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Abstract

Until now there has been no support for specifying and enforcing contracts within a lazy functional program. That is a shame, because contracts consist of pre- and post-conditions for functions that go beyond the standard static types. This paper presents the design and implementation of a small, easy-to-use, purely functional contract library for Haskell, which, when a contract is violated, also provides more useful information than the classical blaming of one contract partner. From now on lazy functional languages can profit from the assurances in the development of correct programs that contracts provide.

Categories and Subject Descriptors D.1.1 [Programming Techniques]: Applicative (Functional) Programming

General Terms Languages, Reliability

Keywords purely functional, lazy, library, Haskell

1. Introduction

Pre- and post-conditions have been important tools for developing correct programs since the early days of programming. A contract for a function comprises both a pre- and a post-condition. Figure 1 shows definitions in Haskell of two functions with contracts that operate on the type \texttt{Formula}, which represents propositional logic formulae. The two functions \texttt{clausalNF'} and \texttt{clause'} have rather non-descriptive types. The function \texttt{clausalNF'} transforms a propositional formula into clausal normal form. To work correctly, the function requires its input to be in conjunctive normal form and to be "right-bracketed", that is, for example \texttt{And (Atom 'a') (And (Atom 'b') (Atom 'c'))} is used instead of \texttt{And (And (Atom 'a') (Atom 'b')) (Atom 'c')}. The output is a list of list of literals, where a literal is an atom or a negated atom. This pre-condition and post-condition are expressed in the contract \texttt{conjNF & right >-> list (list lit)}, which is attached to \texttt{clausalNF'} using the function \texttt{assert} in the definition of the contracted function variant \texttt{clausalNF}. The function \texttt{clause} has a similar contract. For any contract \texttt{c} the function \texttt{assert c} is roughly the identity function, except that it also enforces the contract. The program states the contracts and monitors them at runtime.

Since the work of Findler and Felleisen \cite{12} on contracts for eager functional languages, contracts have become an important item in the toolbox of the Racket/Scheme programmer. Other functional languages, however, have not yet profited from the support of contracts for several reasons:

- Eager functional contracts were introduced as a small library of contract combinators. However, the implementation in Racket uses its powerful macro system to smoothly integrate contracts into the language\footnote{Recently added features concerning mutable data also required modifications of the language implementation\cite{10}.}. Thus contracts are very easy to use, for example, do not require user-supplied program location parameters. Implementors of other programming languages, however, do not have such a powerful macro system and are wary of making the implementation effort and of extending the language.

- In contrast to dynamically typed Racket, many functional programming languages have a static type system based on Hindley-Milner types with parametric polymorphism. Thus contract combinators need to be statically typed too and it is desirable to have type-directed contract combinators such as \texttt{list :: Contract a -> Contract [a]}. We also have to avoid classes in types. To see why, consider the following example usage of the Haskell object observation debugger HOOD \cite{16}:

\begin{verbatim}
data Formula =
  Imp Formula Formula | And Formula Formula | 
  Or Formula Formula | Not Formula | Atom Char

clausalNF =
  assert (conjNF & right >-> list (list lit))
  clausalNF'

clausalNF' :: Formula -> [[Formula]]
  clausalNF' (And f1 f2) = clause f1 : clausalNF' f2
  clausalNF' f = [clause f]

clause = assert (disj & right >-> list lit) clause'

clause' :: Formula -> [Formula]
  clause' (Or f1 f2) = f1 : clause' f2
  clause' lit = [lit]
\end{verbatim}

Figure 1. Contracts of functions for clausal normal form
Here observe "fun" behaves like an identity function but also records input and output of the \texttt{length} function for debugging purposes. However, the observation function does not have the type of the identity function: the type of \texttt{length} includes the class \texttt{Observable} and thus adding an observation may require substantial changes to type annotations in the whole program. To avoid this problem, we have to ensure that contract combinators have simple parametrically polymorphic types, without class contexts.

- Eager functional contracts are strict. Let \texttt{nat} be the contract that holds only for non-negative integers. With strict contracts we would get

\[
\text{assert} \ (\text{list nat}) \ [4,-1,2] = \text{error} "\ldots"
\]

Asserting an eager contract yields either the unchanged argument or an error/exception. In contrast, lazy functional languages demand lazy contracts. Asserting a lazy contract yields those parts of the argument data structure that meet the contract; only those parts that violate the contract are "cut off":

\[
\text{assert} \ (\text{list nat}) \ [4,-1,2] = [4,\text{error} "\ldots",2]
\]

Because lazy evaluation generally only evaluates parts of a data structure, a computation may succeed without any contract violation error, if it only demands those data structure parts that meet the contract. Such lazy contracts preserve the lazy semantics of the program and thus ensure that we can add contracts anywhere in a program without changing its semantics, provided contracts are not violated. For example, the following definition of the infinite list of fibonacci numbers requires lazy contracts:

\[
\text{fibs} :: [\text{Integer}]
\]

\[
\text{fibs} = \text{assert} \ (\text{list nat})
\]

\[
(0 : 1 : \text{zipWith (+)} \ \text{fibs} \ (\text{tail fibs}))
\]

In this paper we develop a library for contracts in Haskell that makes the following contributions:

- The contract combinators have simple parametrically polymorphic types, such that adding a contract does not change the type of a function (Section 2).
- The library provides lazy contract combinators. Adding contracts leaves the semantics of a program unchanged, unless a contract is violated (Sections 2 and 3).
- The library is written in pure, portable Haskell, without any use of side-effecting primitives such as \texttt{unsafePerformIO} that could change the semantics of the program (Section 3).
- All data-type-dependent code is simple and thus easy to write by hand if necessary (Section 3).
- The contract combinators have a nice algebra of properties. Contract assertions are partial identities and we claim that they are idempotent too. Thus contracts are projections, like eager contracts (Section 4).
- If a contract is violated, then the raised exception does not simply blame the server (contracted expression) or its client, but provides additional information about the specific value that caused the violation (Section 5).
- The library can use Template Haskell to derive all data-type-dependent code and include source code locations. Thus the programmer can formulate contracts for new algebraic data types without any additional work (Section 6).

The contract library for Haskell is available on Hackage\(^2\).

2. Simple Contract Combinators with a Problem
Using previous work on eager contracts [11] and typed contracts [19], we can easily design and implement most of a contract library for Haskell.

We implement a parametric type \texttt{Contract} \(a\) and a function

\[
\text{assert} :: \text{Contract} \ a \rightarrow (a \rightarrow \text{a})
\]

that turns a contract into a partial identity, that is, \texttt{assert} \(c \sqsubseteq\) \texttt{id}. Here \(\sqsubseteq\) is the standard information-theoretic partial order on values with least element \(\bot\). For simplicity we consider \(\bot\) to be an expression. It represents both non-termination and an exception raised by a violated contract.

Most of the contract library consists of combinators for building contracts of type \texttt{Contract} \(T\), for various types \(T\).

We start with a combinator that turns a predicate into a contract:

\[
\text{prop} :: \text{Flat} \ a \Rightarrow (a \rightarrow \text{Bool}) \rightarrow \text{Contract} \ a
\]

used for example as

\[
\text{nats} :: \text{Contract} \ \text{Integer}
\]

\[
\text{nats} = \text{prop} \ (\geq 0)
\]

to specify natural numbers as integers greater or equal zero.

We have to restrict \text{prop} by a new class \texttt{Flat} to be used for flat types only. A type is flat if for all values \(v_1\) and \(v_2\) the ordering \(v_1 \sqsubseteq v_2\) implies \(v_1 = \bot\). We cannot use \text{prop} for non-flat types such as lists, because the predicate could be arbitrarily strict and thus violate our aim of building lazy contracts [2]. For example

\[
\text{nats}' = \text{prop} \ (\geq (0)) :: \text{Contract} \ \text{Integer}
\]

would not be lazy and thus would be unusable for our infinite list of fibonacci numbers.

So the class \texttt{Flat} has only a few instances such as

\[
\text{instance} \ \text{Flat} \ \text{Integer}
\]

\[
\text{instance} \ \text{Flat} \ \text{Float}
\]

\[
\text{instance} \ \text{Flat} \ \text{Char}
\]

In our initial example in Figure 1 we already used two combinators for building contracts:

\[
(\text{k}) :: \text{Contract} \ a ightarrow \text{Contract} \ a ightarrow \text{Contract} \ a
\]

\[
(\triangleright\triangleright) :: \text{Contract} \ a ightarrow \text{Contract} \ b ightarrow \text{Contract} \ (a \rightarrow b)
\]

The conjunction combinator \((\&\&\&\&)\) builds a contract that is violated if one of the components is violated. The combinator \((\triangleright\triangleright)\) looks similar to the function type; it does not indicate logical implication. The function combinator combines a pre- and a post-condition to a contract for a function. For a function to be correct, whenever the pre-condition holds, the post-condition must hold too. However, if the pre-condition is violated, then the client (caller) of the function is wrong. A function contract is an agreement between both a function and its client. So neither pre- nor post-condition should be violated. In summary, the function contract combinator is rather like the conjunction combinator, except that values of two possibly different types are monitored.

