1. Introduction

1.1 What is Blue?

Blue is an object-oriented programming language especially designed for teaching. Its emphasis is on conceptual clarity and consistency, suitability for first year teaching and support for “good” software design.

1.2 The Blue Environment

Blue is intended to be used in an integrated programming environment. The environment itself will not be described in this document, but is important for the overall character of Blue programming. In order to obtain a complete picture of programming in Blue, the reader is encouraged also to read “The Blue Programming Environment”.

1.3 About this Document

This document describes the Blue language. The Blue programming language is part of the Blue programming system. The language and some of the most important standard libraries (such as those for standard I/O) will be described in this document. However, most standard libraries will not be included here, but will be described in a separate document “The Blue Libraries”.

It is assumed that the reader is familiar with a Pascal- or C-like language and with object-oriented concepts. This document does not give an introduction to object-oriented programming in general. It is assumed that the most basic ideas of object-orientation are known.

At some points this document also includes thoughts and reasoning beyond pure description of the language. These comments typically include reasons for specific decisions made during language design. They may be skipped by readers interested only in the language, or read by someone interested in the design process itself.

Comments about reasons for specific design decisions are set in greyed boxes like this and may be skipped without losing information about the Blue language.

Descriptions of language constructs always include a formal syntax description in EBNF. A complete syntax description in EBNF can be found in Appendix A.
1.4 **Blue Programming**

Blue does not include a concept of programs. All programming is performed using classes. All code is part of a class. The classical idea of a program is represented in Blue as a collection of classes (called a *project* in Blue). Running a program is done by creating an instance of a class, and calling interface routines of that class.

A simulation of a classical “program” (a single entry point executable) can easily be achieved by having one designated top-level object with one interface routine which in turn creates and calls all other objects involved in the execution of the program. (This is the standard execution model in most existing object-oriented languages.) It is, however, only a small subset of the possible ways to execute code in Blue.

In the Blue environment, *all* classes in the project may be used to interactively create objects of that class and *all* those objects can immediately be used to call their interface routines (that is: without the necessity of writing test program shells). This changes programming in at least two ways: it allows very flexible (incremental) testing and incremental software development. All low level classes may immediately be used and tested for correct behaviour. It also allows the construction of software with more than one entry point, avoiding common awkward dispatch mechanisms invoked by command line options.

Classes are types\(^1\). So all programming comes down to the definition of types. This is why type declarations are the first major part of this document: They specify the overall structure of Blue programs.

2. **Aliases**

All types in Blue are classes and all data are objects. This general rule simplifies the language design. There are, however, a number of data types for which it is convenient to use syntax other than the Blue object call to perform one of their operations. The reason for this can be:

- Another syntax is commonly used and is therefore more intuitive (e.g. 3 + 5 for integer addition, rather than 3.add (5)).
- Another syntax simplifies use of elementary constructs which should be used by beginners before the underlying language concepts need to be understood, e.g.  
  ```
  print ("result\("\, 42)
  ```
  instead of  
  ```
  output.write ("result\=".concat (42.toString))
  ```
- Another syntax is more convenient (usually because it is shorter, see above).

\(^1\) It is not strictly true that classes and types are the same, but most of the time each class represents exactly one type. This is not true when it comes to generic classes (which do not represent a type, but rather a *type pattern*). The difference will be explained in detail in section 17. For most of this document, however, this simplification is not a problem and makes expressing things a bit easier.
For these reasons, several operations on the predefined classes are supported by special syntax. This special syntax is allowed in addition to the standard object-call syntax generally available for all classes, and is referred to as *aliases*.

The prime reason for the introduction of aliases is to make the reading and writing of simple programs performing elementary tasks easy. Aliases provide an easy, intuitive syntax for the most common operations and considerably increase the ease with which Blue can be used by beginners.

Aliases will be learnt as statements or expressions in their own right by beginners, making it unnecessary to understand all underlying concepts right from the start. The expert programmer or compiler implementor, however, will appreciate the unifying concept for all data types.

Note that aliases are a pure *syntactic* addition which does not add any functionality to the language. They do not affect the semantics or the theoretical language description of Blue (although they are part of the language), and are purely intended to increase readability and intuitivity of statements.

Some common aliases and their resolutions are:

<table>
<thead>
<tr>
<th>alias</th>
<th>resolution</th>
</tr>
</thead>
<tbody>
<tr>
<td>3+6</td>
<td>3.add (6)</td>
</tr>
<tr>
<td>b1 or b2</td>
<td>b1.or (b2)</td>
</tr>
<tr>
<td>a[i]</td>
<td>a.getElem (i)</td>
</tr>
<tr>
<td>str (num)</td>
<td>num.toString</td>
</tr>
<tr>
<td>str (a, b, ...)</td>
<td>a.toString.concat (b.toString.concat (...))</td>
</tr>
<tr>
<td>print (a, b)</td>
<td>output.write ( str (a.b) )</td>
</tr>
</tbody>
</table>

The list of existing aliases is short and fixed – programmers cannot define additional aliases.

Where aliases exist for a construct explained in this document, they will be mentioned when that construct is introduced. A full list of aliases is given in Appendix B.

### 3. Classes and their Operations

All types derive from classes. Classes for the most commonly used data types are predefined in the language. The language offers customised syntax for these classes.

The predefined classes are:

- Integer
- Real
- Boolean
- String
- Array `<T>` (generic)

Two different kinds of classes exist: manifest classes and dynamic classes. The following sections will introduce manifest and dynamic classes, describe operations that are available on all classes, and then list the predefined classes and their operations. Finally, user defined classes are described.
3.1 Manifest vs. Dynamic Classes

Manifest classes are classes where all objects are known statically. The objects pre-exist with the definition of the class and do not have to be created. The manifest classes are Integer, Real, Boolean, String and Enumeration classes. The first four of these are predefined and all values are known to the Blue compiler. Enumeration classes are user defined. The definition of such a class consists of an enumeration of all existing objects of that class, simultaneously creating a named reference to each object.

The literal ‘2’, for instance, is a reference to the unique integer object with the value 2. The code segment

\[
\begin{align*}
a & := 2 \\
b & := 2
\end{align*}
\]

assigns references to the same object to a and b. Only one integer object with the value 2 exists. This does not create two distinct objects.

All literals are constant references to objects of manifest classes.

Dynamic classes are classes where, rather than listing all objects, a creation method for objects is specified. Dynamic classes are arrays and user defined general classes.

With the definition of a dynamic class, no object is created automatically. The user has to execute an explicit create operation to create objects of these types.

Care must be taken not to confuse this with pointer and non-pointer types in other languages. In Blue, all variables hold references to objects. An integer variable holds, when assigned a value, the reference to that integer object. Thus the object model is simple: only references to objects exist. The difference between manifest and dynamic objects affects only the time and method of creation of the objects of the class, not the mechanism by which they are referenced.

3.2 General Operations

This section describes operation that are common to all classes.

The general operations are:

\[
\begin{align*}
:= & \quad \text{assignment} \\
= & \quad \text{comparison (equality)} \\
<> & \quad \text{comparison (inequality)}
\end{align*}
\]

3.2.1 Assignment (:=)

The assignment (:=) assigns an object reference to a variable (variables always store object references). The object itself is not copied.

Example:
Consider variables a and b which are declared to be of some user defined class:
Then the code segment
\[
a := \text{create Myclass}
b := a
\]
causes \(a\) and \(b\) to reference the same object. Changes to \(a\) will be visible using \(b\).

### 3.2.2 Comparison (=, <>)

The comparison compares two values (variables or expression results). Values are always object references.

**Example:**

```plaintext
var
a: myclass
b: myclass
result: Boolean

... 
result := b = a
```

\(result\) will be true if \(a\) and \(b\) reference the same object. It will not be true, if \(a\) and \(b\) reference different objects in identical states.

\(a <> b\) is true, if \(a\) and \(b\) do not reference the same object.

### 3.3 Predefined Types

The predefined types and their operations are as follows. Many of the operations mentioned here are aliases. The full list of operations and their aliases are given in the individual sections for each type below.

<table>
<thead>
<tr>
<th>Integer</th>
<th>Real</th>
<th>Boolean</th>
<th>String</th>
<th>Array</th>
</tr>
</thead>
<tbody>
<tr>
<td>(n + m)</td>
<td>(x + y)</td>
<td>not (b)</td>
<td>s.length</td>
<td>a ([n])</td>
</tr>
<tr>
<td>(n - m)</td>
<td>(x - y)</td>
<td>(a \text{ and } b)</td>
<td>str ((s1, s2))</td>
<td>a.size</td>
</tr>
<tr>
<td>(-n)</td>
<td>(-x)</td>
<td>(a \text{ or } b)</td>
<td>s.substr ((n, m))</td>
<td>a.setSize ((n))</td>
</tr>
<tr>
<td>(n \times m)</td>
<td>(x \times y)</td>
<td>(s[n])</td>
<td></td>
<td></td>
</tr>
<tr>
<td>(n \div m)</td>
<td>(x / y)</td>
<td>s1.find ((s2, n))</td>
<td></td>
<td></td>
</tr>
<tr>
<td>(n \mod m)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>(n ^ m)</td>
<td>(x ^ y)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>(\sqrt{x})</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>(n &gt; m)</td>
<td>(x &gt; y)</td>
<td>s1 &lt; s2</td>
<td></td>
<td></td>
</tr>
<tr>
<td>(n &lt; m)</td>
<td>(x &lt; y)</td>
<td>s1 &gt; s2</td>
<td></td>
<td></td>
</tr>
<tr>
<td>(n \geq m)</td>
<td>(x \geq y)</td>
<td>s1 &lt;= s2</td>
<td></td>
<td></td>
</tr>
<tr>
<td>(n \leq m)</td>
<td>(x \leq y)</td>
<td>s1 &gt;= s2</td>
<td></td>
<td></td>
</tr>
<tr>
<td>str ((n))</td>
<td>str ((x))</td>
<td>str ((b))</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

*Integer*, *Real*, *Boolean*, and *String* are manifest classes. *Array* is a dynamic class.
The predefined types are now discussed in more detail. For a full class interface of the predefined types see Appendix D.

### 3.3.1 Integer

The class Integer stores whole numbers. Integer is a manifest class.

The operations defined on integers are:

<table>
<thead>
<tr>
<th>alias</th>
<th>operation</th>
<th>return type</th>
<th>meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>n + m</td>
<td>n.add (m)</td>
<td>Integer</td>
<td>addition</td>
</tr>
<tr>
<td>n - m</td>
<td>n.sub (m)</td>
<td>Integer</td>
<td>subtraction</td>
</tr>
<tr>
<td>-n</td>
<td>n.neg</td>
<td>Integer</td>
<td>negation</td>
</tr>
<tr>
<td>n * m</td>
<td>n.mult (m)</td>
<td>Integer</td>
<td>multiplication</td>
</tr>
<tr>
<td>n div m</td>
<td>n.div (m)</td>
<td>Integer</td>
<td>Integer division</td>
</tr>
<tr>
<td>n mod m</td>
<td>n.mod (m)</td>
<td>Integer</td>
<td>remainder</td>
</tr>
<tr>
<td>n ^ m</td>
<td>n.pow (m)</td>
<td>Integer</td>
<td>power</td>
</tr>
<tr>
<td>n &gt; m</td>
<td>n.greater (m)</td>
<td>Boolean</td>
<td>greater than</td>
</tr>
<tr>
<td>n &lt; m</td>
<td>n.less (m)</td>
<td>Boolean</td>
<td>less than</td>
</tr>
<tr>
<td>n &gt;= m</td>
<td>n.greaterEq (m)</td>
<td>Boolean</td>
<td>greater or equal</td>
</tr>
<tr>
<td>n &lt;= m</td>
<td>n.lessEq (m)</td>
<td>Boolean</td>
<td>less or equal</td>
</tr>
<tr>
<td>str (n)</td>
<td>n.toString</td>
<td>String</td>
<td>conversion to String</td>
</tr>
</tbody>
</table>

In addition to these operations, integer expressions can always be used where real expressions are expected. An implicit conversion to Real takes place in that case.

Integer literals are written as usual:

- 42
- -99

Constants are predefined for the smallest and largest available integer value in every specific implementation:

- MAXINT: largest representable integer
- MININT: smallest representable integer

The actual value of MAXINT and MININT are implementation dependent.

### 3.3.2 Real

The class Real stores floating point numbers. Real is a manifest class.

The operations defined on reals are:
### Real

<table>
<thead>
<tr>
<th>alias</th>
<th>operation</th>
<th>return type</th>
<th>meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>x + y</td>
<td>x.add (y)</td>
<td>Real</td>
<td>addition</td>
</tr>
<tr>
<td>x - y</td>
<td>x.sub (y)</td>
<td>Real</td>
<td>subtraction</td>
</tr>
<tr>
<td>-x</td>
<td>x.neg</td>
<td>Real</td>
<td>negation</td>
</tr>
<tr>
<td>x * y</td>
<td>x.mult (y)</td>
<td>Real</td>
<td>multiplication</td>
</tr>
<tr>
<td>x / y</td>
<td>x.div (y)</td>
<td>Real</td>
<td>division</td>
</tr>
<tr>
<td>x ^ y</td>
<td>x.pow (y)</td>
<td>Real</td>
<td>power</td>
</tr>
<tr>
<td></td>
<td>x.sqrt</td>
<td>Real</td>
<td>square root</td>
</tr>
<tr>
<td></td>
<td>x.trunc</td>
<td>Integer</td>
<td>truncation to integer</td>
</tr>
<tr>
<td></td>
<td>x.round</td>
<td>Integer</td>
<td>round to closest Integer</td>
</tr>
<tr>
<td></td>
<td>str (x)</td>
<td>String</td>
<td>String conversion to String</td>
</tr>
</tbody>
</table>

Real literals are written as:

- 2.0
- 0.111
- -0.001

Note that an implicit conversion takes place when an integer is provided where a real value is expected. Thus it is legal to write 2 instead of 2.0, where the value of 2 will be implicitly converted to the real value 2.0.

