. Besides generators' allocation, one may also be concerned with the optimal
565 maintenance and protection strategy which minimizes the expected system cost per unit time
566 (Peng et al., 2012; Peng et al., 2014). In addition, this paper only considers the distribution of
567 electricity. In the future, the integrated electricity and gas distribution system can be studied.
568 Another possible direction is to incorporate the framework the effects of cascading failures by
569 physical and cyber attacks.

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666 Table 1. The capacity distribution for each generator

Generator/ capacity	4	3	0
a_1	0.8	0.1	0.1
a_2	0.8	0.1	0.1
a_3	0.7	0.2	0.1
a_4	0.7	0.2	0.1
a_5	0.6	0.2	0.2

667 Table 2. The demand distribution for each node

Node/demand	3	4	5
d_1	0.1	0.8	0.1
d_2	0.1	0.8	0.1

d_3	0.2	0.6	0.2
d_4	0.15	0.7	0.15

668

669

Table 3. A specific case for illustration

Node	capacity (g_i)	probability	demand (d_i)	probability
n_1	7	0.26	3	0.1
n_2	3	0.1	5	0.1
n_3	3	0.1	4	0.6
n_4	4	0.7	5	0.15

670

Table 4. Results for conditions of different transmission bandwidths

Bandwidths of transmission lines	The expected value of minimum unsupplied demand	The expected PSR
$c_{12} = c_{21} = 1$, other $c_{ij} = 0$	1.8977	0.0985
$c_{13} = c_{31} = 1$, other $c_{ij} = 0$	1.8384	0.1267
$c_{14} = c_{41} = 1$, other $c_{ij} = 0$	1.8176	0.1365
$c_{23} = c_{32} = 1$, other $c_{ij} = 0$	2.0346	0.0334
$c_{24} = c_{42} = 1$, other $c_{ij} = 0$	2.0466	0.0277
$c_{34} = c_{43} = 1$, other $c_{ij} = 0$	2.0093	0.0455

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Table 5. Results for conditions of different transmission bandwidths

Bandwidths of transmission lines	The expected value of minimum unsupplied demand	The expected PSR
$c_{12} = c_{21} = 2, \text{ other } c_{ij} = 0$	1.8189	0.1359
$c_{13} = c_{31} = 2, \text{ other } c_{ij} = 0$	1.7524	0.1675
$c_{14} = c_{41} = 2, \text{ other } c_{ij} = 0$	1.7245	0.1808
$c_{23} = c_{32} = 2, \text{ other } c_{ij} = 0$	2.0346	0.0334
$c_{24} = c_{42} = 2, \text{ other } c_{ij} = 0$	2.0466	0.0277
$c_{34} = c_{43} = 2, \text{ other } c_{ij} = 0$	2.0093	0.0455

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Table 6. Comparison of solution with or without power loss

	Bandwidth (c_{ij})	No loss	Loss
The minimum value of the unsupplied power demand for the case in Table 3	0	4	4
	1	1	1.15
	2	0	0.2
The expected value of minimum unsupplied power demand of the system	0	2.1050	2.1050
	1	1.0156	1.0774
	2	0.8955	0.9736

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Table 7. Results for conditions of different transmission bandwidths

Bandwidths of transmission lines	The expected value of minimum unsupplied demand	The expected PSR
$c_{12} = c_{21} = 1, \text{ other } c_{ij} = 0$	1.9080	0.0936
$c_{13} = c_{31} = 1, \text{ other } c_{ij} = 0$	1.8517	0.1203

$c_{14} = c_{41} = 1, \text{ other } c_{ij} = 0$	1.8320	0.1297
$c_{23} = c_{32} = 1, \text{ other } c_{ij} = 0$	2.0381	0.0318
$c_{24} = c_{42} = 1, \text{ other } c_{ij} = 0$	2.0496	0.0263
$c_{34} = c_{43} = 1, \text{ other } c_{ij} = 0$	2.0141	0.0432
$c_{12} = c_{21} = 2, \text{ other } c_{ij} = 0$	1.8275	0.1318
$c_{13} = c_{31} = 2, \text{ other } c_{ij} = 0$	1.7622	0.1629
$c_{14} = c_{41} = 2, \text{ other } c_{ij} = 0$	1.7348	0.1759
$c_{23} = c_{32} = 2, \text{ other } c_{ij} = 0$	2.0381	0.0318
$c_{24} = c_{42} = 2, \text{ other } c_{ij} = 0$	2.0496	0.0263
$c_{34} = c_{43} = 2, \text{ other } c_{ij} = 0$	2.0141	0.0432