A contract that is always met is useful as component of a bigger contract to express that some values are irrelevant. The opposite contract that is never met can still be occasionally useful in a lazy functional language. We can use

\[
\text{true} :: \text{Contract} \ a
\]

\[
\text{false} :: \text{Contract} \ a
\]

for example in

\(^2\)\url{http://hackage.haskell.org}
type Contract a = a -> a

assert c = c

class Flat a where
  prop :: (a -> Bool) -> Contract a
  prop p = \x -> if p x then x else error "..."

pNil [] = []
pNil (...) = error "...

pCons c cs [] = error "..."
pCons c cs (x:xs) = Just (assert c x : assertion cs xs)

true = Just id
false = Just Nothing

c1 & c2 = c2 . c1
pre >> post = \f -> post . f . pre

Figure 2. A lazy contract implementation for most combinators

const = assert (true >> false >> true) const'
const' x y = x

to express that the second argument of the function is never demanded. Because that argument is never demanded, its contract will never be used and thus will never be violated.

Finally we need combinators to build contracts for algebraic data types, which generally are not flat. Here we introduce for each data constructor a combinator that is used like a data constructor in pattern matching.

pNil :: Contract [a]
pNil [] = Just []
pNil (...) = Nothing

pCons :: Contract a -> Contract [a] -> Contract [a]
pCons c cs [] = Nothing
pCons c cs (x:xs) = Just (assert c x : assertion cs xs)

true :: Contract a
true = Just Nothing

false :: Contract a
false = Just Nothing

(|>) :: Contract a -> Contract a -> Contract a
(|>) c1 c2 = \x -> c1 x 'mplus' c2 x

(<<>) :: Contract a -> Contract a -> Contract a
(<<>) c1 c2 = \x -> c1 x '>>=' c2

pre >> post = \f -> Just (f 'seq' (assert post . f . assert pre))

Figure 3. Implementation of typed lazy contract combinators

We can combine two functions of type a -> a only by composition and we have done so already for the contract combinator (%). For disjunction we would need to apply both functions separately and then somehow combine the two results: if one is an exception, then we should return the value of the other one. We cannot test for exceptions in a purely functional language.

3. Implementing Lazy Contract Combinators

A simple modification of our contract type definition solves our problem:

type Contract a = a -> Maybe a

assert :: Contract a -> (a -> a)
assert c x = case c x of
  Just y -> y
  Nothing -> error "Contract violated."

class Flat a where
  prop :: (a -> Bool) -> Contract a
  prop p x = if p x then Just x else Nothing

pNil :: Contract [a]
pNil [] = Just []
pNil (...) = Nothing

pCons :: Contract a -> Contract [a] -> Contract [a]
pCons c cs [] = Nothing
pCons c cs (x:xs) = Just (assert c x : assertion cs xs)

true :: Contract a
true = Just True

false :: Contract a
false = Just False

(|>) :: Contract a -> Contract a -> Contract a
(|>) c1 c2 = \x -> c1 x 'mplus' c2 x

(<<>) :: Contract a -> Contract a -> Contract a
(<<>) c1 c2 = \x -> c1 x '>>=' c2

pre >> post = \f -> Just (f 'seq' (assert post . f . assert pre))

3 There is an impure solution [6] that, however, still cannot handle non-termination in one argument.
is only safe, because the function take is strict in its integer parameter, which determines how
many list elements shall be returned.

prop $p_1 \mid> p_2 = prop (\langle x \rightarrow p_1, x \mid \mid p_2, x \rangle)$
prop $p_1 \& p_2 = prop (\langle x \rightarrow p_1, x \& p_2, x \rangle)$

$c_1 \& (c_2 \& c_3) = (c_1 \& c_2) \& c_3$
true $\& c = c$
$c \& true = c$
false $\& c = false$

c_1 \mid> (c_2 \mid> c_3) = (c_1 \mid> c_2) \mid> c_3$
false $\mid> c = c$
c $\mid> false = c$
true $\mid> c = true$
c $\mid> c = c$
c_1 \mid> (c_2 \mid> c_1) = c_1 \mid> c_2$

c_1 \mid> (c_1 \& c_2) = c_1$
true $\mid> true = true$
c_1 \mid> false = c_2 \mid> false$

$(c_1 \mid> c_2) \& (c_3 \mid> c_4) = (c_3 \& c_1) \mid> (c_2 \& c_4)$
$(c_1 \mid> c_2) \mid> (c_3 \mid> c_4) = c_1 \mid> c_2$

Figure 4. Contract properties

$(c_1 \& (c_1 \mid> c_2) = c_1$
$c_1 \& (c_2 \& c_1) = c_1 \& c_2$
c $\& c = c$

Figure 5. Claimed contract properties

4. Properties of Contracts

Our contract type Contract $a$ is a combination of the function type with the \texttt{Maybe} monad. Thus we have a rich set of known properties to work with for establishing an algebra of contracts. Contract itself is \texttt{not} a monad.

4.1 An Algebra of Contracts

Figure 4 lists many simple properties enjoyed by contracts. All of these can be proved by simple equational reasoning, using the monad laws of \texttt{Maybe} $a$.

All the properties of conjunction and disjunction of contracts also hold for conjunction and disjunction of Booleans in a lazy language. Recall that some standard properties of Boolean algebra do not hold for the Boolean type in non-strict languages. For example, $(kk)$ and $\mid\mid$ are not commutative and the standard distribution laws do not hold. These properties do not hold for contracts either, with similar counterexamples. So the non-strict algebra of $(kk)$ and $\mid\mid$ is a good guideline for developing the lazy contract algebra of $k$ and $\mid\mid$.

The "distribution" law for conjunction and function contract may at first surprise. It holds because the function contract combinator is not some kind of implication but more a kind of conjunction. A function contract holds only if both the input and the
output of a function meet the respective subcontracts. From this "distribution" law of conjunction and function plus idempotence of conjunction further laws follow:

\[(c_1 \rightarrow c) \land (c_2 \rightarrow c) = (c_1 \land c_2) \rightarrow c\]

\[(c \rightarrow c_2) \land (c \rightarrow c_1) = c \rightarrow (c_2 \land c_1)\]

Figure 5 lists further properties of contracts that we have not proved but claim also hold. They require stronger proof methods than equational reasoning, but are linked to the idempotence of contracts discussed in the subsequent subsection. The last property in the list, idempotence of conjunction, is a corollary of the preceding property, taking \(c_2 = true\).

4.2 Contracts are Projections

Eager contracts are projections [1, 11], that is, they are idempotent and partial identities.

Lemma 4.1 (A contract is a partial identity).

For any contract \(c\)

\[assert c \subseteq id\]

This can be proved using induction on the contract combinators. Idempotence is more difficult to establish. It would follow from idempotence of conjunction, \(c \land c = c\). In practice, both properties probably need to be established in a single inductive proof. Intuitively idempotence holds because if a contract returns \(true\), then \(true\) is the same as would be returned by the eager contract \(true \rightarrow a\), and pattern contracts only test for the top constructor before returning \(true\) or Nothing.

Claim 4.2 (A contract is idempotent).

For any contract \(c\)

\[assert c . assert c = assert c\]

4.3 Distinct Contract Exceptions

We identified non-termination and any contract exception as the single value \(⊥\). However, we might distinguish them, following [22], such that exceptions are values above \(⊥\) in the information order, but still values of any type. We would change our partial order to consider exceptions as least elements, because a contract replaces some parts of values by exceptions. However, with that choice our contracts are neither partial identities, nor idempotent (the properties of Figure 4 are unaffected). The reason is that contracts such as

\[⊥ :: Contract a\]

\[\text{prop } ⊥ :: \text{Flat a } -> \text{Contract a}\]

\[\text{prop } (\lambda x -> \text{if odd x the True else } ⊥ ) :: \text{ContractInt}\]

exist. They would still introduce \(⊥\) instead of exceptions.