#### 3.3.3 Boolean

Variables of class Boolean store boolean values (i.e. `true` or `false`). Boolean is a manifest class.

The operations defined on booleans are:

<table>
<thead>
<tr>
<th>alias</th>
<th>operation</th>
<th>return type</th>
<th>meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>not b</td>
<td>b.invert</td>
<td>Boolean</td>
<td>negation</td>
</tr>
<tr>
<td>a and b</td>
<td>a.and (b)</td>
<td>Boolean</td>
<td>logical and</td>
</tr>
<tr>
<td>a or b</td>
<td>a.or (b)</td>
<td>Boolean</td>
<td>logical or</td>
</tr>
<tr>
<td>str (b)</td>
<td>b.toString</td>
<td>String</td>
<td>conversion to String</td>
</tr>
</tbody>
</table>

The boolean literals are:

- `true`
- `false`

Note that, since "true" and "false" are keywords, and the case of keywords is insignificant, variation of case such as True or TRUE are also recognised.

#### Partial Evaluation

Boolean expressions are incrementally evaluated. Evaluation starts from the left (subject to operator precedence rules) and stops as soon as the result can be determined.

**Examples:**

Consider the expression

```
a and b
```

If `a` evaluates to `false`, the evaluation of the expression will terminate and return the result `false`. `b` is only evaluated if `a` evaluates to `true`.

Likewise, in the expression
a or b

b is only evaluated if a is false. If a is true, the result of the expression will be determined without evaluating b.

This allows the following code segments to be written:

```
if obj<>nil and obj.size>N then ...
```

or

```
if n<=a.size and a[n]>0 then ...
```

3.3.4 String

Strings hold sequences of characters. Strings are manifest. Thus a String variable can reference any possible String. There is no notion of a specific (maximum) number of characters or a specific memory section associated with a String variable. String literals are written using double quotes (").

Example:

```
s1, s2: String

... 

s1 := "Da steh' ich nun, ich armer Tor"
s2 := s1
```

The effect of class String being a manifest type is that it behaves like having value semantics. In this example, subsequent changes to s1 will not affect s2. (A change of s1 only makes s1 reference another string.)

Characters and substrings can be extracted by using indices. String indices always start at 1.

There is no character type in Blue. Characters are strings with length one.

The most common operations on strings are:

<table>
<thead>
<tr>
<th>alias</th>
<th>operation</th>
<th>return type</th>
<th>meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>str (s1, s2)</td>
<td>s.length</td>
<td>Integer</td>
<td>number of characters in s</td>
</tr>
<tr>
<td>s1.concat(s2)</td>
<td>String</td>
<td>returns concatenation of s1 and s2</td>
<td></td>
</tr>
<tr>
<td>s.substring (n, m)</td>
<td>String</td>
<td>substring of s starting at n with length m</td>
<td></td>
</tr>
<tr>
<td>s[n]</td>
<td>s.substring (n, 1)</td>
<td>String</td>
<td>character in s at position n</td>
</tr>
<tr>
<td>s2.find (s1, n)</td>
<td>Integer</td>
<td>position of s1 in s2 after n</td>
<td></td>
</tr>
<tr>
<td>s1.insert (s2, n)</td>
<td>String</td>
<td>inserts s2 into s1 at pos. n</td>
<td></td>
</tr>
<tr>
<td>s.delete (n, m)</td>
<td>String</td>
<td>delete m characters, starting at n</td>
<td></td>
</tr>
<tr>
<td>s1 &lt; s2</td>
<td>s1.less (s2)</td>
<td>Boolean</td>
<td>s1 less than s2 ?</td>
</tr>
<tr>
<td>s1 &gt; s2</td>
<td>s1.greater (s2)</td>
<td>Boolean</td>
<td>s1 greater than s2 ?</td>
</tr>
<tr>
<td>s1 &lt;= s2</td>
<td>s1.lessEq (s2)</td>
<td>Boolean</td>
<td>s1 less or equal s2 ?</td>
</tr>
<tr>
<td>s1 &gt;= s2</td>
<td>s1.greaterEq (s2)</td>
<td>Boolean</td>
<td>s1 greater or equal s2 ?</td>
</tr>
<tr>
<td>str (s)</td>
<td>s.toString</td>
<td>String</td>
<td>convert to upper case</td>
</tr>
<tr>
<td>s.toUpper</td>
<td>String</td>
<td>convert to lower case</td>
<td></td>
</tr>
<tr>
<td>s.toLower</td>
<td>String</td>
<td>equality (ignore case)</td>
<td></td>
</tr>
<tr>
<td>s.caseEqual (s2)</td>
<td>Boolean</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
Special characters
A number of escape sequences are provided to specify non-printable characters:

- "\t" tab character
- "\n" end-of-line character
- "\\" escape character
- "\\nnn" ISO character nnn (decimal)

Long strings
If a string is too long to be written in one line in the editor, it can be broken into several parts:

s := "This is one string"
    "(and this is still the same string)"

String literals that are separated by whitespace only are processed as one string. The neighbouring quotes and the whitespace in between are ignored. The strings

"one" "two" "three"

and

"one"
"two"
"three"

and

"one two three"

are identical. Every string literal (every part of the string) must end on the same line it starts on.

For a full interface of class String, see Appendix D.

The decision to make the class String a manifest class may surprise initially. Programmers tend to think about strings as memory buffers holding characters. Objects of manifest classes can not change - this contradicts the model programmers used to some other languages have of strings: they can change. After careful examination, however, it turns out that we usually think of strings not as different objects. Making strings a dynamic class would lead to the effect that in the code fragment

```plaintext
a := "marvin"
if a = "marvin" then ...
```

the expression `a = "marvin"` results in false (because a literal string is then an object creator and `a` and "marvin" are two distinct objects). Intuitively, we do not think about strings as independent objects, where two distinct objects with the same value can exist. If they have the same value, then they are the same. And this is exactly the behaviour of manifest objects.

All string functions that change a string (such as toUpper) do not change the current string, but rather return a reference to another string that contains the required text.
3.3.5 **Array**

Arrays are the only predefined dynamic class. Therefore array objects need to be explicitly created.

Arrays are sequences of homogenous elements. The class specification is generic. (For details about generic classes, see section 17.) The element type is specified when declaring a variable; the index type is always Integer. The size of an array is dynamic. The boundaries are not statically fixed and are not part of the type.

Example:

```plaintext
var a: Array <Boolean>

... a := create Array<Boolean> (20)
```

This example creates an array of 20 booleans. Array indices always start at 1, so this array contains elements from a[1] to a[20].

Arrays can then be assigned like all dynamic classes.

Operation on arrays:

<table>
<thead>
<tr>
<th>alias</th>
<th>operation</th>
<th>return type</th>
<th>meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>a [n]</td>
<td>a.getElem (n)</td>
<td>t</td>
<td>element at index n</td>
</tr>
<tr>
<td>a [n]</td>
<td>a.putElem (n, t)</td>
<td>–</td>
<td>set element at index n</td>
</tr>
<tr>
<td>a.init</td>
<td>–</td>
<td>Integer</td>
<td>set all elements in a to v</td>
</tr>
<tr>
<td>a.size</td>
<td>–</td>
<td>Integer</td>
<td>number of elements in a</td>
</tr>
<tr>
<td>a.setSize (n)</td>
<td>–</td>
<td>–</td>
<td>change array size to n. If n &gt; old n, the new elements are undefined.</td>
</tr>
</tbody>
</table>

Apart from these operations, literal Array constructors can be specified.

Example:

```plaintext
var a: Array <Integer>

... a := [23, 2, 42]
```


Elements of an array can be of any type. In particular, arrays of arrays are possible:

```plaintext
var
    matrix: Array <Array<Integer>>

i1, i2: Integer

... i1 := matrix [2][2]
```

The array bracket alias is special in that it can appear on the right or the left side of an assignment. Dependent on its position, it translates to either `getElem` or `putElem`.

```plaintext
num := a[i] translates to num := a.getElem(i)
a[i] := num translates to a.putElem(i, num)
```

The bracket alias therefore is the only alias that is not a simple macro.
3.4 User Defined Types

User defined classes can be either general classes (user defined dynamic classes) or enumeration classes (user defined manifest classes).

3.4.1 General classes

The definition of a general class is the closest thing to a “program” in procedural languages. A class typically consists of some internal data (encapsulated), internal routines, a creation routine and interface routines. The interface cannot contain variables.

**BNF:**

```
class-decl ::= class identifier ["<" generics-list "] is class-definition

generics-list ::= gen-param { , gen-param }
gen-param ::= identifier [ is identifier ]

class-definition ::= [ identifier ] general-class-decl
                          | ...

general-class-decl ::= class-comment
                          uses [ ident-list ]
                          [ internal
                          [ const const-decls ]
                          [ var var-decls ]
                          [ routines routine-decls ] ]
                          interface
                          [ creation ["(" parameter-list ")" ]
                          routine-body ]
                          [ routines routine-decls ]
                          [ invariant condition ]
                          end class
```

Example:

```plaintext
class Rectangle is
   == Author: M. Kölling
   == Date: April 1995
   == Version: 1.0
   == Short: Graphical representation of a rectangle.
   ==
   == Class Rectangle represents a rectangle with specified
   == coordinates and colour that can be ...
   uses Point, Colour
   
   internal
   
   var
       top_left: Point
       bottom_right: Point
```

11
fill_col: Colour
border_col: Colour

interface

creation (tl: Point, br: Point) is
  == Create rectangle at coordinates defined by tl (top-left)
  == and br (bottom-right). Default colours are: fill white,
  == border black.
  do
    top_left := tl
    bottom_right := br
    fill_col := white
    border_col := black
  end creation

routines

move (dx: Integer, dy: Integer) is
  == Move rectangle by distance defined by dx, dy.
  do
    top_left.move (dx, dy)
    bottom_right.move (dx, dy)
  post
    size = old size
  end move

size -> (width: Integer, height: Integer) is
  == Return size of rectangle in width and height.
  do
    width := bottom_right.x - top_left.x
    height := bottom_right.y - top_left.y
  end size

invariant
  (top_left.x < bottom_right.x) and
  (top_left.y < bottom_right.y)

end class

There is no separate interface definition for a class. The interface of a class is a restricted view of
the class attributes (the internals are hidden) and it is presented to the user as a special view that is produced by a tool in the Blue
Programming Environment from the full class definition.

For details on generic classes, see section 17.

- The name is always first in all declarations. The reason for this is that a class definition is mostly used to look up interface features (members) and their characteristics (parameters, types). The feature should be easy to find. This is achieved by placing the name first, making it easy to read down a column of identifiers. In C-type syntax (type and some keywords first) the identifier is shifted somewhere towards the middle of the line, which makes it harder to find. The name is followed by the accessor
definition (type, parameters) showing the syntax for usage, and the implementation definition (routine body), which is class internal information.

- Creation is separate to make clear it is not a normal routine. Create looks similar to a routine but has a special syntax and semantics. It is separated from normal interface procedures and put into a special position in the class declaration to make it easy to find by enforcing a certain position for it and to symbolise its special meaning.

- Comments are compulsory. Writing comments is not seen as a luxury but as part of programming. Therefore the compiler also deals (as well as it can) with proper commenting. This includes a class comment, containing certain keywords, used by the library browser, and comments for routines.

- Eliminating separate interface files avoids the danger of inconsistency and repetition of code.

3.4.2 Enumeration classes

Enumeration classes are the only user defined manifest classes. In a manifest class all objects exist automatically and no other objects of this class can be created. Every object is referenced by a named constant.

```
class-decl ::= class identifier [ "<" generics-list ">"] is class-definition
class-definition ::= Enumeration enum-class-decl | ...
enum-class-decl ::= class-comment manifest ident-list
                 end class
ident-list ::= identifier ["," identifier ]
```

Example:
```
class Colour is Enumeration
    == Simple enumeration type for basic colours
    manifest red, white, blue
end class
```

This example defines and creates three objects which are referenced through the constant identifiers red, white and blue. In another class a variable can be defined of class Colour, and one of the constant references can be assigned to that variable.

Example:
```
var col1, col2: Colour
...
```
col := red
col2 := blue

Note that, because no new objects are created at any time, after assigning red to both col1 and col2, \((\text{col1} = \text{col2})\) is true.

Enumeration classes inherit from the predefined class `Enumeration`. `Enumeration` defines two routines `pred` and `succ` and the standard routine `toString`, which can then be applied to all enumeration types. It also provides a routine `ord` for conversion to Integer. For a full interface definition of `Enumeration`, see appendix D.

No further attributes or routines can be defined for enumeration objects. The only characteristics of enumerations are: they exist, they are distinct, and they are ordered.