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Table 8. Results for conditions of different transmission bandwidths

Bandwidths of transmission lines	The expected value of minimum unsupplied demand	The expected PSR
$c_{12} = c_{21} = 1, \text{ other } c_{ij} = 0$	2.0018	0.0211
$c_{13} = c_{31} = 1, \text{ other } c_{ij} = 0$	1.8312	0.1046
$c_{14} = c_{41} = 1, \text{ other } c_{ij} = 0$	1.9866	0.0285
$c_{23} = c_{32} = 1, \text{ other } c_{ij} = 0$	1.8312	0.1046
$c_{24} = c_{42} = 1, \text{ other } c_{ij} = 0$	1.9866	0.0285

$c_{34} = c_{43} = 1, \text{ other } c_{ij} = 0$	1.7487	0.1499
$c_{12} = c_{21} = 2, \text{ other } c_{ij} = 0$	2.0018	0.0211
$c_{13} = c_{31} = 2, \text{ other } c_{ij} = 0$	1.7528	0.1429
$c_{14} = c_{41} = 2, \text{ other } c_{ij} = 0$	1.9866	0.0285
$c_{23} = c_{32} = 2, \text{ other } c_{ij} = 0$	1.7528	0.1429
$c_{24} = c_{42} = 2, \text{ other } c_{ij} = 0$	1.9866	0.0285
$c_{34} = c_{43} = 2, \text{ other } c_{ij} = 0$	1.6561	0.1902

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Table 9. Results for conditions of different transmission bandwidths

Bandwidths of transmission lines	The expected value of minimum unsupplied demand	The expected PSR
$c_{12} = c_{21} = 1, \text{ other } c_{ij} = 0$	2.0040	0.0201
$c_{13} = c_{31} = 1, \text{ other } c_{ij} = 0$	1.8419	0.0993
$c_{14} = c_{41} = 1, \text{ other } c_{ij} = 0$	1.9896	0.0272
$c_{23} = c_{32} = 1, \text{ other } c_{ij} = 0$	1.8419	0.0993
$c_{24} = c_{42} = 1, \text{ other } c_{ij} = 0$	1.9896	0.0272
$c_{34} = c_{43} = 1, \text{ other } c_{ij} = 0$	1.7635	0.1377
$c_{12} = c_{21} = 2, \text{ other } c_{ij} = 0$	2.0040	0.0201
$c_{13} = c_{31} = 2, \text{ other } c_{ij} = 0$	1.7618	0.1385

$c_{14} = c_{41} = 2, \text{ other } c_{ij} = 0$	1.9896	0.0271
$c_{23} = c_{32} = 2, \text{ other } c_{ij} = 0$	1.7618	0.1385
$c_{24} = c_{42} = 2, \text{ other } c_{ij} = 0$	1.9896	0.0272
$c_{34} = c_{43} = 2, \text{ other } c_{ij} = 0$	1.6668	0.1849

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Table 10. The capacity distribution for each generator

Generator/ capacity	4	3	0
a_1	0.8	0.1	0.1
a_2	0.8	0.1	0.1
a_3	0.8	0.1	0.1
a_4	0.8	0.1	0.1
a_5	0.7	0.2	0.1
a_6	0.7	0.2	0.1
a_7	0.7	0.2	0.1
a_8	0.6	0.2	0.2

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Table 11. The demand distribution for each node

Node	Demand 1	Demand 2	Demand 3
	2	3	4
d_1	0.1	0.8	0.1

	2	3	4
d_2	0.15	0.7	0.15
d_3	4	5	6
	0.1	0.8	0.1
d_4	2	3	4
	0.1	0.8	0.1
d_5	4	6	8
	0.15	0.7	0.15
d_6	3	4	5
	0.1	0.8	0.1