For related reasons other works [1, 8] restricted the definitions of contracts such that a contract can never introduce \(⊥\) itself. However, the desirable freedom to use the whole language to define contracts and the fact that we are just defining a library makes this an impractical choice.

5. Informative Contract Violation

A contract is concluded between two partners, a server and a client. If a contract is violated, one of the two partners is to blame for it. A major contribution of Findler and Felleisen’s functional contracts [12] is its system for choosing whom to blame. In a higher-order language function arguments can themselves be functions. If such a functional argument is used within the function such that the precondition of the functional argument is violated, then the function itself has to be blamed for contract violation, not the caller that passed the functional argument.

```
type Contract a = a -> Bool -> Either Bool a

assert :: Contract a -> (a -> a)
assert = monitor True

monitor :: Bool -> Contract a -> (a -> a)
monitor b c x =
case c x b of
  Right x -> x
  Left b -> error("Contract violated. Blame " ++ if b then "server." else "client.")

( >>= ) :: Contract a -> Contract b -> Contract (a->b)
pre >>= post = \f b -> Right (f 'seq'
  (monitor b post.f..monitor (not b) pre))

true :: Contract a
true = \x b -> Right x

false :: Contract a
false = \x b -> Left (not b)
```

Figure 6. Implementing blaming

5.1 Blaming

The blaming system for higher-order functional languages applies to both eager and lazy languages equally, and thus we can easily add it to our lazy contract library. For eager languages several equivalent implementations for handling blame are known [11, 12, 19]. Here we simply extend a contract by a Boolean state that indicates whether the server or the client of the contract are to blame in case of violation. The Maybe monad is replaced by Either Bool a so that blame information is available when a sub-contract is violated. Figure 6 shows the most interesting extended definitions. Contract monitoring starts by potentially blaming the server, that is, the expression for which the contract is asserted. The function contract combinator \(>>->\) negates the Boolean blame indicator for monitoring the co-variant argument, but passes it unchanged for monitoring the co-variant result.

Now there are two different possible implementations of the contract \(false\) that can never be met: The contract either always blames the party indicated by the given Boolean argument, or it always blames the opposite party by negating the Boolean value. So let us look back at our example of Section 2:

\[\text{const } = \text{assert (true } >>- \text{false } >>- \text{true) const}\]

Any client of \(\text{const}\) will provide some second argument, but if that second argument is actually demanded, then clearly \(\text{const}\) is wrongly defined and has to be blamed. In this example \(false\) is in a co-variant position of the whole contract and hence to blame the server, \(false\) has to negate its Boolean parameter. So on its own, \(false\) always blames its client, never its server. We do not provide the server-blaming variant in the library, because it does not seem to be of any practical use.

5.2 Witness Tracing

Blaming alone, however, is rather unsatisfactory. It just points the finger at one partner without providing any evidence that would explain in which way a complex contract was violated. Blaming hardly provides a good starting point for debugging. Furthermore, blaming can be misleading. Often when a contract is violated neither server nor client are wrong, but the contract itself! Specifying
type Contract a =
(String->String) -> a -> Either String a

assert :: Contract a -> (a -> a)
assert = monitor id

monitor :: (String->String) -> Contract a -> (a->a)
monitor wc c x =
    case x of
        Right v -> v
        Left w   -> error ("Contract violated. Witness:" ++ wc ("{" ++ w ++ "}") )

(>>>) :: Contract a -> Contract b -> Contract (a->b)
pre >> post = \wc f -> Right \f 'seq'
    (monitor (wc . \w->("{->"++w++"}"))) post . f .
    monitor (wc . \w->("{++w++->}_{}")) pre)

pNil :: Contract [a]
pNil = \wc x -> case x of
    []   -> Right x
    _:_  -> Left "::"           

pCons :: Contract a -> Contract [a] -> Contract [a]
pCons cx cxs = \wc x -> case x of
    (y:ys) -> Right
        (monitor (wc . \w->("{++w++->}_{}")) cx y :)
    []     -> Left "[]"

 pCons cx cxs = \wc x -> case x of
    (y:ys) -> Right
        (monitor (wc . \w->("{++w++->}_{}")) cx y :)
    []     -> Left "[]"

Figure 7. Implementing witness tracing

the right contract is challenging and contract monitoring just checks
whether specification and implementation agree.

Hence our lazy contracts report, when they are violated, the top
data constructor, or whole flat value, that causes the contract viola-
tion, plus all data constructors in the path above it. For example

> :> clausalNF form
[[Atom 'a'],[Atom 'b'],Not
*** Exception: Contract at ContractTest.hs:101:3
violated by
((And _ (Or _ (Not _))) ) ->

The client is to blame.

6. Deriving Contract Combinators

For every data constructor Con that we want to pattern match in a
contract we have to define a pattern contract pCon. These defini-
tions are simple, even with handling of location, blame and witness
information, but they are still tedious. Hence our contract library al-
 lows their automatic derivation using Template Haskell [23]. Tem-
plate Haskell is a meta-programming extension of Haskell that
the Glasgow Haskell compiler, the only Haskell system used for
 professional Haskell program development, provides. Template
Haskell allows us to define in the contract library functions that
will generate Haskell code at compile time, type check that code
and compile it.

The user no longer needs to define these pattern contracts at
all, but can basically derive them on demand where needed, that is,
directly write

conjNF = $(p 'And) conjNF conjNF | > disj
     = $(p 'Or) disj disj | > lit
lit    = $(p 'Not) atom | > atom
atom   = $(p 'Atom) true

Here p is a Template Haskell function that receives the name of a
data constructor as argument. The $ and the single quote in front of
the data constructor are syntax required by Template Haskell. The
definition of a pattern contract is short so that repeated derivation
is not a problem.

Alternatively, the programmer can also write the declaration

$ (deriveContracts 'Formula)

to derive all pattern contract definitions for the type Formula.

Finally, assert is also a Template Haskell function:

clausalNF =
    $assert (conjNF & right >> list (listlit))
    clausalNF'

Template Haskell allows the definition of assert to determine its
own location in the file and then generate code for calling the real
assertion function with that location as parameter.

7. Further Contract Features

Initial experience of using contracts raises new questions and de-
mand for additional contract combinators.

7.1 Negation

We have conjunction, &, and disjunction, |>, of contracts. However,
we cannot have negation

neg :: Contract a -> Contract a
for contracts. General negation would violate basic semantic prop-
eries of contracts [2].

Nonetheless, in practice we often want to express that the top
data constructor of a monitored value is not a specific given data
constructor. Hence we introduce additional combinators such as the
following for every data type.

pNotImp :: Contract Formula
pNotAnd :: Contract Formula
pNotOr :: Contract Formula
pNotNot :: Contract Formula
pNotAtom :: Contract Formula

These negated pattern contracts provide nothing new. In fact

\[ p\text{NotImp} = p\text{And} \text{true} \text{true} \rightarrow p\text{Or} \text{true} \text{true} \rightarrow p\text{Not} \text{true} \rightarrow p\text{Atom} \text{true} \]

However, for types with many data constructors these combina-
tors are certainly substantial abbreviations and they are needed fre-
quently. Additionally, our implementation can perform an efficient
single pattern match instead of many repeated ones.

We use these negated pattern contracts in the definition of con-
tracts for our initial propositional formulae example. They sub-
tially simplify our definition of “right-bracketedness”.

\[ \text{conjNF, disj, lit, atom, right, rightConjNF :: Contract Formula} \]

\[ \text{conjNF} = p\text{And conjNF conjNF} \rightarrow \text{disj} \]
\[ \text{disj} = p\text{Or disj disj} \rightarrow \text{lit} \]
\[ \text{lit} = p\text{Not atom} \rightarrow \text{atom} \]
\[ \text{atom} = p\text{Atom true} \]

\[ \text{right} = p\text{Imp (right & pNotImp) right} \rightarrow \]
\[ p\text{And (right & pNotAnd) right} \rightarrow \]
\[ p\text{Or (right & pNotOr) right} \rightarrow \]
\[ p\text{Not right} \rightarrow p\text{Atom true} \]

\[ \text{rightConjNF} = \text{conjNF} \& \text{right} \]

Even for data types with few constructors they can express an idea
more clearly. So

\[ \text{head'} = \text{assert (pNotNil} \rightarrow \text{true}) \text{head} \]
is more direct than

\[ \text{head'} = \text{assert (pCons true true} \rightarrow \text{true}) \text{head} \]
to express that the function only works on non-empty lists.

7.2 Contracts for the IO monad

We have contract combinators for flat types, algebraic data types
and the function type constructor. However, a real programming
language has more types, especially abstract data types. The most
notorious in Haskell is the IO monad that is required for any input
or output actions.