**Qualified Enumeration Constants**

In some situation it may be necessary to qualify enumeration values with their type. This situation arises when a class uses two enumeration types that define one or more common enumeration constants. Consider the following two enumeration classes:

```plaintext
class Colour is Enumeration
  == Simple enumeration type for basic colours
  manifest red, white, blue
end class

class Program is Enumeration
  == Simple enumeration type for basic colours
  manifest emacs, vi, blue
end class
```

If a third class now uses both `Colour` and `Program`, the constant `blue` is ambiguous. Trying to use it will result in a compile time error. This conflict can be resolved by preceding the enumeration value with its class name and an exclamation mark (!):

**Example:**

```plaintext
var col : Colour
  prog : Program

... 

col := Colour!blue
  prog := Program!blue
```

### 3.5 Operator Precedence

(to be written)
4. Identifiers

Identifiers are used as names for classes, variables, constants and routines. Identifiers are strings that consist only of letters, digits and the underscore (_), where the first character is a letter or an underscore.

Example of legal identifiers are:

- abc
- Num23
- addAtEnd
- add_at_beginning
- _count

Examples of illegal identifiers are:

- hit cnt \textit{-- error: space in identifier}
- 23add \textit{-- error: digit at start of identifier}
- #elem \textit{-- error: illegal symbol (#) in identifier}

Case of Characters

Blue identifiers are case sensitive. That means that \textit{abc} and \textit{ABC} are two different identifiers.

By convention, class names are often written with a leading capital, whereas routine names start with a lowercase character.

5. The Uses Clause

If a class A declares variables or parameters of another class B (we then say "A uses B"), and B is not one of the predefined classes, then A must declare the use of B.

BNF:

\[
\begin{align*}
\text{class-decl} & ::= \text{class identifier [ "<" generics-list ">"] is class-definition} \\
\text{class-definition} & ::= [ \text{identifier} ] \text{general-class-decl} \\
\text{general-class-decl} & ::= \text{class-comment} \\
& \quad\quad\quad\quad \text{uses [ ident-list ]} \\
& \quad\quad\quad\quad \text{...} \\
& \quad\quad\quad\quad \text{end class}
\end{align*}
\]

The following example shows a class \textit{Rectangle} that uses two classes \textit{Point} and \textit{Colour}. 

Example:

```pascal
class Rectangle is
  == ...
  uses Point, Colour
  ...
end class
```

The effect of the `uses` clause is that the classes `Point` and `Colour` are known inside the class definition of `Rectangle` and can then be used as types for variables, parameters and return values. The classes `Point` and `Colour` must exist in the current project.

Uses clauses are inherited. If a superclass lists a class in its uses clause, the subclass does not need to (and indeed must not) repeat the same class in its uses clause. A superclass itself is automatically known to the subclass – it does not need to be listed as a used class. For more details about inheritance see section 16.

The `uses` keyword itself is not optional. If a class does not use other classes the uses list is empty, but the keyword must appear in the source.

The predefined classes (Integer, Real, Boolean, String, Array) do not need to be listed in the uses clause. They are automatically known in every class without being explicitly mentioned as being used. This is, in fact, the meaning of the term `predefined` – those classes are always considered used by every class.

### 6. Variables

Variables can be declared of any class. This is done in separate variable declaration sections. Such a section exists for the whole class (instance variables) and once for each routine (local variables) and is preceded by the keyword `var`. The scope for instance variables is the class definition, the scope for local variables is the routine they are declared in.

The section declaring the instance variables must be in the `internal` section and must precede the internal routine declarations. No variables are allowed in the class interface.

The section declaring local variables follows the preconditions and constant declarations in a routine declaration and precedes the routine statement block.

**BNF:**

```plaintext
var-decls ::= var-decl { var-decl }  
var-decl ::= ident-list "::" type-def [ initialisation ]  
type-def ::= class-ident [ "<" ident-list "]" ]  
class-ident ::= identifier  
initialisation ::= ":=" expression
```
Examples:

```pascal
var
  num1, num2: Integer
  count : Integer := 0
  int_arr: Array <Integer>
  f: Figure
  s: Stack <Figure>

printList (l: List) is == print all list elements
  var
    cnt: Integer
    nm: String := ""
  do
    ...
  end printList
```

**Lifetime**

The lifetime of local variables is the time of a routine execution. In other words: scope and lifetime are connected (as in Pascal). There is no mechanism to extend the lifetime of local variables beyond the procedure execution (such as the "static" construct for local variables in C).

The lifetime of instance variables is the time of existence of the object they are part of.

**States and Values**

Variables are a combination of type, state and value. The type of a variable is specified statically and never changes. At any time, each variable is in one of two states: *undefined* or *defined*. If a variable is *defined*, it also has a value.

The value of a variable is always a reference to an object of its type (or *nil*, see below).

A variable can be initialised on declaration by any legal expression, using the assignment instruction (see example above). If a variable is not initialised at its declaration, its initial state is *undefined*. It is an error to use an undefined variable. An attempt to do so results in a runtime error.

No function exists to check for the state *undefined*. It is the responsibility of the programmer at the time of writing the program to ensure that variables are not used in the undefined state. (A good programmer will always know when a variable is potentially undefined. Parameters can never be undefined – that would have resulted in a runtime error at the time of the call.)

Assignment of a value to a variable changes the state to *defined*. After a variable leaves the state *undefined*, it never returns to that state.

**Nil**

A special value called *nil* exists. The value *nil* indicates that the variable does not reference any object. A variable of any type can hold the value *nil*. *nil* can be assigned, passed as a parameter and used in comparisons. Trying to use the object held by a nil variable (e.g. in an object call) is an error, since the variable does not refer to an object.

**Encapsulation**
Variables are encapsulated and cannot appear in the interface of a class. This means that a class cannot directly access a variable of another class. (It can indirectly access it if that class offers a function that returns or changes the value of a variable, but this is the choice of the class owning the variable.)

The only variables directly accessible are local variables of the current routine and instance variables of the current class.

Note, however, that encapsulation is class based, not object based. As a result, a class can access variables in another object, if that object is of the same class.

Example:

```pascal
class Point is
  ...
  var
    x, y: Integer
  ...
  add (other: Point) is
    == Add another point to this one
    do
      x := other.x
      y := other.y
    end add

This technique is useful in situations where access to the internals is needed, and no interface access functions are defined. (An object clone function is such a case where an object might want to create a new object of its own class and set all its instance variables to the same values as itself. We would not want to make all internal variables publicly accessible only to be able to provide a clone function.)
```

The possibility of accessing instance variables of another object, even though they are of the same class, seems to some people to go against the information hiding (or encapsulation) principle. And in some respects it does. To understand the reason for this access to be allowed, we must understand why information hiding is a good idea in the first place.

Information hiding is a software engineering technique, that decouples the implementation of different modules (here: classes) from each other. Accessing variables directly in arbitrary places is a maintenance disaster, because changes to the implementation of one module might require changes in unexpected other places in the program. This is why we do not allow direct access of instance variables.

If the access comes from the same class, though, (even if it is from another object) the argument does not hold anymore. If I change the implementation of that class, I will change all accesses to a variable within that class. So allowing access to instance variables of other objects from within the same class is no problem.
7. Constants

Constants in Blue can be literals or named constants. Literals are available for the predefined manifest classes. Named constants can be created from every legal expression.

7.1 Literals

Literals can be written for all predefined types. They have been introduced in the sections describing the individual types. Here, we just provide a brief summary of possible literal values for the predefined types. For details, refer to the section introducing each type (sections 3.3.1 to 3.3.5).

<table>
<thead>
<tr>
<th>Class</th>
<th>Integer</th>
<th>Real</th>
<th>Boolean</th>
<th>String</th>
</tr>
</thead>
<tbody>
<tr>
<td>Examples</td>
<td>17</td>
<td>2.0</td>
<td>true</td>
<td>&quot;hello&quot;</td>
</tr>
<tr>
<td></td>
<td>0</td>
<td>0.0</td>
<td>false</td>
<td>&quot;&quot;</td>
</tr>
<tr>
<td></td>
<td>-5</td>
<td>-0.122</td>
<td>TRUE</td>
<td>&quot;ab\n&quot;</td>
</tr>
</tbody>
</table>

7.2 Named Constants

Named constants are defined in a separate constant definition section, which is preceded by the reserved word const. The constant definition section can precede every variable section. Scope rules are the same as those for variables, and the syntax of the declaration is the same as for variables, except that the initialisation is not optional.

BNF:

```
const-decls ::= const-decl { const-decl }
const-decl   ::= ident-list "::" type-def initialisation
initialisation ::= ":=" expression
```

Examples:

```blue
const
  size: Integer := 99
  default_title: String := "untitled"

printList (l: List) is
  == print all list elements
  const
    empty: Boolean := l.isEmpty (false)
  var
    ...
  do
    ...
end printList
```
Note that literals as well as arbitrary function calls may be used for the initialisation of constants. Instance variables cannot be used for the initialisation of instance constants — they are still uninitialised. Care must be taken with function calls: If the function used for initialisation uses instance variables, a runtime error will be generated for using an uninitialised variable. The initialisation of instance constants and variables takes place on object creation before the creation routine is executed.

The exact order of object initialisation is:

- Initialisation of instance constants
- Initialisation of instance variables
- Execution of the creation routine, including (in this order)
  - Initialisation of local constants
  - Initialisation of local variables
- Execution of routine body

All constants and variables are initialised sequentially in the order in which they are defined. Later definitions may use preceding ones.

7.3 Set Constants

Blue allows sets of literals to be specified. These sets can only appear as literals. No variables of these set types can be created. The only available operator on these sets is \textit{in}, testing a value for membership in a given set. The \textit{in}-operation returns a boolean value. (See section 10.4 for more details on the \textit{in} operation.) Constant sets are also used for \textit{case} statements (section 11.2).

BNF:

\begin{verbatim}
set_expr ::= "{ [ set_elem { "," set_elem ] } "}
set_elem ::= expression | subrange
subrange ::= expression ".." expression
\end{verbatim}

Examples:

\begin{verbatim}
if choice in \{1, 2, 3 \} then ...
if result in \{0..9\} then ...
whitespace := ch in \{" ", "\n", "\t\} 
if ch in \{"a".."z", "A".."Z", "_\} then ...
\end{verbatim}

As the examples show, sets can be created by enumeration of values by comma separated lists, subranges, or a combination of both.

Subranges can only be taken from ordinal types. The ordinal types are \texttt{Integer}, \texttt{String} and \texttt{Enumeration}.

This construct does not replace a general purpose set type. It is merely a convenience notation that abbreviates relational comparisons. A general purpose set class can be found in the Blue Standard Collection Library.
There are two types of routines: procedures and functions. Both have an optional parameter list. Functions return one or more values.

All parameters are passed by value. (Note that variables always hold references, so a reference is passed by value, resulting in semantics of pass by reference for the objects themselves.)

Parameters cannot be used to return values directly. (There are no “var” parameters.) The only way to return a value is via a function result. Thus procedures do not return any values (but they can change the state of an object they reference). A function can return more than one result.

BNF:

```
routine-decl ::= \textit{identifier} \{ \textit{"(" parameter-list ")\text{"}} \} \{ \textit{"->"} \{ \textit{"(" parameter-list ")\text{"}} \} \} \textit{is routine-impl}

parameter-list ::= param-decl \{ \textit{",”} param-decl \}

param-decl ::= \textit{identifier} \?\textit{":"} \textit{type-def}

routine-impl ::= \textit{deferred routine-spec} \textit{|} \textit{builtin routine-spec} \{ \textit{redefined} \} routine-body

routine-body ::= \textit{routine-comment} \{ \textit{pre condition} \} \{ \textit{const const-decls} \} \{ \textit{var var-decls} \} do \textit{statement-list} \{ \textit{post condition} \} end \textit{identifier}

routine-spec ::= \textit{routine-comment} \{ \textit{pre condition} \} \{ \textit{post condition} \} end \textit{identifier}
```

Examples:

A short function is

```
count \rightarrow (val: Integer) is
\quad == Return number of elements
\quad do
\quad \quad val := i\_count
\quad end count
```

A simple procedure without parameters, pre or post conditions:
**Blue Specification, Version 1.0, 24.7.98**

```
show is
  == displays this element on standard output
  do
    print ("Name: ", name, "\n")
    print ("Count: ", count, "\n")
  end show

A function with two parameters and two return values:

`lookup (n: Integer, show: Boolean) → (nm: String, cnt: String)` is
  == returns name and count for entry number n
  == n: number of entry to lookup, must be valid
  == show: if true, value is also printed to standard out
  == nm: returns the name of the entry
  == cnt: returns the number of the entry

pre
  n >= 1
  and n <= nr_of_entries
  do
    nm := entries[n].name
    cnt := entries[n].count
    if show then
      entries[n].show
    end if
  end do
post
  nr_of_entries = old nr_of_entries
end lookup
```

The comment after the function header is part of the language. Comments for routines are not optional and are part of the routine interface.

Pre- and post conditions are optional.

Values are returned by assigning to the result variable (`nm` and `cnt` in this example). Result variables are initially undefined and must be assigned a value during function execution. It is an error to return from a function with undefined result variables. A function without an assignment statement to a result variable will produce a compile time error; a function that contains but does not execute such a statement results in a runtime error.

**Visibility**

Routines are visible in the whole class scope. This means that a routine call may precede the routine definition, allowing mutually recursive routines.

> Many people hold the view that functions should not have side effects. Suggestions have been made to us that the language should not allow side effects in functions (by inhibiting assignments to instance variables and routine calls to objects held in instance variables). While this is technically possible, it is impractical. We agree that it is a valuable design rule that functions should not have side effects, but side effects in this context can only mean changes to the external state of the object.