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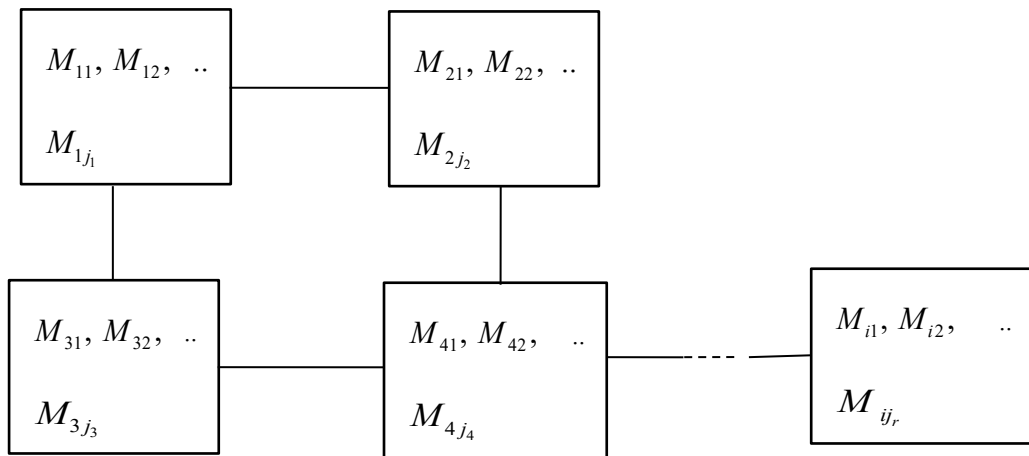
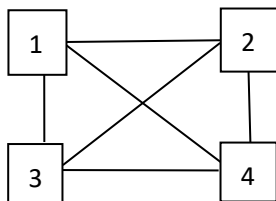


Fig 1. The power grid system structure diagram

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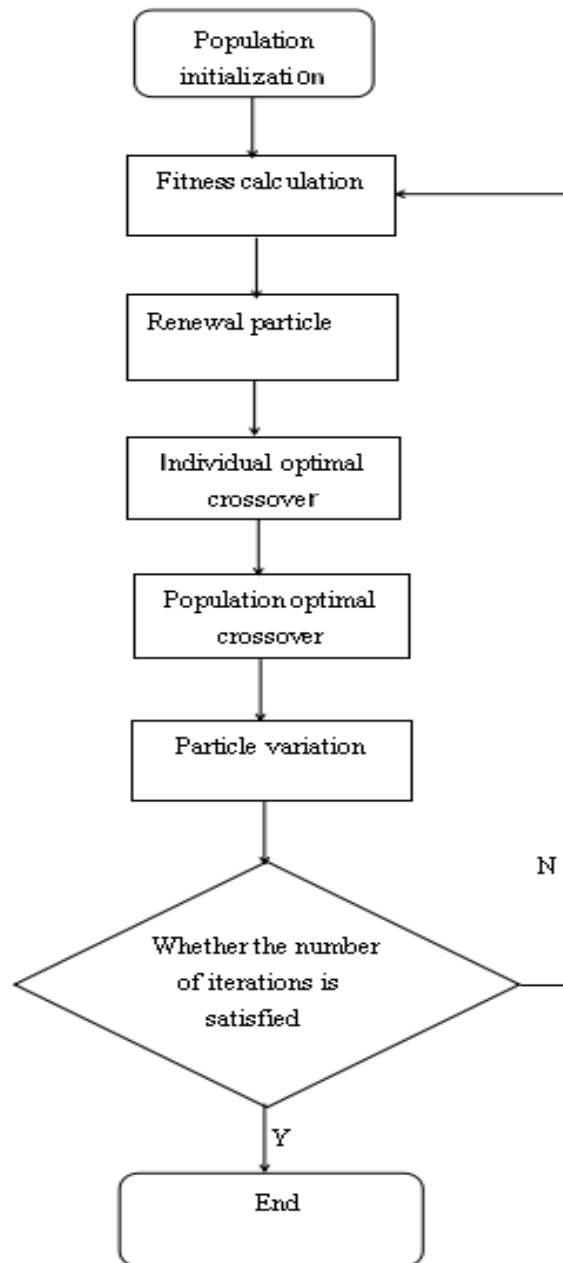
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Fig 2. The topology diagram of power grid with four nodes



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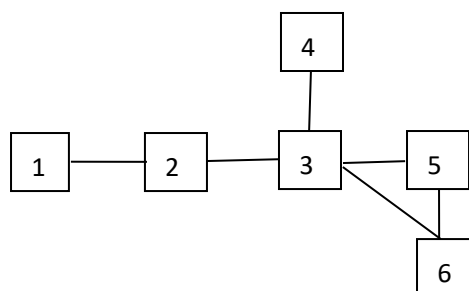
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Fig 3. Flow chart of hybridized particle swarm optimization algorithm

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Fig 4. The topology diagram of power grid with six cities