For example, we may want to write a contract for an IO action
that gets a natural number from standard input:

\[ \text{getNat :: IO Integer} \]
\[ \text{getNat} = \text{assert} (\text{io nat}) \text{getNat'} \]

The choice of contract combinator is natural, following our
general approach of type-directed contract combinators. How do
we define the IO contract combinator?

\[ \text{io :: Contract a} \rightarrow \text{Contract (IO a)} \]
\[ \text{io c} = (\text{io} \rightarrow \text{Just (io >>= return . assert c}) \]

Our definition simply follows the scheme we are already using
for the function contract combinator \(\rightarrow\). After all, the function
type is “just” an abstract data type as well\(^3\). With this definition our

\(^3\)The forced evaluation with \text{seq} by the function contract combinator is
required because of the peculiar semantics of functions in Haskell. It is
not needed for other types. For example, if \text{io} = \bot, then also \text{io >>=
return . assert c} = \bot.

IO contract combinator also has the same properties as the function
contract combinator:

\[ \text{io c} = \text{io c_1} \& \text{io c_2} = \text{io (c_1} \& \text{c_2)} \]
\[ \text{io c_1} \rightarrow \text{io c_2} = \text{io c_1} \]
\[ \text{io true} = \text{true} \]

Our definition of the contract combinator for the abstract data
type \text{IO} a raises the question whether we should do the same for
other data types. For example, we have

\[ \text{list :: Contract a} \rightarrow \text{Contract [a]} \]
\[ \text{list c} = p\text{Nil} \rightarrow p\text{Cons c} \text{ (list c)} \]

Alternatively we could follow our definition of \text{io}:

\[ \text{list'} = \text{xs} \rightarrow \text{Just (xs >>= return . assert c)} \]

which is the same as

\[ \text{list'} = \text{xs} \rightarrow \text{Just (map (assert c) xs)} \]

It turns out that the two definitions are equivalent\(^4\), thus confirm-
ing our original definition. Hence we prefer to define a contract
combinator for a non-abstract algebraic data type such as \text{list} in
terms of the only primitive contract combinators, the pattern con-
tract combinators, such as \text{pNil} and \text{pCons}. For an abstract data
type we define a contract combinator using the scheme above with
the respective map function for the type.

7.3 Strict data types

Haskell allows the definition of strict data types. The strictness flag
\(!\) in a data type definition states that the data constructor is strict in
that argument. For example,

\[ \text{data SListBool} = \text{SNil} | \text{SCons !Bool} !\text{SListBool} \]
defines the type of finite Boolean lists that are either \(\bot\) or fully
defined.

Luckily we do not need to adapt our definition of contract
combinators. As usual we have

\[ \text{pSNil} :: \text{Contract SListBool} \]
\[ \text{pSNil SNil} = \text{Just SNil} \]
\[ \text{pSNil} (\text{SCons b bs}) = \text{Nothing} \]

\[ \text{pSCons :: Contract Bool} \rightarrow \text{Contract SListBool} \rightarrow \]
\[ \text{Contract SListBool} \]
\[ \text{pSCons c cs SNil} = \text{Nothing} \]
\[ \text{pSCons c cs (SCons b bs)} = \]
\[ \text{Just (SCons (assert c b) (assert cs bs))} \]

A contract traverses the strict list and builds a new strict list.
Thus demanding the top data constructor of a contracted strict
list automatically forces checking the whole list. The result will
be either a contract violation (\(\bot\)) or the whole list. So on strict
data types lazy contracts behave like eager contracts. It would be
possible to define more expressive contract combinators for data
types with strictness flags, that, for example, ensure that a list
is ordered; but because strictness flags are rarely used in Haskell
programs, such an extension does not seem worthwhile.

In the definition of \text{SListBool} all constructor arguments are
strict and the only other types used are flat types. In such a case
the data type is actually a flat type. We can declare it an instance of
the class \text{Flat} and use expressive \text{prop} contracts.

\(^4\)They are not equal, because \text{list \bot} = \bot, but \text{list'} \bot = \text{Just \bot}.
However, in the context of a contract with \text{assert} they always yield the
same result.
8. Related Work

This paper builds firmly on three sets of previous work: Findler and Felleisen’s work on eager contracts for higher-order functions [12], Hinze, Jeuring and Löh’s work on typed contracts for functional programming, and our own previous work on lazy functional contracts.

Eager Higher-Order Contracts  Findler and Felleisen’s paper on contracts for higher-order functions [12] made contracts popular for eager functional programming languages. All interesting properties of functional values, which are passed around by higher-order functions, are undecidable; it is impossible to monitor a function contract for all argument-result-pairs. However, Findler and Felleisen realised that it is sufficient to monitor both pre- and post-condition of a functional value only when this function is applied. The resulting contract system is sound.

The second major contribution of that paper is a system for correctly attributing blame in case of contract violation and its implementation. We easily added blaming to our lazy contract library. However, additionally our contracts report a witness, a partial value, that caused a contract violation.

Findler and Felleisen’s contract system also provides dependent function contracts, where the contract for the function codomain can use the actual argument value. Such dependent contracts are more expressive but easily change a non-strict function into a strict function; hence our lazy contract combinators do not provide them.

Subsequent work [11, 13] stresses that contracts are projections and thus they can be implemented in a simple, modular and efficient way. Our first implementation of Figure 2 copies that work and our full implementation with disjunction is an extended variant.

The papers do not discuss algebraic data types, because in strict languages these domains are flat and hence contracts for algebraic data types are predicates like for other flat types. Consequently disjunction is not considered either. Because of the universality of predicate contracts in strict and dynamically typed functional languages, type-directed contracts such as \texttt{list} are also of little interest.

Although disjunction is not discussed in the papers, the contract system of Racket does provide a disjunctive contract combinator \cite{racket-disjunction}. To support disjunction, a Racket contract for \texttt{type a} contains both a function of type \texttt{a -> a} and a function of type \texttt{a -> Bool}. Together they are used similarly to our type \texttt{a -> Maybe a}. In particular, the disjunctive combinator applies the \texttt{a -> Bool} functions of all its direct sub-contracts, checks that at most one of the results is \texttt{True} (otherwise it fails) and then applies the corresponding sub-contract further \cite{racket}. So disjunction behaves similarly to \texttt{||} but is not sequential.

Blume et al. proposed and studied several semantic models of eager higher-order contracts \cite{blume95,blume97,blume09}. To prove soundness, the definition of contracts is first restricted, to avoid e.g. having a contract \texttt{a -> a}. Later, recursive contracts are added to regain expressivity. In a discussion of the most permissive contract, \texttt{true}, Findler and Blume point out that the contract \texttt{true}, to be the most permissive contract, should always report contract violation and blame the client. However, they also note that such a contract would be useless in practice. In contrast, our \texttt{true}, which cannot be violated, is very useful to leave parts of a contract unconstrained. Similarly, the least permissive contract should always blame the server, but we demonstrated in Section 5.1 that our definition of \texttt{false}, which always blames the client, is more useful. As a consequence in our library \texttt{false} = \texttt{prop (const False)} does not hold for flat types whereas \texttt{true} = \texttt{prop (const True)} does.

Typed Contracts for Functional Programming  Hinze, Jeuring and Löh [19] transferred contracts for higher-order functions to the statically typed language Haskell. Hence they proposed contract combinators with parametrically polymorphic types; we have adopted all of them except for dependent contracts. Typed contracts also emphasise type-directed contract combinators such as \texttt{list}. However, the work disregards the lazy semantics of Haskell, defining contracts with a seemingly random mixture of eager and lazy monitoring. Predicate contracts can be applied to expressions of all types, not just flat types, thus breaking laziness. However, these predicate contracts are required for expressing many interesting properties, because a type-directed contract combinator such as \texttt{list} can only express a uniform property over all list elements: our pattern contracts and disjunction are missing.

Contracts are not projections, because generally they are not idempotent. Idempotence is lost because of the eagerness of predicate contracts. Hinze et al. make the point that if contract conjunction \texttt{&} was commutative, then idempotence would be a simple consequence. However, our lazy contracts demonstrate that commutativity of conjunction is not necessary for idempotence; we can have the latter without the former.

Hinze et al. also provide an interesting technique for providing more informative error messages than standard blaming. Their library provides several source locations as explanation of a single violated contract. However, these sets of source locations are still hard to understand for a programmer and the system requires a source code transformation to insert source locations into the program. Otherwise the programmer would have to do this substantial work.