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Sometimes side effects that change the internal state of the object are desirable. An example is the implementation of a query object with caching. A query (a function) should not change the external state of the object, but it may well result in some information being cached inside the object (thus changing the state of the object – a side effect). If this change does not affect the externally visible behaviour (other than in efficiency), we say that the external state has not been changed. Our design rule is thus satisfied. The compiler cannot distinguish whether changes of state are internal or external. Thus this valuable design rule remains unenforced by the compiler.

For details on pre- and postconditions, see section 12.
For details on redefined routines, see section 16.2.
For details on deferred routines, see section 16.4.

8.1 Builtin Routines

Builtin routines are written with the keyword `builtin` after the routine header:

```
BNF:

routine-decl ::= identifier [ "(" parameter-list ")" ]
                [ "->" "(" parameter-list ")" ] is routine-impl

routine-impl  ::= builtin routine-spec
                | ...

routine-spec  ::= routine-comment
                [ pre condition ]
                [ post condition ]
end identifier
```

Example:

```
cursorTo (x: Integer, y: Integer) is builtin
  == Set the terminals cursor to screen position (x,y).
  pre
  x >= 0 and x <= width and
  y >= 0 and y <= height
end cursorTo
```

The builtin keyword is used only by system programmers writing classes that are part of the Blue environment. It cannot be used by application programmers. Builtin classes are usually implemented in a system-specific way (often in another language than Blue).
9. Statements

Statements are:
- assignment
- assignment attempt
- procedure call
- return
- assertion
- control structures

Control structures are described in section 11. This section only describes simple (non-compound) statements.

BNF:

```
statement ::= assignment
            | assignment-attempt
            | procedure-call
            | return
            | assertion
            | ...
```

9.1 Assignment

Values can be assigned to variables if the type of the value conforms to the type of the variable (see below for definition of conformance).

BNF:

```
assignment ::= ident-list "::=": expression-list
ident-list ::= identifier {"","identifier"}
expression-list ::= expression {"","expression"}
```

Examples:

- `a := 42`
- `a, b, c := x, y, z`  // multi-assignment (see below)
- `a, b, c := f()`

Multi-Assignment

The multi-assignment is defined as follows: Each of the expressions on the right hand side is evaluated to one or more values. (An identifier always evaluates to one value, a function evaluates to the list of its return values.) The order of evaluation is undefined. The values are then assigned to the identifiers on the left hand side. The number of identifiers must be equal to the number of values produced by the expressions on the right hand side. The first value is assigned to the first identifier, the second value to the second identifier, and so on. The evaluation of all expressions on the right hand side
is always completed before any assignment is executed. This allows a statement like

```
a, b := b, a
```

to swap two values.

**Conformance**

A type \( B \) conforms to a type \( A \) if

- \( B \) is the same type as \( A \)
- \( B \) is a (direct or indirect) subclass of \( A \)

Thus, the code segment

```
var
x: A
y: B -- B is subclass of A
...
y := create B
x := y
```

is legal, because \( B \) is a subclass of \( A \), and so \( B \) _conforms to_ \( A \).

The variable \( x \) is said to have the _static type_ \( A \) and the _dynamic type_ \( B \).

An assignment

```
a := b
```

is legal, if the static type of \( b \) conforms to the static type of \( a \).

### 9.2 Assignment Attempt

The assignment attempt allows an assignment from a variable \( a \) of some class to a variable \( b \) of a subclass. Such an assignment can be safely executed without risking type errors only if the _dynamic type_ of \( a \) conforms to the static type of \( b \). This can not be statically determined and causes a dynamic check to be executed.

**BNF:**

```
assignment-attempt ::= ident-list "?=" expression-list
ident-list ::= ident {"," ident }
expression-list ::= expression {"," expression }
```

**Example:**

```
var
x: A
y: B -- B is subclass of A
...
y ?= x
```

Statically, this assignment cannot be guaranteed to be successful (since the type of \( x \) does not conform to the type of \( y \)). \( x \) could reference an object of the dynamic type \( A \). In that case the assignment cannot be executed, because \( y \)
cannot hold references of type $A$. But $x$ could also reference an object of the
dynamic type $B$. In that case the assignment can be executed.

The assignment attempt results in a dynamic (runtime) check of the dynamic
type of $x$. If the dynamic type of $x$ conforms to the static type of $y$, the
assignment will be executed, otherwise $y$ will be assigned the value \textit{nil}.

An assignment attempt
\[ a := b \]
is statically legal if the static type of $b$ conforms to the static type of $a$, or the
static type of $a$ conforms to the static type of $b$ (in other words: if there is any
chance that the assignment could be successful).

9.3 Procedure Call

A procedure call is a call to a routine that does not return parameters.

\textbf{BNF:}

```
procedure-call ::= call-chain
call-chain ::= [ super"!" ] unqualified-call { "." unqualified-call }
unqualified-call ::= identifier [ "(" expression-list ")" ]
```

\textbf{Examples:}

```
proc
proc (a,b)
obj.proc (a, 23+17)
alist.get (22).put ("hello")
```

A procedure call can be a call of either an internal or external procedure. An
internal procedure is a procedure that is defined in the same class containing
the procedure call. An external procedure is a procedure defined in another
class. External procedure calls are preceded by an object identifier, separated
from the procedure name by a dot (\texttt{.}). If a procedure has parameters, the call
has a list of the actual parameters in parenthesis. If no parameter list exists,
no parenthesis are written.

9.4 Return

The return instruction causes a return from a routine.

\textbf{BNF:}

```
return ::= return
```

\textbf{Example:}

```
show is
  == displays this element on standard output
```
do
  if name = "" then
    return
  ...
end show

return does not take any parameters. If the routine is a function, all return values must be assigned before the return is executed.

Routine exits from anywhere in a routine (even in the middle of a loop!) are seen by some as “unstructured programming”. For a good summary of the reasons that show the advantages (in terms of programming and teaching) of this construct, see Roberts, E., Loop Exits & Structured Programming: Reopening the Debate, SIGCSE Bulletin, 27, 1, March 1995, pp. 268-272

9.5 Assertion

Assertions are used as a tool for correctness assurance and debugging.

BNF:

| assertion       | ::= assert "(" condition ")"
| condition       | ::= boolean-expression |

Example:

assert (name <> nil)

If the boolean expression in the assertion is not true, program execution is interrupted and the user is notified about failing of the assertion. If the expression is true, the statement has no effect.

10. Expressions

Expressions are:

- function call
- equality
- type equality
- in-expression (sets)
- create
- this

BNF:

expression ::= function-call |
             operator-expression
10.1 Function Call

A function call is a call to a routine which returns one or more values.

**BNF:**

```
function-call ::= call-chain
call-chain ::= unqualified-call { "." unqualified-call }
unqualified-call ::= identifier [ "(" expression-list ")" ]
```

**Examples:**

```
x := func
a,b := func (1)
if obj.func (a,b+c) then ...
found, element := list.search (22);
if found then
  print (element)
end if
```

A function call is syntactically similar to a procedure call. Since a function returns values, it is used in an expression rather than as a statement. A function can return one or more values, which can then be assigned to variables in a (multi-)assignment.

If a function returns exactly one value, it can also be used in expressions such as the condition of an `if` statement or in an actual parameter list. Functions returning multiple values can only be used in a multi-assignment statement.

If a function has no parameters, a call to that function consists of its name only (no parentheses are written to indicate an empty parameter list).

10.2 Equality

**BNF:**

```
reference-equality ::= expression comparison expression
comparison ::= "=" | "<>"
```
Examples:

    a = b
    found = true
    a <> 42

See section 3.2.2 (Comparison).

10.3 Type Equality

BNF:

    type-equal ::= expression is type

Examples:

    var
        f: Figure
        r: Rect -- inherits from figure
    ...
    if f is Rect then
        r := f
    end if

The keyword is checks whether a variable is of a given type. The result is a boolean value. The expression

    ident is Mytype

is true, if the dynamic type of ident is Mytype or a subclass of Mytype. In other words: The expression is true, if the dynamic type of ident conforms to Mytype.

10.4 In

BNF:

    in ::= expression in set_expr
    set_expr ::= "{ " set_elem { "", set_elem } }"
    set_elem ::= expression subrange
    subrange ::= expression ".." expression

Example:

    if num in {1, 3, 5, 7} then ...
    if n in {1..10} then ...
    valid := x in {" ", \
    \n    "", "\t", "0"..'9" }
In expressions are aliases and are defined by combinations of comparisons:

\[
\begin{align*}
n & \in \{a, b\} \quad \text{is defined as} \quad n = a \text{ or } n = b \\
n & \in \{a..b\} \quad \text{is defined as} \quad n \geq a \text{ and } n \leq b \\
n & \in \{a, b..c\} \quad \text{is defined as} \quad n = a \text{ or } (n \geq b \text{ and } n \leq c)
\end{align*}
\]

### 10.5 Create

**BNF:**

\[
\begin{align*}
\text{create} & : \text{=} \text{ general-create | array-create} \\
\text{general-create} & : \text{=} \text{ create class-ident [ "(" expression-list ")" ]} \\
\text{array-create} & : \text{=} \text{ [" expression } , \text{ expression ]}
\end{align*}
\]

**Example:**

```plaintext
var
  r : Rectangle
  b : Buffer
...
  r := create Rectangle (p1, p2, c)
  b := create Buffer
```

The `create` keyword creates an object of the specified class and executes its creation routine. Parameters to the creation routine are passed in a parameter list after the class name. If the creation routine has no parameters, the create call has no parameter list. (If the class does not specify a creation routine, a default routine without parameters is generated).

### 10.6 This

**BNF:**

\[
\begin{align*}
\text{expression} & : \text{=} \text{ this}
\end{align*}
\]

**Example:**

```
list.add (this)
```

The expression `this` is a reference to the currently active object. It can be used to pass the current object reference to other objects.
11. Control Structures

The remaining group of statements, not described so far, are control structures. Control structures are:

- conditional
- multi branch
- loop

**BNF:**

\[
\text{statement} ::= \ldots | \text{conditional} | \text{selection} | \text{loop}
\]

### 11.1 Conditional: The If Statement

**BNF:**

\[
\text{conditional} ::= \text{if} \ \text{condition} \ \text{then} \\
\text{statement-list} \\
\{ \ \text{elseif} \ \text{expression} \ \text{then} \\
\text{statement-list} \} \\
[ \ \text{else} \\
\text{statement-list} ] \\
\text{end if}
\]

\[
\text{condition} ::= \text{boolean-expression}
\]

\[
\text{statement-list} ::= \{ \text{statement} \}
\]

**Examples:**

```plaintext
if val>0 then
    handle_positiv
elseif val=0 then
    handle_zero
else
    print(“The value was negative!”)
    handle_error
end if
if n < 0 then
    n := –n
end if
```

The if statement is similar to conditionals in many other languages. Note that the if and else parts contain statement lists. This makes an explicit grouping.
symbol at the beginning of the statement list unnecessary. The “end if” is always required.

11.2 Selection: The Case Statement

The case statement allows a multi-way branch depending on an expression. It is similar to the case statement in Pascal or the switch statement in C. The case labels are set constants and selection of a case is done by applying an in operation to the expression and the case labels. The expression is evaluated and the resulting value is checked for membership in the label sets. The statements at the first matching label are executed. Only the statements between the label and the next label (if any) are executed. Then execution resumes after the end case instruction. The else part is executed if the value was not a member of any of the label sets. If no set contains the value and the case statement does not contain an else part, no nested instructions are executed and execution continues after the case statement.

BNF:

\[
\text{selection} ::= \text{case expression of} \\
\text{ \{ set_expr ":" statement-list \} \\
\text{[ else statement-list ]} \\
\text{end case}
\]

Examples:

case colour of
  {red}:
    label := "danger"
  {yellow}:
    label := "warning"
  {green}:
    label := "okay"
else
  handle_error
end case

case value of
  {0..33}:
    print ("low")
  {34..66}:
    print ("medium")
  {67..99}:
    print ("high")
  {100}:
    print ("top")
end case
case ch of
  {" ", "\n", "\t"}:
    handle_whitespace
  {"a"..'z', "A"..'Z', "_"}:
    handle_letter
  {"0"..'9"}:
    handle_number
else
  handle_error
end case

If the sets used as case labels overlap and the test value is a member of more
than one of the sets, only the instructions at the first match are executed.

11.3 Iteration: The Loop Statement

There is only one kind of loop structure in Blue. The loop is defined with the
keywords `loop ... end loop`.

Exit from the loop is explicit by using the `exit on` keyword, which is part of
the loop. Every loop must have at least one exit. Multiple exits are possible.
By placing the exit at the beginning or end of the loop, behaviour of a pre-test
or posttest loop (while and repeat in Pascal) can be implemented.

\[
iteration ::= \text{loop} \quad \text{statement-list} \\
\quad \text{exit on} \quad \text{condition} \quad \text{statement-list} \\
\quad \{ \quad \text{exit on} \quad \text{condition} \quad \text{statement-list} \quad \} \\
\quad \text{end loop}
\]

Note that a statement list can be empty.

Examples:

The first example shows a loop with the exit instruction at the beginning.
This construct has semantics similar to a while loop (where the condition in
the while loop would be the negation of the exit condition in this example).

```
list.first
loop
  exit on list.atEnd
  print (list.current)
  list.advance
end loop
```

Similarly, a post-test loop (`repeat` in Pascal) can be constructed by placing the
exit instruction at the end of the loop.