Lazy Contracts for Functional Languages  Lazy contracts were first discussed and several implementations presented in 2004 \cite{casner04}. That paper makes the point that while eager contracts must be \texttt{True}, lazy contracts must not be \texttt{False}. This means that unvaluated parts of a data structure can never violate a lazy contract. The paper uses predicates on values of all types and hence, despite some technical tricks using concurrency, the contracts are lazy but neither idempotent nor prompt. The paper itself gives examples of where contract violations are noticed too late. This problem was later rectified \cite{casner05,findler06}. Both these papers implement lazy assertions as libraries that require only the commonly provided non-pure function \texttt{unsafePerformIO}, which performs side effects within a purely functional context. The first lazy and idempotent implementation \cite{findler06} uses patterns contracts similar to those in this paper to express contracts over algebraic data types. However, a non-deterministic implementation of disjunction leads to semantic problems. Later \cite{casner05} provided a more user-friendly language for expressing contracts and improved the internal structure of the implementation, but the implementation principles were identical and hence the non-deterministic disjunction remained.

A semantic investigation \cite{branchey08} developed contracts that are pure and implementable within the functional language. However, for every algebraic data type its contracts requires a different implementation type. Thus disjunction is not a parametrically polymorphic combinator but requires a class context. Furthermore, the implementations of some combinators are large and complex. Disjunction is more powerful than in the lazy contracts described in the present paper, for example

\begin{verbatim}
assert (pCons nat pNil |> pCons true pNil) |-3
=assert (pCons (nat |> true) (pNil |> pNil)) [-3] =[-3]
\end{verbatim}

but this additional expressibility does not seem to be needed in practice.

Comparing Contracts  Degen, Thiemann and Wehr \cite{degen11,thiemann11} classify existing contract systems for Haskell as eager (straight translation of [12]), semi-eager \cite{wadler04} and lazy \cite{casner04,findler06}. They check whether the systems meet their desirable properties of meaning preservation
and completeness. Each contract system meets at most one of these properties. The authors show that it is impossible to meet both properties. Our lazy contracts are meaning preserving but not complete. The notion of completeness seems to be biased towards a strict semantics, contradicting the principle that unevaluated parts can never violate a lazy contract. Our lazy contracts have limited expressibility, but they have a clear semantics.

**Generic Programming** We use Template Haskell to derive pattern contracts and to enable the assertion function to determine its own location in the source code [23]. The derivation of pattern contracts is an instance of generic programming. Many generic programming systems have been proposed and even been implemented for Haskell [17, 18, 20, 21]. All of these have two disadvantages that make them unsuitable for being used for our pattern contracts: First, they introduce one or more classes that will then appear in the type of every derived pattern contract. Thus pattern contracts will not be parametrically polymorphic. Second, they consider functions as second class values. That means that either they can only generically define code for types that do not involve function types at all, or they can recognise a functional value within an algebraic data type, but cannot do anything with it, that is, apply any transformation to it.

Template Haskell provides few static guarantees and thus requires us to ascertain that our contract library will derive typeable and correct code for any data constructor. However, Template Haskell provides all the functionalities needed in the contract library.

9. Conclusions and Future Work

This paper describes the design of a practical contract library for lazy typed functional languages and its implementation for Haskell. The library meets many essential criteria, such as combinators with simple parametric types, a lazy semantics, a rich algebra of properties, informative exceptions in case of contract violation and automatic code generation to make it easy to use.

Interestingly the resulting contract system reminds strongly of a subtyping system, especially with a definition of sub-contracts/types for algebraic data types that looks very similar to the actual type definition of algebraic data types. Defining subtypes of algebraic data types is also where we see the main application area of the contract system. Many programs require several variants of some big algebraic data types. In practice programmers then simply ignore the subtyping and define a single algebraic data type that encompasses all variants, because they want to reuse functions that work on several subtypes and have the flexibility to exchange some code without having to change between numerous similar but separate data types. The classical example is a compiler: it consists of a long sequence of passes, each of which works with a slightly differently structured abstract syntax tree. In practice, subtle differences are ignored and only a few different abstract syntax tree structures are used in one compiler. Lazy contracts provide a new solution.

Our next step is to develop the algebra of contract combinators further and thus also prove our claim that these contracts are idempotent. The main current shortcoming and that biggest challenge for future development of the contract library is its lack of a dependent function contract combinator that allows using the function argument in the post-condition. We can define

\[( pre >> >>= post ) >>= Contract a \rightarrow (a \rightarrow Contract b) \rightarrow Contract a \rightarrow b ]

\pre >>= post = \nf \rightarrow \Just (f \ 'seq' \ \hx \rightarrow \Let y = assert pre x \in assert (post y) (f y))\]

and use it for example in

```haskell
contractTake :: Int -> [a] -> [a]
contractTake =
  assert (prop (>=0) >>=
    \n -> lengthAtLeast n >>= listOfLength n)
  take
```

It is easy to extend this pithy implementation to use indy monitoring [9], which may blame the contract itself, not just the server or the client. However, >> >>= is not a lazy contract combinator; the post-condition may force evaluation of too much of the function argument and thus the contract may change the semantics of the program. A definition of a lazy dependent function combinator is still an open problem. Meanwhile the existing contract library can be used in practice.

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**References**


A. Proofs

The following are proofs for Section 4. All proofs use the implementation of contracts given in Figure 3.

A.1 General properties of contracts

Lemma A.1. Let \( c \) be any contract. If \( c \downarrow = \text{Just} \ v' \) implies that \( v' \sqsubseteq v \), then \( \text{assert} \ c \downarrow \sqsubseteq v \).

Proof. Case analysis:

\[
c \downarrow = \bot: \text{assert} \ c \downarrow = \bot \sqsubseteq v.
\]

\[
c \downarrow = \text{Just} \ v': \text{assert} \ c \downarrow = v' \sqsubseteq v.
\]

\[
c \downarrow = \text{Nothing}: \text{assert} \ c \downarrow = \text{error} "..." = \bot \sqsubseteq v.
\]

Lemma A.2. Let \( c \) be any contract. If \( c \downarrow = \text{Just} \ v' \), then \( v' \sqsubseteq v \).

Proof. Induction on the contract \( c \), using its definition.

\[
\bot = \text{Just} \ v'; \text{Impossible}.
\]

\[
\text{true} = \text{Just} \ v': \text{Then} \ v' = v \sqsubseteq v.
\]

\[
\text{false} = \text{Just} \ v'; \text{Impossible}.
\]

\[
\text{prop} \ p = \text{Just} \ v': \text{Then} \ v' = v \sqsubseteq v.
\]

\[
\text{\( c_1 \land c_2 \downarrow = \text{Just} \ v' \):}
\]

\[
\text{Then} \ c_1 \downarrow = \text{Just} \ v'' \text{ and } c_2 \downarrow = \text{Just} \ v'''. \text{ With ind.}
\]

\[
\text{hypothesis follows} v' \sqsubseteq v'' \sqsubseteq v.
\]

\[
\text{\( c_1 \rightarrow c_2 \downarrow = \text{Just} \ v' \):}
\]

\[
\text{Then either } c_1 \downarrow = \text{Just} \ v' \text{ or } c_1 \downarrow = \text{Nothing} \text{ and } c_2 \downarrow = \text{Just} \ v'. \
\]

\[
\text{With ind. hypothesis follows} v' \sqsubseteq v \text{ in either case}.
\]

\[
\text{\( c_1 \Rightarrow c_2 \downarrow = \text{Just} \ v' \):}
\]

\[
\text{If} \ v' = \bot, \text{then} \ v' = \bot \text{ and thus} \ v' \sqsubseteq v.
\]

\[
\text{Otherwise} \ v' = \text{assert} \ c_2 \cdot v \cdot \text{assert} \ c_1. \text{ Let} \ v'' \text{ be any possible argument value of the function} \ v'. \text{ If the first assertion} \text{assert} \ c_1 \downarrow = \text{Just} \ ... \text{ or the second assertion} \text{assert} \ c_2 \downarrow = \text{Just} \ ...
\]

\[
\text{then} \ v' = v'' = \bot. \text{ Otherwise by induction hypothesis and Lemma A.1 we have} \text{assert} \ c_1 \downarrow = \text{Just} \ v'' \text{ and likewise also} \text{assert} \ c_2 \downarrow = \text{Just} \ v'''. \text{ With continuity of all functions follows} v \text{ (assert} \ c_1 \downarrow = v'') \sqsubseteq v \text{ and likewise also}\text{assert} \ c_2 \downarrow = \text{Just} \ (v \text{ (assert} \ c_1 \downarrow = v'')) \sqsubseteq v \text{ and} \text{assert} \ c_2 \downarrow = \text{Just} \ v'''. \text{ Altogether} \text{assert} \ c_2 \downarrow = \text{Just} \ (v \text{ (assert} \ c_1 \downarrow = v'') \sqsubseteq v \text{ and likewise also} \text{assert} \ c_2 \downarrow = \text{Just} \ v'''.
\]

Hence in all cases \( v' \sqsubseteq v \text{ and} \text{thus} \ v' \sqsubseteq v \).

\[
\text{\( \text{pCon} \ c_1 \ldots c_n \downarrow = \text{Just} \ v' \text{ where} \ p\text{Con} \text{ is the contract combinator for data constructor} Con.
\]

\[
\text{From the definition follows} v = \text{Con} \ v_1 \ldots v_n \text{ and} v' = \text{Con} \ (\text{assert} \ c_1 \downarrow v_1) \ldots (\text{assert} \ c_n \downarrow v_n). \text{ By ind.}
\]

\[
\text{hypothesis and Lemma A.1 we get that} \text{assert} \ c_1 \downarrow = v_1, \ldots, \text{assert} \ c_n \downarrow = v_n. \text{ So} v' \sqsubseteq v.
\]

Corollary A.3 (Contracts are partial identities).

\[
\text{assert} \ c \downarrow = v.
\]


Corollary A.4. \( \text{assert} \ c \) is strict for any contract \( c \).

Proof. From Corollary A.3 follows that \( \text{assert} \ c \downarrow = \bot. \text{ Hence} \text{assert} \ c \downarrow = \bot.
\]

Lemma A.5 (Asserting \&).

\[
\text{assert} \ (c_1 \land c_2) = \text{assert} \ c_2 \cdot \text{assert} \ c_1
\]

Proof.

\[
\text{assert} \ (c_1 \land c_2)
\]

\[
= \lambda x \rightarrow \text{case} \ c_1 \downarrow x >> c_2 \downarrow of
\]

\[
\text{Just} \ z \rightarrow z
\]

\[
\text{Nothing} \rightarrow \text{error} "..."
\]

\[
= \lambda x \rightarrow \text{case} \ (\text{case} c_1 \downarrow x \text{ of}
\]

\[
\text{Just} \ y \rightarrow c_2 \downarrow y
\]

\[
\text{Nothing} \rightarrow \text{Nothing}
\]

\[
\text{of} \text{Just} \ z \rightarrow z
\]

\[
\text{Nothing} \rightarrow \text{error} "..."
\]

\[
= \lambda x \rightarrow \text{case} \ c_1 \downarrow x \text{ of}
\]

\[
\text{Just} \ y \rightarrow \text{case} c_2 \downarrow y \text{ of}
\]

\[
\text{Just} \ z \rightarrow z
\]

\[
\text{Nothing} \rightarrow \text{error} "..."
\]

A.2 Properties of predicate contracts

Lemma A.7 (Disjunction).

\[
\text{prop} \ p_1 \ || \text{prop} \ p_2 = \text{prop} \ (\lambda x \rightarrow p_1 \downarrow x || p_2)
\]

Proof.

\[
\text{\( (c_1 \downarrow = \text{Just} \ v', c_2 \downarrow = \text{Just} \ v'', v' \sqsubseteq v, v'' \sqsubseteq v \) implies} \text{assert} \ c_1 \downarrow \sqsubseteq v, \text{assert} \ c_2 \downarrow \sqsubseteq v
\]

\[
\text{Thus} \text{assert} \ (c_1 \downarrow = \text{Just} \ v', c_2 \downarrow = \text{Just} \ v'', v' \sqsubseteq v, v'' \sqsubseteq v \) \rightarrow \text{assert} \ c_1 \downarrow \sqsubseteq v, \text{assert} \ c_2 \downarrow \sqsubseteq v
\]

\[
\text{Thus} \text{assert} \ (c_1 \downarrow = \text{Just} \ v', c_2 \downarrow = \text{Just} \ v'', v' \sqsubseteq v, v'' \sqsubseteq v \) \rightarrow \text{assert} \ c_1 \downarrow \sqsubseteq v, \text{assert} \ c_2 \downarrow \sqsubseteq v
\]

\[
\text{Thus} \text{assert} \ (c_1 \downarrow = \text{Just} \ v', c_2 \downarrow = \text{Just} \ v'', v' \sqsubseteq v, v'' \sqsubseteq v \) \rightarrow \text{assert} \ c_1 \downarrow \sqsubseteq v, \text{assert} \ c_2 \downarrow \sqsubseteq v
\]

\[
\text{Thus} \text{assert} \ (c_1 \downarrow = \text{Just} \ v', c_2 \downarrow = \text{Just} \ v'', v' \sqsubseteq v, v'' \sqsubseteq v \) \rightarrow \text{assert} \ c_1 \downarrow \sqsubseteq v, \text{assert} \ c_2 \downarrow \sqsubseteq v
\]
Proof.
\[ \text{prop } p_1 \triangleright \text{prop } p_2 \]
\[ = \lambda x \rightarrow \begin{cases} \text{Just } y & \rightarrow \text{Just } y \\ \text{Nothing} & \rightarrow \text{prop } p_2 x \end{cases} \]
\[ = \lambda x \rightarrow \begin{cases} \text{if } p_1 x \text{ then Just } x \text{ else Nothing} & \rightarrow \text{prop } p_2 x \\ \text{Nothing} & \rightarrow \text{prop } p_2 x \end{cases} \]
\[ = \lambda x \rightarrow \begin{cases} \text{if } p_1 x \text{ then Just } x \text{ else Nothing} & \rightarrow \text{prop } p_2 x \\ \text{Nothing} & \rightarrow \text{prop } p_2 x \end{cases} \]
\[ = \text{prop } (\lambda x \rightarrow p_1 x \triangleright p_2 x) \]

Lemma A.8 (Conjunction).
\[ \text{prop } p_1 \& \text{prop } p_2 = \text{prop } (\lambda x \rightarrow p_1 x \& p_2 x) \]

Proof.
\[ \text{prop } p_1 \& \text{prop } p_2 \]
\[ = \lambda x \rightarrow \text{prop } p_1 x >>\text{prop } p_2 x \]
\[ = \lambda x \rightarrow \begin{cases} \text{Just } y & \rightarrow \text{prop } p_2 y \\ \text{Nothing} & \rightarrow \text{Nothing} \end{cases} \]
\[ = \lambda x \rightarrow \begin{cases} \text{if } p_1 x \text{ then Just } x \text{ else Nothing} & \rightarrow \text{prop } p_2 x \\ \text{Nothing} & \rightarrow \text{Nothing} \end{cases} \]
\[ = \lambda x \rightarrow \text{if } p_1 x \text{ then prop } p_2 x \text{ else Nothing} \]
\[ = \lambda x \rightarrow \text{if } p_1 x \text{ then prop } p_2 x \text{ else Nothing} \]
\[ = \lambda x \rightarrow \text{if } p_1 x \text{ then Nothing} \text{ else Nothing} \]
\[ = \lambda x \rightarrow \text{prop } p_1 x \& \& \text{prop } p_2 x \]
\[ = \text{prop } (\lambda x \rightarrow p_1 x \& p_2 x) \]

A.3 Properties of the conjunction contract combinator &

Lemma A.9 (Associativity of \&).
\[ c_1 \& (c_2 \& c_3) = (c_1 \& c_2) \& c_3 \]

Proof.
\[ c_1 \& (c_2 \& c_3) \]
\[ = \lambda x \rightarrow (c_1 x \triangleright c_2 y \triangleright c_3 y) x \]
\[ = (\text{associativity of } \triangleright \text{ for the Maybe monad}) \]
\[ = \lambda x \rightarrow ((c_1 x \triangleright c_2 y) \triangleright c_3 y) x \]
\[ = \lambda x \rightarrow (c_1 x \triangleright c_2 y) \triangleright c_3 y \]
\[ = \lambda x \rightarrow c_1 x \triangleright c_2 y \triangleright c_3 y \]
\[ = \lambda x \rightarrow c_1 x \triangleright (c_2 y \triangleright c_3 y) \]
\[ = \lambda x \rightarrow c_1 x \triangleright (c_2 y \triangleright c_3) \]
\[ = \lambda x \rightarrow (c_1 x \triangleright c_2 y) \triangleright c_3 \]
\[ = (c_1 \& c_2) \& c_3 \]

Lemma A.10 (Left neutral element of \&).
\[ \text{true} \& c = c \]