In Blue, however, exit instructions can appear at any stage in the loop,
increasing flexibility:
loop
  readInt (i)
  exit on i>0
  print (“The number has to be greater than 0. Enter again.”)
end loop

The last example shows the use of multiple exit instructions.

loop
  s := getSelection
  exit on s=0
  error := process (s)
  exit on error
end loop

A loop without an exit statement causes a compile time error. Exit statements are illegal outside a loop. In nested loops, exit always exits only the innermost loop. It is not possible to exit an outer loop from within a nested loop.

12. Pre/Post Conditions

Preconditions and postconditions are part of routine definitions. They are optional. If present, they are automatically checked at runtime and an error is reported if a condition is not met.

Preconditions are written at the beginning of a routine body (before the variable declarations) and are checked on entry of that routine. It is the responsibility of the caller to ensure that the precondition is met (i.e. that it evaluates to true on routine entry). The code inside the routine body can then safely assume that the condition is true.

Postconditions are written at the end of a routine body and are checked on exit of the routine. It is the responsibility of the routine to ensure that the postcondition is true, and the caller can assume that the condition is met on return from the routine.

BNF:

```
routine-body ::= routine-comment
               [ pre [ condition ]
               [ comment ]]
               ...
               do
               ...
               [ post [ condition ]
               [ comment ]]
               end identifier
```

Pre- and postconditions consist of one condition and/or a comment each. If multiple conditions are to be defined, they can be combined with a logical and to form one condition. The comment is intended to express conditions that cannot be expressed in Blue expressions. The comment has no runtime effect.
(no error is generated if a comment is false), but it is good programming practice to include documentation about pre- and postconditions even if they cannot be written in code.

Examples:

```plaintext
func (n: Integer, m: Integer) -> (res: Integer) is
pre
    n >= 0 and
    m >= 1
var
    x : Integer
    do
        ...
    post
        res <> nil
end func

printName is
    do
        ...
    post
        == the name has been printed on screen
end printName
```

Both pre- and postconditions are part of the interface of a routine. Redefinitions of routines may alter these conditions only in restricted ways: preconditions may be weakened, postconditions may be strengthened. If a redefined routine does not define pre- or postconditions, the conditions of the original (parent) routine apply. If a precondition is redefined, the actual condition tested at runtime is

```
precondition or parent-precondition
```

If a postcondition is redefined, the condition tested at runtime is

```
postcondition and parent-postcondition
```

This ensures that the redefined routine guarantees at least as much as its parent.

Two special expression are available in conditions: \( \Rightarrow \) (implies) and the reserved word \texttt{old}.

Example:

```plaintext
post
    found => index > 0
```

The implies symbol can be read as

```
if (found) then assert (index > 0)
```
The old expression is available only in postconditions. It is used on exit of a routine to refer to the value that an expression had on entry to that routine.

The following example ensures that the global variable num has not been altered in the routine:

Example:

\[\text{doSomething (n: Integer) is} \]
\[\quad \text{pre} \]
\[\quad \quad \text{n} \neq \text{nil} \]
\[\quad \text{post} \]
\[\quad \quad \text{num} = \text{old num} \]
\[\text{end doSomething} \]

It is an error to use “\(=\)“ outside of a pre- or postcondition or class invariant or to use “old” outside a postcondition.

13. Class Invariants

Class invariants are conditions that have to be met by any stable state of an object. Stable states exist before and after every execution of an interface routine from the outside of an object. (Note that when an interface routine is called locally from within an object, it does not have to be in a stable state.)

BNF:

\[
\begin{align*}
\text{class-decl} & \quad ::= \text{class identifier [ "<" generics-list "]} \text{ is} \\
\quad & \quad \text{class-definition} \\
\text{class-definition} & \quad ::= \ldots \\
\quad & \quad [ \text{identifier} ] \text{ general-class-decl} \\
\text{general-class-decl} & \quad ::= \ldots \\
\quad & \quad \text{[ invariant condition ]} \\
\quad & \quad \text{end class} \\
\end{align*}
\]

Invariants are checked at runtime before and after every external interface routine call and after execution of the creation routine. A runtime error is generated if an invariant evaluates to false.

14. Comments

Blue recognises two types of comments, marked by a double equals sign (==) and a double hyphen (––). Comments always extend to the end of the line. To have several lines as a comment, every line has to be preceded by the comment symbol.
14.1 Interface Comments

Comments starting with == are part of the class interface and part of the language definition. They are used to describe the class itself and routine semantics, and are allowed to appear only in strictly defined locations. These locations are:

• after the class header
• after a routine header

Example

```pascal
class Rectangle is
  == Author: M. Kölling
  == Date: April 1995
  == Version: 1.1
  == Short: Graphical representation of a rectangle.
  ==
  ==Class Rectangle represents a rectangle with
  ==specified coordinates and colour that can be ...
  ==...
  ...

  interface

  creation (tl: Point, br: Point) is
    == Create rectangle at coordinates defined by tl (top-left)
    == and br (bottom-right). Default colours are: fill white,
    == border black.
    do
      ...
    end creation

  routines

  move (dx: Integer, dy: Integer) is
    == Move rectangle by distance defined by dx, dy.
    do
      ...
    end move

  size -> (width: Integer, height: Integer) is
    == Return size of rectangle in width and height.
    do
      ...
    end size

  invariant
    ...

end class
```
The example shows interface comments for the class and for each routine. Interface comments may not appear anywhere else. They are part of the class interface and therefore displayed in interface view (see [1] for a description of the environment, including the interface view of classes). They are also used by the class browser. (By convention, some lines of the class comments begin with certain keywords which are recognised by the class browser. These lines define the author, version number, date of creation and a short description of the class. See the documentation of the browser for details.)

14.2 Implementation Comments

All other comments, starting with the symbol – – (double hyphen), are implementation comments. They may appear anywhere and are not included in the interface view. Implementation comments are ignored by the compiler.

Note that implementation comments may appear as part of the routine comment. If they do, they should describe the routine implementation and will not be displayed as part of the interface.

Example:

```blue
class WidgetManager is
  ...
  add (new_widget: Widget) is
    == Add ‘new_widget’ to the set of managed widgets.
    -- -- This is done by adding ‘new_widget’ to ‘wl’, the
    -- -- internal widget list. Note that ‘wl’ may be undefined
    -- -- if this is the first widget to be added. We have to test for
    -- -- that and possibly create the list before adding to it.
    do
      if wl = nil then -- -- test whether list exists
        wl := create List <Widget>
      end if
      wl.add (new_widget) -- -- add widget to widget list
    end if
  end add
  ...
end class
```

The interface of this routine is

```blue
add (new_widget: Widget)
  == Add ‘new_widget’ to the set of managed widgets.
```

There is no block comment in Blue (every comment is for one line only). Instead, the Blue environment provides tools to add or remove comments on every line in a block of text.
I/O is implemented in Blue by providing a number of standard classes and objects that can be called to perform input/output operations. There are two groups of I/O classes: one group for text based (or "standard") I/O and one group for more sophisticated graphical user interface building. The graphical user interface classes are found in the Blue GUI library and are not described here. The following sections describe what is known as standard I/O, a simple, text based I/O facility.

15.1 Standard I/O

Blue provides four environment variables for standard I/O: input, output, terminal and format. These variables are predefined and accessible in all Blue classes.

It also provides two predefined objects: one object of class TextTerminal and one object of class OutputFormat. The predefined variables input, output and terminal typically all refer to the predefined TextTerminal object, the format variable refers to the OutputFormat object. It is possible, though, to bind input or output to other objects, for instance output to an object representing a printer or input to a file. This can be done externally before a routine is executed, or dynamically by the program.

The predefined variables are of the following classes:

```plaintext
var
    input : IO_Channel
    output : IO_Channel
    terminal : TextTerminal
    format : OutputFormat
```

IO_Channel is a superclass of TextTerminal, so the variables input and output provide a view onto the predefined TextTerminal object as an IO_Channel object.

The variables input and output exist to provide a means do redirect input or output to and from other channels (such as a text file, which is also a subclass if IO_Channel).

The variable terminal exists so that a program can take advantage of TextTerminal specific operations that are not part of IO_Channel (such as positioning the cursor on screen, etc.)

For a full description and interface, see the "Blue Standard Library Manual" [2].

15.1.1 Standard Output

Output to the terminal is usually done using the alias print.

Examples:

```plaintext
i : Integer
m : Myclass
```
print (i)
print (m)
print ("The value of ", i, " is ", m, "\n")

The alias \textit{print} is defined as
\begin{verbatim}
print (a, b, ...)  <=>  output.write (str (a, b, ...) )
\end{verbatim}
\textit{str} itself is an alias, which is defined as
\begin{verbatim}
str (a, b, ...)  <=>  a.toString.concat (b.toString.concat (...))
\end{verbatim}
The function \textit{concat} is defined in class \textit{String} and returns a string that is the concatenation of two other strings. The function \textit{toString} is defined for all predefined classes (except \texttt{Arrays}) and may be defined for user defined classes. It returns a printable representation of an object in a string. If a user defined class wants to define a \textit{toString} routine, it has to conform to the following signature:
\begin{verbatim}
toString  ->  (s: String)
\end{verbatim}
\[== Return a string representation of this object.\]

Every object that defines a function \textit{toString} conforming to this specification, can be printed using the alias \textit{print}.

Furthermore, apart from being used for output, the \textit{str} alias can be used for conversion to string from another type and for concatenation of string representations of any type, including strings.

\textbf{Examples:}
\begin{verbatim}
s, s1, s2 : String
i : Integer
m : Myclass
s := str (i)  -- conversion Integer to String
s := str ("The value is ", i)  -- conversion and concatenation
s := str (s1, s2)  -- concatenation of Strings
\end{verbatim}

\section{Output Formatting}

The \textit{toString} routines of the predefined classes \texttt{Integer}, \texttt{Real}, \texttt{Boolean} and \texttt{Enumeration} use the predefined \textit{format} object to determine their output format. \texttt{(format} is a predefined constant referring to a standard object of class \texttt{OutputFormat}.) Calls to \texttt{format} may be made to influence the formatting of number, Boolean and enumeration output. The (incomplete) interface of \texttt{OutputFormat} is as follows:

\begin{verbatim}
class interface OutputFormat is
setWidth (n : Integer)
== Set field width for output. Default is 0. If the string
== representation of a value consists of more characters than
== the field width, the width is ignored.
==
== Used by Integer, Real, Boolean and Enumeration.
==
== Note that the class \texttt{String} does not use the \texttt{format} object for
== output. String output is unaffected by the field width. To
\end{verbatim}
alignRight (right : Boolean)

== Specify alignment of output in the field defined by 'setWidth'.
== The default is 'left' (spaces will be added to the right of the
== output as appropriate). If 'right' is set to true, output will
== be right-aligned within the specified field width.
==
== Used by Integer, Real, Boolean and Enumeration.

pre
  n >= 0

right <> nil

scientific (useScientific : Boolean)

== If useScientific is true, scientific notation is used for output
== of real numbers (e.g. 1.23456e+02). Otherwise fixed point
== notation is used (e.g. 123.456). The default is
== useScientific=false.
==
== Used by Real.

pre
  useScientific <> nil

roundTo (n : Integer)

== Round output of fixed point real numbers to n digits after the
== decimal point. Default is 6. Used by Real.

pre
  n >= 0

The format object may also be used by user defined classes to determine the
behaviour of the toString routine (so that user defined output can, for
instance, honour field width or justification). For a full interface of
OutputFormat, see Appendix D8.

15.1.3 Standard Input

Input from standard input is defined for integers, reals, booleans and strings.
It is done using a group of aliases which map to routine calls to the standard
object input. The aliases are:

readInt  <=>  input.readInt
readReal  <=>  input.readReal
readStr   <=>  input.readString
readChar  <=>  input.readChar

All of these routines are functions returning objects of the according type.

Examples:

num := readInt
name := readStr
line [i] := readChar
readInt and readReal skip whitespace and end-of-line characters. If the next characters cannot be interpreted in the required format, they return nil. readStr reads until an end-of-line character is found, or the end of file is reached. The relevant part of the interface of the class IO_Channel is shown below:

class interface IO_Channel is
  . . .

  readStr -> (s : String) is deferred
    == Read a String. Reads characters until
    == a line break ("\n") is read or an error condition (such as
    == "end-of-file" is encountered). The resulting String does
    == not include the line break character.

  readChar -> (s : String) is deferred
    == Read the next character. Every character (including line
    == break characters) are returned as entered. The line break
    == character is returned as "\n".

  readInt -> (i : Integer) is deferred
    == Read an Integer. Skips white space (spaces, tabs and
    == newlines) before the Integer. Returns "nil" if the
    == next non-white characters do not represent a number.
    == Numbers are written with an optional - or + sign and digits.

  readReal -> (r : Real) is deferred
    == Read an Real (floating point number). Skips white space
    == (spaces, tabs and newlines) before the number. Returns "nil"
    == if the next non-white characters cannot be interpreted as a
    == number. Real numbers are written with an optional - or +
    == sign, digits and a decimal point (.).

  endLine is deferred
    == Read and discard all characters up to (and including) the next
    == NewLine character.

  atEnd  -> (isAtEnd : Boolean) is deferred
    == Return true if the current position is the end of the channel.
    == Some channels might not have an end (and atEnd always
    == returns false) some might reach an end at a specific point (like
    == a file, where the end of the channel might be defined by the
    == end of the file). In some channels the end might be
    == dynamically defined (i.e. when a terminal allows the user to
    == generate and at-end condition via a control key). See the
    == concrete subclasses for their individual handling of the at-end
    == condition.

end class
15.1.4 The TextTerminal Class

As mentioned above, the variables *input*, *output* and *terminal* by default refer to an object of class TextTerminal, which is created by the Blue environment on system startup (a project does not need to create this object explicitly). The TextTerminal class defines additional routines apart from those defined in IO_Channel. Those routines can be called only through the variable *terminal*, since it is the only one to be declared of class TextTerminal. There are no aliases for these routines. They are called using the standard routine call syntax. For example, to position the cursor on the screen, a class might include the statement

```plaintext
terminal.cursorTo (4, ln)
```

See Appendix D7 for the complete interface of class TextTerminal.

15.2 File I/O

Working with files is supported by standard classes, available in the standard library (group Standard_IO). These are the classes FileSystemHandle and TextFileHandle. Objects of both of these classes are handles – they give access to an underlying object. This indicates that, for instance, creating a file system handle does not create a file system, but rather creates an object that gives the user access to the underlying filesystem which exists independently from the handle object. Equally, creating a TextFileHandle does not create a text file – it just allows access to a text file that was created earlier.

**Class FileSysHandle**

The class FileSysHandle is used to access the file system. Access to the file system allows operations such as creating or removing files, reading directories, checking for the existence of files, etc.

The (incomplete) definition of the class FileSysHandle is

```plaintext
class interface FileSysHandle is
  ...

  creation (fileSys : String)
    == Create a handle to a file system. The parameter "fileSys"
    == identifies the file system to be accessed. On systems that
    == have only one file system (e.g. Unix), the parameter is
    == ignored.

  routines

  createFile (name : String) -> (done : Boolean)
    == Create a file with file name "name". If the file exists, its
    == length is truncated to 0.
    == ...

  deleteFile (name : String) -> (done : Boolean)
    == Delete the file with file name "name".
    == ...
```


copyFile (name : String, newName : String) -> (done : Boolean)
  == Copy the file with name "name" to a second file with name
  == "newName".
  == ...

createDirectory (name : String) -> (done : Boolean)
  == Create a directory with name "name".
  == ...

deleteDirectory (name : String) -> (done : Boolean)
  == Delete the directory with name "name". The directory must
  == exist and must be empty.
  == ...

readDirectory (name : String) -> (dir : Array <String>)
  == Read the directory with name "name". Returns an array with
  == the directory entries. Each entry is the name of a file or
  == directory. ...

exists (name : String) -> (result : Boolean)
  == Return result=true if an entry with name "name" exists in
  == the file system.
  ...

copyFile (name : String, newName : String) -> (done : Boolean)
  == Copy the file with name "name" to a second file with name
  == "newName".
  == ...

cREATEDIRECTORY (name : String) -> (done : Boolean)
  == Create a directory with name "name".
  == ...

deleteDirectory (name : String) -> (done : Boolean)
  == Delete the directory with name "name". The directory must
  == exist and must be empty.
  == ...

readDirectory (name : String) -> (dir : Array <String>)
  == Read the directory with name "name". Returns an array with
  == the directory entries. Each entry is the name of a file or
  == directory. ...

exists (name : String) -> (result : Boolean)
  == Return result=true if an entry with name "name" exists in
  == the file system.
  ...

copy (source: String, target: String)
  == Copy the file "source" to file "target" character by character.
  var
    file : TextFileHandle
    filesys : FileSysHandle
    file1 : TextFileHandle
    file2 : TextFileHandle
  do
    filesys := create FileSysHandle (nil)
    -- create the target file
    if (not filesys.createFile (target)) then
      print ("Could not create target file
      return

This interface is incomplete. Look at the interface online in the Blue class
browser, or have a look at "The Blue Libraries" [2] to see the full interface.

Class TextFileHandle
A file handle is used to access the contents of a file. A file handle is not the
file itself. This is an important distinction. Creating a file handle does not
create a file. Creation of a file handle rather corresponds to opening a file.

Example:
This example creates a copy of an text file.

copy (source: String, target: String) is
  == Copy the file "source" to file "target" character by character.
  var
    filesys : FileSysHandle
    file1 : TextFileHandle
    file2 : TextFileHandle
  do
    filesys := create FileSysHandle (nil)
    -- create the target file
    if (not filesys.createFile (target)) then
      print ("Could not create target file!
      return

44
end if

– – open both files

file1 := create TextFileHandle (source)
file2 := create TextFileHandle (target)

if (not file1.isReadable) or (not file2.isWritable) then
    print ("Could not open files with necessary access.\n")
    return
end if

– – do the copying

loop
    exit on file1.endOfFile
    file2.write (file1.readChar)
end loop

– – close the files

file1.close
file2.close
print ("Done.\n")

end copy

After creating a TextFileHandle, the state of the file handle should be checked. A file is always opened for reading and writing, if the underlying file system allows this access. If access rights in the file system deny read or write access, the file is opened in read-only or write-only mode. The access can be checked with the \texttt{isReadable} and \texttt{isWritable} functions. If no access is possible at all (either no access was granted by the file system, or the file does not exist) then \texttt{isReadable} and \texttt{isWritable} are both false, and \texttt{isBad} is true. All attempts to read from a file that is not readable, or to write to a file that is not writable result in runtime errors.

The function \texttt{endOfFile} can be used to check whether the current file position is at the end of the file. It must be checked before a read operation is attempted. Trying to read from a file while \texttt{endOfFile} is true is an error.

The (incomplete) definition of class \texttt{TextFileHandle} is

\begin{verbatim}
class interface TextFileHandle is IO_Channel
    creation (fileName : String)
        == Create a handle to a text file. The file "fileName" must exists
        == in the file system.
        ==
        == If a file with the specified name does not exist or if access to
        == the file is denied, no file will be opened and the status of the
        == TextFileHandle will be set to "isBad" (see routine "isBad"
        == below).
        == This routine does not create a file. To create a new file, use
        == the createFile routine in class FileSystemHandle.
    pre
\end{verbatim}
fileName <> nil

routines

inherited from IO_Channel

write (s : String) is redefined
== Write 's' to the file.
pre
  isWritable

readStr -> (s : String) is redefined
== Read a String from the file. Reads characters until
== a line break ("\n") is read or the end of the file is reached.
== The resulting string does not include the line break character.
== ...
pre
  isReadable and not atEnd

readChar -> (s : String) is redefined
== Read the next character from the file. Every character
== (including line break characters) are returned as entered
== The line break characte is returned as "\n".
pre
  isReadable and not atEnd

readInt -> (i : Integer) is redefined
== Read an Integer from the file. Skips white space (spaces,
== tabs and newlines) before the Integer. Returns "nil" if the
== next non-white characters do not represent a number.
== Numbers are recognised with an optional - or + sign and
== digits.
pre
  isReadable and not atEnd

atEnd -> (isAtEnd: Boolean)
== Return true if the end of input has been reached. While input
== from a terminal does not have a natural end, end end-of-input
== condition can be generated by pressing CTRL-D.
== If the next character entered is a CTRL-D, 'atEnd' removes
== the character from the input and returns true. If it is any other
== character, 'atEnd' leaves it pending in the input queue and
== returns false.
pre
  not isBad

...
pre
not isBad

setPosition (pos : Integer)
== Set the read/write position in the file to "pos".
== 0 is the beginning of the file, -1 is the end of the file.
== Positive numbers count from the beginning of the file,
== negative numbers count backwards from the end of the file.
pre
not isBad and pos <> nil

close
== Close the file associated with this handle. In general all files
== should be closed by the user as soon as they are not needed
== anymore. Closing files is important, because the number of
== files a user can have open at the same time is limited. If files
== are not closed, opening of further files might fail. The exact
== number of files that can be open at any one time is system
== dependent.
== Also, output to files is buffered. Text written to the file is
== not immediately written to disk. Only after closing the file or
== calling an explicit "flush" can the user rely on the file content
== being visible to other processes.
==
== Files not being closed are automatically closed when the file
== handle object is garbage collected. Since the user has no
== influence on the timing of garbage collection, it is bad practice
== to rely on this.
==
== After closing, the file handle is not usable any more, all read
== and write operations cause runtime errors, and "isBad" returns
== true.

isReadable -> (readable : Boolean)
== Returns true if the file can be read. If access rights in the
== file system specify write-only access, isReadable returns
== false.

isWritable -> (writable : Boolean)
== Returns true if the file can be written. If access rights in the
== file system specify read-only access, isWritable returns false.

isBad -> (bad : Boolean)
== "isBad" returns true if this file handle cannot be used to access
== a file. This can happen if the file specified in the creation
== routine does not exist, or if neither read nor write access was
== permitted. If "isBad" is true, "isReadable" and "isWritable"
== are both false.

end class
This interface is incomplete. Look at the interface online in the Blue class browser, or have a look at "The Blue Libraries" [2] to see the full interface.

16. Inheritance

Inheritance can be used to define specialisations of previously defined classes. These specialisations are called subclasses (or “children”) of their superclasses (or “parents”).

16.1 Defining Subclasses

Classes can be subclasses of at most one other class. The superclass (if any) is listed in the head of the class definition.

BNF:

\[
\begin{align*}
class-decl & ::= \text{class identifier } [ "<" \text{generics-list } ">" ] \text{ is class-definition} \\
class-definition & ::= [ \text{identifier } ] \text{general-class-decl} \\
\end{align*}
\]

Example

\[
\begin{align*}
class \text{ Car is Vehicle} \\
& == \text{Defines a car for a traffic simulation.} \\
& \ldots \\
end \text{ class}
\end{align*}
\]

A subclass inherits all routines and variables from its parent. All interface routines of the parent are also interface routines of the subclass, and all internals (routines and variables) are available inside the subclass. The creation routine is the only routine that is not inherited by the subclass.

Inherited interface routines cannot be "hidden" (removed from the interface). Inheritance is intended to be used to express real “is-a” relationships. If such a relationship is true, there is no need to hide interface routines of the parent. In all situations where a programmer wants to hide parts of an interface in a child, inheritance should not be used, but a client-server model (“uses” relationship) should be used instead.
16.2 Redefinition

Children may redefine the implementation of an inherited routine. This is done by providing a routine implementation in the child with the keyword redefined in the routine header.

BNF:

```
routine-decl ::= identifier [ "(" parameter-list ")" ]
                 [ "->" "(" parameter-list ")" ] is
                 routine-impl

routine-impl ::= [ redefined ] routine-body
```

Example:

```
move (dest: Location) is redefined
  == Move the position of this ...
  do
    ...
  end move
```

The explicit keyword redefined protects a programmer from accidental reuse of a routine name of the parent. The parameter list in the redefined function must be the same as the one of the original function. Changing of parameter lists (covariance or contravariance) is illegal.

Only the implementation of routines may be redefined. It is not allowed to redefine the parameter or result list of an existing routine.

16.3 Calling Superclass Functions

Sometimes is is desirable to call the original of a function that has been redefined (usually as part of the redefinition of the function). Consider the following example:

Example:

```
class Person is

  ...

  print is
    == Write this person’s details to the screen
    do
      print (name, “\n”)
      print ("date of birth: ", date, "\n")
    end print

end class

class Student is Person
```
\begin{verbatim}
print is redefined
  == Write this student’s details to the screen
  do
    super!print -- write Person details
    print ("student number: ", stud_num, ";n")
  end print
end class
\end{verbatim}

The class \textit{Person} has an interface routine \textit{print} that prints out that person’s details. The class \textit{Student}, which inherits from \textit{Person}, redefines print to print out the details for the student, which include the details for a person. \textit{Student}'s print routine can print the person details by calling the original print routine in its superclass. This is done by preceding the routine call with the keyword \textbf{super} and an exclamation mark (!).

\textbf{BNF:}

\begin{verbatim}
procedure-call ::= call-chain
call-chain ::= [ super! ] unqualified-call { "." unqualified-call }
unqualified-call ::= identifier [ "." expression-list ]
\end{verbatim}

This construct is often used to call the superclass’s creation routine from a child’s creation routine.

\textbf{Example:}

class Student is Person

\begin{verbatim}
  ... interface
  creation (name: String, dob: Date, sid: Integer) is
    == Create student instance
    do
      super!creation (name, dob)
      student_num := sid
    end creation
  ...
end class
\end{verbatim}

\textit{I have often been asked why we use an exclamation mark rather than the usual dot notation in this situation. In other words: why do we write super!creation instead of super.creation}

The answer is that, although similar on first glance, these are two very different instructions. The dot notation specifies an object before the dot. The routine call is performed on that other object. The exclamation mark
(or "bang notation") specifies the execution of a routine that was defined in another class on this object. The super keyword is a scope specifier, not an object specifier. We regard this distinction as so important that we want the programmer to be aware of the difference. The different syntax is an attempt to emphasise this difference.

16.4 Deferred Routines

Routine implementations can be deferred. Deferred routines are marked with the keyword deferred.

BNF:

```
routine-decl ::= identifier ["(" parameter-list ")"] ["->" "(" parameter-list ")"] is routine-impl

routine-impl ::= deferred routine-spec | ...

routine-spec ::= routine-comment [ pre condition ] [ post condition ] end identifier
```

Example:

```
move (dest: Location) is deferred
  == Move the position of this ...
  pre
  dest <> nil
  end move
```

If a routine is deferred, the class does not provide an implementation of the routine. Classes that contain one or more deferred routines are called abstract classes. No instances can be created of abstract classes. Children of abstract classes can provide implementations for deferred routines. If they do not provide implementations for all deferred routines, they remain abstract themselves.

Implementations for deferred routines are provided by redefining the routine in the child. See section 16.2 for a description of redefinition.