Proof.
\[ \text{true} \& c \]
\[ = \lambda x \rightarrow \text{true } x \triangleright= c \]
\[ = \lambda x \rightarrow \text{Just } x \triangleright= c \]
\[ = (\text{Just } \text{return} \text{ is left identity for Maybe monad}) \]
\[ = \lambda x \rightarrow c x \]
\[ = c \]

Lemma A.11 (Right neutral element of \&).
\[ c \& \text{true} = c \]

Proof.
\[ c \& \text{true} \]
\[ = \lambda x \rightarrow c x \triangleright \text{true} \]
\[ = \lambda x \rightarrow c x \triangleright \text{Just } \text{return} \]
\[ = (\text{Just } \text{return} \text{ is right identity for Maybe monad}) \]
\[ = \lambda x \rightarrow c x \]
\[ = c \]

A.4 Properties of the disjunction contract combinator \triangleright

Lemma A.12 (Associativity of \triangleright).
\[ c_1 \triangleright (c_2 \triangleright c_3) = (c_1 \triangleright c_2) \triangleright c_3 \]

Proof.
\[ c_1 \triangleright (c_2 \triangleright c_3) \]
\[ = \lambda x \rightarrow c_1 x \triangleright (\text{c_2 y mplus c_3 y}) x \]
\[ = (\text{associativity of mplus for Maybe}) \]
\[ = \lambda x \rightarrow (c_1 x \triangleright \text{c_2 y mplus c_3 y}) x \]
\[ = \lambda x \rightarrow (c_1 y \triangleright (c_2 x \triangleright \text{c_3 y}) x \]
\[ = (c_1 \triangleright c_2) \triangleright c_3 \]

Lemma A.13 (Left neutral element of \triangleright).
\[ \text{false} \triangleright c = c \]

Proof.
\[ \text{false} \triangleright c \]
\[ = \lambda x \rightarrow \text{const Nothing } x \triangleright \text{mplus' c x} \]
\[ = \lambda x \rightarrow \text{Nothing } \text{mplus' c x} \]
\[ = (\text{Nothing } \text{mzero is left identity for Maybe}) \]
\[ = \lambda x \rightarrow c x \]
\[ = c \]

Lemma A.14 (Right neutral element of \triangleright).
\[ c \triangleright \text{false} = c \]

Proof.
\[ c \triangleright \text{false} \]
\[ = \lambda x \rightarrow \text{c x mplus Nothing' c x} \]
\[ = \lambda x \rightarrow \text{Nothing' c x} \]
\[ = (\text{Nothing' mzero is left identity for Maybe}) \]
\[ = \lambda x \rightarrow c x \]
\[ = c \]
Lemma A.15 (Idempotence of \( \triangleright \)).
\[
c \triangleright c = c
\]

Proof.
\[
c \triangleright c \equiv \lambda x \rightarrow c \ x \ 'mplus' \ const \ Nothing \ x \\
\equiv \lambda x \rightarrow c \ x \ 'mplus' \ Nothing \\
=(\text{Nothing} = \text{mzero} \text{ is right identity for Maybe}) \\
\equiv \lambda x \rightarrow c \ x \\
=c
\]

A.5 Guarded commutativity and absorption

Lemma A.16 (Guarded commutativity of disjunction).
\[
c_1 \triangleright (c_2 \triangleright c_1) = c_1 \triangleright c_2
\]

Proof.
\[
c_1 \triangleright (c_2 \triangleright c_1) \equiv \lambda x \rightarrow \text{case } c_1 \ x \ of \\
\text{Just } y \rightarrow \text{Just } y \\
\text{Nothing} \rightarrow \text{case } c_2 \ x \ of \\
\text{Just } y \rightarrow \text{Just } y \\
\text{Nothing} \rightarrow c_1 \ x \\
=(\text{case distinction } c_1 \ x = \bot, \text{Just } y, \text{Nothing}) \\
\lambda x \rightarrow \text{case } c_1 \ x \ of \\
\text{Just } y \rightarrow \text{Just } y \\
\text{Nothing} \rightarrow \text{case } c_2 \ x \ of \\
\text{Just } y \rightarrow \text{Just } y \\
\text{Nothing} \rightarrow \text{Nothing} \\
\equiv \lambda x \rightarrow \text{case } c_1 \ x \ of \\
\text{Just } y \rightarrow \text{Just } y \\
\text{Nothing} \rightarrow c_2 \ x \\
=c_1 \triangleright c_2
\]

Lemma A.17 (Absorption 1).
\[
c_1 \triangleright (c_1 \& c_2) = c_1
\]

Proof.
\[
c_1 \triangleright (c_1 \& c_2) \equiv \lambda x \rightarrow \text{case } c_1 \ x \ of \\
\text{Just } y \rightarrow \text{Just } y \\
\text{Nothing} \rightarrow \text{case } c_1 \ x \ of \\
\text{Just } y \rightarrow \text{Just } y \\
\text{Nothing} \rightarrow \text{Nothing} \\
=(\text{case distinction } c_1 \ x = \bot, \text{Just } y, \text{Nothing}) \\
\lambda x \rightarrow \text{case } c_1 \ x \ of \\
\text{Just } y \rightarrow \text{Just } y \\
\text{Nothing} \rightarrow \text{Nothing} \\
\equiv c_1
\]

A.6 Corollaries

Corollary A.18.
\[
\text{false} \& c = \text{false}
\]

Proof.
\[
\text{false} \& c \equiv (\text{Lemma A.13, left neutral element of } \&) \\
\text{false} \triangleright (\text{false} \& c) \equiv (\text{Lemma A.17, absorption 1}) \\
\text{false}
\]

Corollary A.19.
\[
\text{true} \triangleright c = \text{true}
\]

Proof.
\[
\text{true} \triangleright c \equiv (\text{Lemma A.10, left neutral element of } \&) \\
\text{true} \triangleright (\text{true} \& c) \equiv (\text{Lemma A.17, absorption 1}) \\
\text{true}
\]

A.7 Properties of the function contract combinator \(\triangleright\triangleright\)

Lemma A.20 (Relation with true).
\[
\text{true} \triangleright\triangleright \text{true} = \text{true}
\]

Proof.
\[
\text{true} \triangleright\triangleright \text{true} \equiv \lambda f \rightarrow \text{Just } (f \ 'seq' \ (\text{assert true } . f . \text{assert true})) \\
\equiv \lambda f \rightarrow \text{Just } (f \ 'seq' \ (\text{id } . f . \text{id})) \\
\equiv \lambda f \rightarrow \text{Just } f \\
=\text{true}
\]

Lemma A.21 (false postcondition).
\[
c_1 \triangleright\triangleright \text{false} = c_2 \triangleright\triangleright \text{false}
\]

Proof.
\[
c_1 \triangleright (c_1 \& c_2) \equiv \lambda x \rightarrow \text{case } c_1 \ x \ of \\
\text{Just } y \rightarrow \text{Just } y \\
\text{Nothing} \rightarrow \text{case } c_2 \ z \ of \\
\text{Just } z \rightarrow c_2 \ z \\
\text{Nothing} \rightarrow \text{Nothing} \\
=(\text{case distinction } c_1 \ x = \bot, \text{Just } y, \text{Nothing}) \\
\lambda x \rightarrow \text{case } c_1 \ x \ of \\
\text{Just } y \rightarrow \text{Just } y \\
\text{Nothing} \rightarrow \text{Nothing} \\
\equiv c_1
\]
Proof.
\[ c_1 \rightarrow \text{false} \]
\[ = \lambda f \rightarrow \text{Just} (f \ '\text{seq'} (\text{assert}\ \text{false} . f . \text{assert}\ c_1)) \]
\[ = \lambda f \rightarrow \text{Just} (f \ '\text{seq'} (\text{const} (\text{error} "\ldots") . f . \text{assert}\ c_1)) \]
\[ = \lambda f \rightarrow \text{Just} (f \ '\text{seq'} \text{const} (\text{error} "\ldots")) \]
\[ = \lambda f \rightarrow \text{Just} (f \ '\text{seq'} (\text{assert}\ \text{false} . f . \text{assert}\ c_2)) \]
\[ = c_2 \rightarrow \text{false} \]

Lemma A.22 (Distribution with \&).
\[ (c_1 \rightarrow c_2) \& (c_3 \rightarrow c_4) = (c_3 \& c_1) \rightarrow (c_2 \& c_4) \]