17. Genericty

Classes may be generic. Genericty (often called parametric polymorphism) allows classes to include unspecified types in their class description, which are instantiated when an object of this class is created. This allows a more
general form of code reuse. A list, for instance, can be defined to hold elements of some type, where the type is not specified at the time the class is implemented. The actual type of the list elements is then specified when an object of class list is created. Several instances of the list can exist, one holding elements of type Integer, another one elements of type String, and so on.

17.1 Unconstrained Genericity

Most of the BNF rules involved in genericity have been mentioned above. Here we repeat the relevant parts:

**BNF:**

```
class-decl ::= class identifier [ "<" generics_list "]"] is class-definition

generics-list ::= generics-decl { "," generics-decl }
generics-decl ::= form-actual-generic-param [ is class-type ]
formal-generic-param ::= identifier

var-decl ::= ident-list ":" type [ initialisation ]
type ::= class-type
    | formal-generic-param
    | ...

class-type ::= identifier [ "<" actual-generics-list "]"]
actual-generics-list ::= type { "," type }
```

The following example shows part of a class definition for a generic stack and code fragments declaring and using the stack class, instantiating it once as an Integer Stack and once as a Stack of Figures.

**Example:**

```markdown
class Stack <ELEM_TYPE> is
  internal
  var
    st: Array <ELEM_TYPE>
  interface
    ...
    routines
      push (elem: ELEM_TYPE) is
        == push elem on stack
      pop -> (elem: ELEM_TYPE) is
        == pop top of stack and return it in elem
    ...
```

52
end class  

class client
    uses Stack, Figure
    internal
    var
        s1: Stack <Integer>
        s2: Stack <Figure>
        f: Figure
    ...
        s1.push (42)
        s2.push (f)
end class

Within the generic class (here: Stack) the formal generic parameter (ELEM_TYPE) is a valid type. It can be used for the declaration of variables, parameters and return values.

Classes may have more than one generic parameter. The list of types in the instantiation (in the variable declaration), the actual generic parameters, must have the same number of elements as the list of identifiers in the class header (the formal generic parameters).

Example:

class List <T1, T2> is
    ...
end class

class Client
    uses List
    internal
    var l: List <String, Integer>
    ...
end class

17.1.1 Classes and Types

It is time to revisit the relationship between classes and types. Up to now, we did not distinguish very carefully between classes and types. This was no problem because, as long as no generic classes are involved, each class corresponds to exactly one type, and the type is referred to by using the class's name. Consider

    var
        p : Person

We sometimes said "p is of class Person" or "p is of type Person". A more accurate (but rather clumsy) way to describe this situation is "p is declared of the type that is defined by class Person".
This now changes with generic classes.
Consider the class `Stack` shown above. `Stack` is a class, but because it is
generic, it does not correspond to one type, but potentially to many types.
Since variable declarations need types (not classes), the following declaration
is wrong:

```plaintext
var
  s : Stack  -- ERROR  ERROR  ERROR!
```

Instantiating the generic parameters of a generic class generates a type. In
other words: `Stack` is not a type, but `Stack<Integer>` is. The following
declaration is legal:

```plaintext
var
  s : Stack <Integer>
```

Almost everywhere in Blue code where we have seen class names so far, it is
actually a type that is expected. (The BNF has always shown this correctly.)
For example, the actual generic parameters themselves are types, not classes.
Thus the following declaration is an error:

```plaintext
var
  n: Array <Stack>  -- ERROR!
```

The declaration should read

```plaintext
var
  n : Array <Stack <Integer>>
```

The same is true for a supertype of a class, e.g.

```plaintext
class PersonList is List <Person> ...
```

and for type specifications in the `is` expression (see section 10.3).
The only place where class names, rather than a type, are specified is in the
uses clause of a class.

17.1.2 Operations on Formal Generic Types

Since the dynamic type (the real type at runtime) of an object held in a variable
declared of a formal generic type is not known, no special operation can be
allowed on that object.
Consider:

```plaintext
class List <ELEM_TYPE> is
  ...
  var
    head : ELEM_TYPE
  ...
end class
```

At compile time of class `List` the dynamic type of `head` cannot be known.
Only operations that are allowed on all objects of any class are allowed on
variables of this type. These operations are listed in section 3.2. They are
assignment and equality check (:=, =, <>).
17.2 Constrained Genericity

The generic type in generic classes can be constrained.

Example:

```haskell
class List <T is Comparable> is
  ...
end class
```

In this example, all actual generic parameters for List in instantiations of this class have to be of a subclass of Comparable (or of class Comparable itself).

As a result of this, operations defined for Comparable can be used within the class List on variables declared of type T.

18. Concepts not included in Blue

This section discusses some concepts included in some other modern object-oriented languages that are not included in Blue.

These concepts are not explained in detail. The discussion rather focuses on their relationship to Blue and explains the motivation for their exclusion from the Blue specification.

18.1 Multiple Inheritance

While being a valuable concept for serious production languages, multiple inheritance is considered not important on the beginners' level. It is not a concept we want to teach in the first year, and its introduction into the language would have raised a series of related problems that would have considerably complicated the language and its implementation (such as repeated inheritance, solving of name clashes, etc.)

18.2 Routine Parameters

Passing routines as parameters is, while sometimes nice for elegant algorithms, not compatible with object-oriented programming. It assumes that pieces of code have an existence of their own, independent from classes.

In a pure object-oriented language code does only exist as part of a class. Passing a routine can be achieved (with a bit of overhead) by defining a class with that routine and passing an object of that class. Also, often constrained genericity offers an alternative solution to problems that can be solved by passing routines.
18.3 User Defined Infix Operators

User defined infix operators serve only for convenience in writing – they add no functional value. More importantly, they can be misused and can easily lead to code that is more confusing and less readable than code written without them. Reading and correctness is more important than writing. While there are a few (frequently mentioned) examples where they can sensibly be used, there seem to be many more where their use is tempting but confusing.

18.4 Function Overloading

In Blue, function overloading is included only for the predefined arithmetic operators (+, -, *, /), which are applicable to both Real and Integer numbers. User defined overloading is not allowed.

The argument is similar as for user defined infix operators: while there are some nice applications, the gain through the use of function overloading seems to be too small and infrequent to make up for the added complexity and potential source of confusion and misunderstanding.

18.5 Union Type

A union type is not needed since similar constructs can be expressed with inheritance and supertypes.

18.6 Explicit Blocks

Explicit (user defined) blocks are not needed since in an object-oriented language functions are expected to be relatively short. Explicit blocks serve mainly as scope restrictions for variables and should be superfluous in that case.
### Appendix A: EBNF

<table>
<thead>
<tr>
<th>Rule</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>class-decl</code></td>
<td><code>::= class identifier [ &quot;&lt;&quot; generics_list &quot;&gt;&quot;] is class-definition</code></td>
</tr>
<tr>
<td><code>class-definition</code></td>
<td>`::= Enumeration enum-class-decl</td>
</tr>
<tr>
<td></td>
<td>[ [ class_type ] general-class-decl</td>
</tr>
<tr>
<td><code>enum-class-decl</code></td>
<td>`::= class-comment</td>
</tr>
<tr>
<td></td>
<td>manifest ident-list</td>
</tr>
<tr>
<td></td>
<td>end class</td>
</tr>
<tr>
<td><code>general-class-decl</code></td>
<td>`::= class-comment</td>
</tr>
<tr>
<td></td>
<td>uses [ ident-list ]</td>
</tr>
<tr>
<td></td>
<td>[ internal</td>
</tr>
<tr>
<td></td>
<td>[ const const-decls ]</td>
</tr>
<tr>
<td></td>
<td>[ var var-decls ]</td>
</tr>
<tr>
<td></td>
<td>[ routines routine-decls ]</td>
</tr>
<tr>
<td></td>
<td>interface</td>
</tr>
<tr>
<td></td>
<td>[ creation [ &quot;(&quot; parameter-list &quot;)&quot; ] is routine-impl ]</td>
</tr>
<tr>
<td></td>
<td>[ routines routine-decls ]</td>
</tr>
<tr>
<td></td>
<td>[ invariant condition ]</td>
</tr>
<tr>
<td></td>
<td>end class</td>
</tr>
<tr>
<td><code>generics-list</code></td>
<td>`::= generics-decl { &quot;,&quot; generics-decl }</td>
</tr>
<tr>
<td><code>generics-decl</code></td>
<td>`::= formal-generic-param [ is class-type ]</td>
</tr>
<tr>
<td><code>formal-generic-param</code></td>
<td><code>::= identifier</code></td>
</tr>
<tr>
<td><code>const-decls</code></td>
<td>`::= const-decl { const-decl }</td>
</tr>
<tr>
<td><code>const-decl</code></td>
<td><code>::= ident-list &quot;::&quot; type initialisation</code></td>
</tr>
<tr>
<td><code>var-decls</code></td>
<td>`::= var-decl { var-decl }</td>
</tr>
<tr>
<td><code>var-decl</code></td>
<td>`::= ident-list &quot;::&quot; type [ initialisation ]</td>
</tr>
<tr>
<td><code>ident-list</code></td>
<td>`::= identifier [ &quot;,&quot; identifier ]</td>
</tr>
<tr>
<td><code>initialisation</code></td>
<td><code>::= &quot;:=&quot; expression</code></td>
</tr>
<tr>
<td></td>
<td></td>
</tr>
</tbody>
</table>
Blue Specification, Version 1.0, 24.7.98

```plaintext

```
type ::= class-type
      | formal-generic-param

class-type ::= identifier [ "<" actual-generics-list "]"

actual-generics-list ::= type { "," type }

```

```plaintext

```
routine-decls ::= routine-decl { routine-decl }
routine-decl ::= identifier [ "(" parameter-list ")" ]
               [ "," column ]
               is
               routine-impl

parameter-list ::= param-decl { "," param-decl }
param-decl ::= identifier ":;" type

routine-impl ::= deferred routine-spec
               builtin routine-spec
               [ redefined ] routine-body

routine-body ::= routine-comment
               [ pre condition ]
               [ const const-decls ]
               [ var var-decls ]
               do
               statement-list
               [ post condition ]
               end identifier

routine-spec ::= routine-comment
               [ pre condition ]
               [ post condition ]
               end identifier

statement ::= assignment
             assignment-attempt
             procedure-call
             return
             assertion
             conditional
             selection
             loop

assignment ::= indexed-ident-list ":=" expression-list

assignment-attempt ::= indexed-ident-list ":?=" expression-list

```
### Syntax Definitions

- **indexed-ident-list** ::= indexed-ident `,` indexed-ident
- **expression-list** ::= expression `,` expression
- **procedure-call** ::= call-chain
- **return** ::= return
- **assertion** ::= assert `(" condition ")`
- **conditional** ::= if condition then statement-list { elseif expression then statement-list } [ else statement-list ] end if
- **condition** ::= boolean-expression
- **statement-list** ::= { statement }
- **selection** ::= case expression of { set_expr `:` statement-list } [ else statement-list ] end case
- **loop** ::= loop statement-list exit on condition statement-list { exit on condition statement-list } end loop
**boolean-expression** ::= expression

**expression** ::= function-call
operator-expression
reference-equality
type-equality
in
create
old
this

**function-call** ::= call-chain

**call-chain** ::= [ super "!" ] unqualified-call { "." unqualified-call }

**unqualified-call** ::= identifier [ "(" expression-list ")" ]

**operator-expression** ::= . . .
expression ">=" expression
entity

**reference-equality** ::= expression comparison expression

**comparison** ::= "=" | "<>"

**type-equality** ::= expression **is** type

**in** ::= expression **in** set_expr

**set_expr** ::= 
expression | subrange

**set_elem** ::= expression | subrange

**subrange** ::= expression "." expression
create ::= general-create | array-create
general-create ::= create class-ident ["(" expression-list ")] 
array-create ::= 
old ::= old expression

entity ::= indexed-ident
indexed-ident ::= identifier { 
"[" expression "]" } 
manifest-constant ::= integer-constant 
real-constant
boolean-constant
string-constant 
enum-constant
enum-constant ::= identifier
qualified-ident ::= class-ident ":!" identifier
## Appendix B: Complete List of Aliases

<table>
<thead>
<tr>
<th>alias</th>
<th>original</th>
<th>applicable to</th>
</tr>
</thead>
<tbody>
<tr>
<td>n + m</td>
<td>n.add (m)</td>
<td>Integer, Real</td>
</tr>
<tr>
<td>n – m</td>
<td>n.sub (m)</td>
<td>Integer, Real</td>
</tr>
<tr>
<td>– n</td>
<td>n.neg</td>
<td>Integer, Real</td>
</tr>
<tr>
<td>n * m</td>
<td>n.mult (m)</td>
<td>Integer, Real</td>
</tr>
<tr>
<td>i1 div i2</td>
<td>i1.div (i2)</td>
<td>Integer</td>
</tr>
<tr>
<td>n mod m</td>
<td>n.mod (m)</td>
<td>Integer</td>
</tr>
<tr>
<td>r1 / r2</td>
<td>r1.div (r2)</td>
<td>Real</td>
</tr>
<tr>
<td>n ^ m</td>
<td>n.power (m)</td>
<td>Integer, Real</td>
</tr>
<tr>
<td>a &lt; b</td>
<td>a.greater (b)</td>
<td>Integer, Real, String</td>
</tr>
<tr>
<td>a &gt; b</td>
<td>a.less (b)</td>
<td>Integer, Real, String</td>
</tr>
<tr>
<td>a &lt;= b</td>
<td>a.greaterEq (b)</td>
<td>Integer, Real, String</td>
</tr>
<tr>
<td>a &gt;= b</td>
<td>a.lessEq (b)</td>
<td>Integer, Real, String</td>
</tr>
<tr>
<td>not a</td>
<td>a.invert</td>
<td>Boolean</td>
</tr>
<tr>
<td>a or b</td>
<td>a.or (b)</td>
<td>Boolean</td>
</tr>
<tr>
<td>a and b</td>
<td>a.and (b)</td>
<td>Boolean</td>
</tr>
<tr>
<td>s[i]</td>
<td>s.substring (i,1)</td>
<td>String</td>
</tr>
<tr>
<td>a[i]</td>
<td>a.getElem (i)</td>
<td>Array*</td>
</tr>
<tr>
<td>a[i] :=</td>
<td>a.putElem (i, ...)</td>
<td>Array*</td>
</tr>
<tr>
<td>str (a, b, ...)</td>
<td>a.toString.concat (b.toString.concat (...))</td>
<td>any type</td>
</tr>
<tr>
<td>print (a, b, ...)</td>
<td>output.write ( str (a, b, ...) )</td>
<td>any type</td>
</tr>
</tbody>
</table>

* context sensitive
Appendix C: Implementation-Dependent Definitions

This Appendix lists all characteristics of Blue that are explicitly allowed to be different in different implementations. Blue programs should not rely on specific definitions for these features.