Proof.
\[ (c_1 \rightarrow c_2) \& (c_3 \rightarrow c_4) \]
\[ = \lambda f \rightarrow (c_1 \rightarrow c_2) \ f \rightarrow (c_3 \rightarrow c_4) \]
\[ = \lambda f \rightarrow \text{Just} (f \ '\text{seq'} (\text{assert}\ c_2 . f . \text{assert}\ c_1)) \]
\[ = \lambda f \rightarrow (c_3 \rightarrow c_4) \ (f \ '\text{seq'} (\text{assert}\ c_2 . f . \text{assert}\ c_1)) \]
\[ = \lambda f \rightarrow \text{Just} (f \ '\text{seq'} (\text{assert}\ c_4 . \text{assert}\ c_2 . f . \text{assert}\ c_1)) \]
\[ = \text{Lemma A.5} \]
\[ \lambda f \rightarrow \text{Just} (f \ '\text{seq'} (\text{assert}\ (c_2 \& c_4) . f . \text{assert}\ (c_3 \& c_1))) \]
\[ = (c_3 \& c_1) \rightarrow (c_2 \& c_4) \]

Lemma A.23 (Non-distribution with |>),
\[ (c_1 \rightarrow c_2) |> (c_3 \rightarrow c_4) = c_1 |> c_2 \]

Proof.
\[ (c_1 \rightarrow c_2) |> (c_3 \rightarrow c_4) \]
\[ = \lambda f \rightarrow (c_1 \rightarrow c_2) \ f \ '\text{mplus'} (c_3 \rightarrow c_4) \ f \]
\[ = \lambda f \rightarrow \text{Just} (f \ '\text{seq'} (\text{assert}\ c_2 . f . \text{assert}\ c_1)) \]
\[ '\text{mplus'} (c_3 \rightarrow c_4) \ f \]
\[ = \text{MonadPlus law for Just = return} \]
\[ \lambda f \rightarrow \text{Just} (f \ '\text{seq'} (\text{assert}\ c_2 . f . \text{assert}\ c_1)) \]
\[ = c_1 \rightarrow c_2 \]

B. Properties of the contract combinator \text{io}

The following are proofs for Section 7.2.

Lemma B.1.
\[ \text{io} c_1 \& \text{io} c_2 = \text{io} (c_1 \& c_2) \]

Proof. Structural induction on the value \( V \).

C. Two definitions of the list contract combinator

The two list contract combinators are not equal, but we can prove that they are equal in all contract contexts surrounded by \text{assert}. Basically the results \( \bot \) and \text{Just} \( \bot \) of a contract cannot be distinguished. We denote this equivalence by \( \equiv \).

The next lemma is not only of interest for the subsequent lemma. However, the two expressions will give different fault explanations if a contract is violated.

Lemma C.1.
\[ \text{assert} (\text{list}\ c) \ V = \text{map}\ (\text{assert}\ c) \ V \]

Proof. Structural induction on the value \( V \).
\[ V = \perp : \]
\[
\text{assert (list c)} \perp = \text{assert (pNil |> pCons c (list c))} \perp
\]
\[
\text{=map (assert c)} \perp
\]
\[ V = [] : \]
\[
\text{assert (list c)} [] = \text{assert (pNil |> pCons c (list c))} [] = []
\]
\[
\text{=map (assert c)} []
\]
\[ V = W:WS : \]
\[
\text{assert (list c) (W:WS)} = \text{assert (pNil |> pCons c (list c)) (W:WS)}
\]
\[
\text{=assert c W : assert (list c) WS}
\]
\[
\text{=(induction hypothesis)}
\]
\[
\text{assert c W : map (assert c) WS}
\]
\[
\text{=map (assert c) (W:WS)}
\]

\[ \text{Lemma C.2.} \]
\[
\text{list c V} \equiv \text{list’ c V}
\]

\[ \text{Proof.} \] Structural induction on the value \( V \).
\[ \text{V} = \perp : \]
\[
\text{list c} \perp = (\text{pNil |> pCons c (list c)} \perp \perp
\]
\[
\equiv \text{Just} \perp
\]
\[
\equiv \text{list’ c} \perp
\]
\[ \text{V} = [] : \]
\[
\text{list c} [] = (\text{pNil |> pCons c (list c)} [])
\]
\[
\equiv \text{Just}]
\[
\equiv \text{list’ c} []
\]
\[ \text{V} = W:WS : \]
\[
\text{list c (W:WS)} = (\text{pNil |> pCons c (list c)) (W:WS)}
\]
\[
\equiv \text{Just (assert c W : assert (list c) WS)}
\]
\[
\equiv (\text{Lemma C.1)}
\]
\[
\text{Just (assert c W : map (assert c) WS)}
\]
\[
\equiv \text{Just (map (assert c) (W:WS))}
\]
\[
\equiv \text{list’ c (W:WS)}
\]

\[ \text{D. Some Properties that do not hold} \]

Contracts do not form a distributive lattice with \& and \|. So a number or properties that might be expected do not actually hold. Comparing with the properties of the lazy Booleans is helpful in identifying which properties hold and which do not.

\[ \text{Lemma D.1 (Commutativity).} \]
\[ & \text{is not commutative.} \]

\[ \text{Proof.} \] For example
\[
\perp \& \text{false} = \lambda x \rightarrow \perp
\]
\[
\text{false} \& \perp = \lambda x \rightarrow \text{Nothing}
\]

So
\[
(\perp \& \text{false}) \| \text{true} = \lambda x \rightarrow \perp
\]
\[
(\text{false} \& \perp) \| \text{true} = \lambda x \rightarrow \text{Just x}
\]

Hence
\[
\text{assert ((\perp \& \text{false}) \| \text{true}) e} = \perp
\]
\[
\text{assert (\text{false} \& \perp) \| \text{true}) e} = e
\]

Similarly \&\& is not commutative:
\[
\perp \&\& \text{False} = \perp \neq \text{False} = \text{False} \&\& \perp
\]

\[ \text{Lemma D.2 (Commutativity).} \]
\[ |\| \text{is not commutative.} \]

\[ \text{Proof.} \] For example
\[
\perp |\| \text{true} = \lambda x \rightarrow \perp
\]
\[
\text{true} |\| \perp = \lambda x \rightarrow \text{Just x}
\]

Hence
\[
\text{assert (\perp |\| \text{true}) e} = \perp
\]
\[
\text{assert (\text{true} |\| \perp) e} = e
\]

Similarly \|\| is not commutative.
\[
\perp \|\| \text{True} = \perp \neq \text{True} = \text{True} \|\| \perp
\]

\[ \text{Lemma D.3 (Distributivity 1).} \]
\[ (c_1 \& c_2) \| \text{true} \neq (c_1 \| c_2) \& (c_2 \| c_3) \]

\[ \text{Proof.} \]
\[
(\text{false} \| \text{true}) \& (\perp \| \text{true}) = \lambda x \rightarrow \perp
\]
\[
(\text{false} \& \perp) \| \text{true} = \lambda x \rightarrow \text{Just x}
\]

Note that also
\[
(b_1 \| b_2) \& \& b_3 \neq (b_1 \& \& b_3) \| (b_2 \& \& b_3)
\]

because
\[
(\text{False} \| \text{True}) \& \& (\perp \| \text{True}) = \perp
\]
\[
\equiv \text{True}
\]
\[
\equiv (\text{False} \& \perp) \| \text{True}
\]

\[ \text{Lemma D.4 (Distributivity 2).} \]
\[ (c_1 \& c_2) \| (c_2 \& c_3) \neq (c_1 \| c_2) \& c_3 \]

\[ \text{Proof.} \]
\[
(\text{true} \& \text{false}) \| (\perp \& \text{false}) = \lambda x \rightarrow \perp
\]
\[
(\text{true} \| \perp) \& \text{false} = \lambda x \rightarrow \text{Nothing}
\]

\[ \]

Practical Typed Lazy Contracts, ICFP 2012 15 2012/7/23
Note that also

\[(b_1 \&\& b_3) || (b_2 \&\& b_3) \neq (b_1 || b_2) \&\& b_3\]

because

\[(\text{True} \&\& \text{False}) || (\bot \&\& \text{False})\]

\[= \bot\]

\[\square \text{False}\]

\[= (\text{True} || \bot) \&\& \text{False}\]

**Lemma D.5.**

\[\text{false} \Rightarrow \text{false} \neq \text{false}\]

**Proof.**

\[(\text{false} \Rightarrow \text{true}) \Leftarrow \Rightarrow \equiv f\]

\[((\text{false} \Rightarrow \text{false}) \Rightarrow \text{true}) \Leftarrow \Rightarrow \equiv \text{const (error "..."} )\]

\[\square\]