The implementation dependent features are

• The values of MAXINT and MININT.
Appendix D: Interfaces of Predefined Classes

D.1 Integer

class interface Integer is

== Author:  Michael Kölling
== Version:  1.0
== Date:    9 October 1996
== Short:   Blue standard Integer class
==
== "Integer" is a standard class of the Blue language. It is used to store
== integer numbers. "Integer" is predefined in Blue and thus does not
== have to be imported (i.e. explicitly listed in the "uses" clause). It
== is automatically known in all classes.
==
== User defined classes cannot inherit from "Integer".
==
==

routines

    neg  -> (result: Integer) is
    == Return the negative value of this number.
    == Alias:  – (prefix) (e.g. –4)

    add (other: Integer) -> (sum: Integer) is
    == Return the sum of this number and the number other.
    == Alias:   + other    (e.g. 3 + 4)

    sub (other: Integer) -> (diff: Integer) is
    == Return the difference of this number and the number other.
    == Alias:   – other    (e.g. 3 – 4)

    mult (other: Integer) -> (prod: Integer) is
    == Return the product of this number and the number other.
    == Alias: * other      (e.g. 3 * 4)

    div (other: Integer) -> (quot: Integer) is
    == Return the integer part of the quotient of this number and other.
    == Alias:  div other    (e.g. 3 div 4)

    mod (other: Integer) -> (rem: Integer) is
    == Return the remainder of the integer division of this number and
    == other.
    == Alias:  mod other    (e.g. 3 mod 4)

    pow (exp: Integer) -> (result: Integer) is
    == Return this number raised to the power specified by exp.
    == Alias:  ^ other      (e.g. 3 ^ 4)
greater (other: Integer) -> (result: Boolean) is  
== Return true if this number is greater than other, otherwise false.  
== Alias: > other (e.g. 3 > 4)

greaterEq (other: Integer) -> (result: Boolean) is  
== Return true if this number is greater or equal than other, otherwise false.  
== Alias: >= other (e.g. 3 >= 4)

greaterEq (other: Integer) -> (result: Boolean) is  
== Return true if this number is less than other, otherwise false.  
== Alias: < other (e.g. 3 < 4)

less (other: Integer) -> (result: Boolean) is  
== Return true if this number is less than other, otherwise false.  
== Alias: < other (e.g. 3 < 4)

lessEq (other: Integer) -> (result: Boolean) is  
== Return true if this number is less or equal than other, otherwise false.  
== Alias: <= other (e.g. 3 <= 4)

toString -> (s: String) is  
== Return a string representation of this number. The standard format object (of class OutputFormat) is used to determine details of the appearance.  
== Alias: str ()

end class

D.2 Real

class interface Real is  
==================================
== Author: Michael Kölling  
== Version: 1.0  
== Date: 9 October 1996  
== Short: Blue standard Real class  
==  
== "Real" is a standard class of the Blue language. It is used to store floating point numbers. "Real" is predefined in Blue and thus does not have to be imported (i.e. explicitly listed in the "uses" clause). It is automatically known in all classes.  
==  
== User defined classes cannot inherit from "Real".  
==  
==================================

creation  
== Never used explicitly. "Real" is a manifest class, and all real values exist automatically during every execution.  
== No further reals can be created at runtime.

routines

neg -> (res: Real)  
== Negation. Return the real that represents the value
== 
== Alias: -
res <> nil

**add** (other: Real) -> (res: Real)
== Addition. Return the real that represents the value
== "this + other".
== Alias: +
  pre
other <> nil
  post
res <> nil

**sub** (other: Real) -> (res: Real)
== Subtraction. Return the real that represents the value
== "this - other".
== Alias: -
  pre
other <> nil
  post
res <> nil

**mult** (other: Real) -> (res: Real)
== Multiplication. Return the real that represents the value
== "this * other".
== Alias: *
  pre
other <> nil
  post
res <> nil

**divide** (other: Real) -> (res: Real)
== Division. Return the real that represents the value
== "this / other".
== Alias: /
  pre
other <> nil and
other <> 0.0
  post
res <> nil

**power** (n: Real) -> (res: Real)
== Power. Return the real that is represents the value
== "this ^ n" ("this to the power of n").
== Alias: ^
  pre
n <> nil
  post
res <> nil

**sqrt** -> (res: Real)
== Square root. Return the real that is represents the value
== "sqrt (this)" ("square root of this").
  post
res <> nil
trunc -> (res: Integer)
== Truncate to integer. Return the integer that is represents the
== whole number part of this number (i.e. "this" is rounded towards
== zero).
  post
res <> nil

round -> (res: Integer)
== Round to integer. Return the integer that is nearest to the
== value of "this". (Fractions of value 0.5 or greater will be
== rounded away from zero, fractions less than 0.5 will be rounded
== towards zero.)
  post
res <> nil

greater (other: Real) -> (res: Boolean)
== Greater than. Return true if the value of this real is greater
== than the value of "other". Return false otherwise.
== Alias: >
  pre
other <> nil
  post
res <> nil

greaterEq (other: Real) -> (res: Boolean)
== Greater or equal. Return true if the value of this real is
== greater than or equal to the value of "other". Return false
== otherwise.
== Alias: >=
  pre
other <> nil
  post
res <> nil

less (other: Real) -> (res: Boolean)
== Less than. Return true if the value of this real is less
== than the value of "other". Return false otherwise.
== Alias: <
  pre
other <> nil
  post
res <> nil

lessEq (other: Real) -> (res: Boolean)
== Less or equal. Return true if the value of this real is
== less than or equal to the value of "other". Return false
== otherwise.
== Alias: <=
  pre
other <> nil
  post
res <> nil

toString -> (s: String)
== Conversion to String. Returns a string with a printable
== representation of this real number.
D.3 Boolean

class interface Boolean is

invert -> (res: Boolean)
  == Negation. Return "not this", i.e. "false" if this is "true",
  == or "true" if this is "false".
  == Alias: not
  post
  res <> nil

and (other: Boolean) -> (res: Boolean)
  == Logical and. Return the boolean value "this and other".
  == Alias: and
  pre
  other <> nil
  post
  res <> nil

or (other: Boolean) -> (res: Boolean)
  == Logical or (inclusive). Return the boolean value
  == "this or other".
  == Alias: or
  pre
  other <> nil
  post
  res <> nil
**String**

class interface String is

    =-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=
    == Author:  Michael Kölling
    == Version:  1.0
    == Date:    24.9.96
    == Short:   Blue standard String class
    ==
    == This class implements the standard String type for Blue.
    == Strings are a sequence of characters. The first character in a string
    == has the index 1, the last index is equal to length(string).
    ==
    == Blue does not have a separate character type. Characters are represented
    == by a string of length 1.
    ==
    =-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=

creation is

    == Never used - strings are created by writing string literals
    == in double quotes.
    == Example: "This is a string"

routines

    length -> (l: Integer) is
    == Return the length of the string (number of characters).

    concat (s: String) -> (newstring: String) is
    == Return the string which is the concatenation of this string
    == and 's'. This string and 's' remain unchanged.
    == Alias: str ()

    substring (start: Integer, len: Integer) -> (s: String) is
    == Return the substring from this string which starts at 'start'
    == with length 'len'.
    == If "len" is nil, the substring from "start" to the end of the
    == string is returned.
    ==
    pre
    (start >= 1) and (start <= length)
    and
    (len <> nil => (start+len <= length+1))
**find** (s: String, n: Integer) -> (pos: Integer) is
== Return the position of 's' in this string, starting the search at
== 'n'. 'pos' is the index where 's' was found or nil if not found.
  pre
    (s <> nil) and (n >= 1) and (n <= length)

**insert** (s: String, pos: Integer) -> (newstring: String) is
== Return a string that is like this string with 's' inserted into this
== string at position 'pos'.
  pre
    (s <> nil) and (pos >= 1) and (pos <= length)

**delete** (start: Integer, cnt: Integer) -> (newstring: String) is
== Return a string that is like this string with 'cnt' characters deleted,
== starting at position 'start'.
== If "cnt" is nil, all characters from "start" to the end of the
== string are deleted.
  pre
    (start >= 1) and (cnt >= 0) and (cnt <= length)

**less** (s: String) -> (is_less: Boolean) is
== Returns true, if this string is less than 's', else false.
== See class comment (above) for ordering of strings.
  pre
    s <> nil

**greater** (s: String) -> (is_greater: Boolean) is
== Returns true, if this string is greater than 's', else false.
== See class comment (above) for ordering of strings.
  pre
    s <> nil

**lessEq** (s: String) -> (is_less_eq: Boolean) is
== Returns true, if this string is less or equal to 's', else
== false.
== See class comment (above) for ordering of strings.
  pre
    s <> nil

**greaterEq** (s: String) -> (is_greater_eq: Boolean) is
== Returns true, if this string is greater or equal to 's', else
== false.
== See class comment (above) for ordering of strings.
  pre
    s <> nil

**toString** -> (s: String)
== Conversion to String. This function returns the string itself.
== It is an identity function. There is usually no need to call
== this function explicitly. It is provided to make the "str()"
== alias work with all standard classes, including string.
  pre
    s <> nil
post

\texttt{s <> nil}

toUpper -> (s: String)
== Return a string that is a copy of this string with all lower case
== letters replaced by their upper case equivalent. Other character are
== unchanged.

toLower -> (s: String)
== Return a string that is a copy of this string with all upper case
== letters replaced by their lower case equivalent. Other character are
== unchanged.

strip -> (s: String)
== Return a string that is a copy of this string with leading and
== trailing whitespace removed. Whitespace are spaces, \texttt{TAB}
== characters (\texttt{"\t"}) and newlines (\texttt{"\n"}).

fill (fill_char : String, front : Boolean, length : Integer) -> (s: String)
== Return a copy of this string that has characters added to the
== front or back. 'fill_char' is a character that is added to the
== front (if 'front'=true) or back (if 'front'=false) of this string
== until the string has length 'length'. If the length of the
== initial string is already equal to or greater than length, the
== string remains unchanged.
==
== This function can be used to align string output. E.g.
==
==\begin{verbatim}
== \texttt{print (s.fill (" ", true, 20))}
== \end{verbatim}
==
== prints the string s right aligned in a field 20 characters wide
pre
(fill_char.length = 1) and (length >= 0)

caseEqual (other: String) -> (is_equal: Boolean)
== "caseEqual" is a case-insensitive string equality test.
== It returns true, if this string and "other" are equal except
== for possible differences in the case of letters.

ord -> (val: Integer)
== Return the ordinal value of the first character of this string.
== The ordinal value of a character is its internal byte
== representation, usually its ASCII (or, more correctly ISO) code.
== If the length of the string is 0, the result is 0.

hash (limit: Integer) -> (hash_val: Integer)
== Return a hash value for a string. The value is between 0
== and limit.
pre
(limit > 0)
post
hash_val >= 0 and hash_val < limit
end class
class interface Array <ELEM_TYPE> is

creation (size: Integer)
== Creation an array with 'size' elements.

routines

getElem (pos: Integer) -> (elem: ELEM_TYPE)
== Return the element at position 'pos'.
  pre
  pos >= 1  and
  pos <= size

putElem (pos: Integer, elem: ELEM_TYPE)
== Assign 'elem' to position 'pos'.
  pre
  pos >= 1  and
  pos <= size

init (val: ELEM>Type)
== Set all array elements to 'val'.

size -> (sz: Integer)
== Return the current size of the array.
  post
  sz >= 0

setSize (sz: Integer)
== Set the size of the array to 'sz'. If the current size is larger
  == the elements at positions greater than 'sz' will be lost. If the
  == current size is smaller, the elements at the new positions will
  == be undefined.
  pre
  sz >= 0
class interface Enumeration is

===================================================

== Author: Michael Kölling
== Version: 1.0
== Date: 30.11.96
== Short: Abstract superclass for enumerations.
==
== This class serves as a superclass for all enumeration classes. It is
== abstract — no objects can be created of this class directly.
==
===================================================

routines

pred -> (previous: SelfType) is
  == Return the predecessor in this enumeration type. If there is no
  == predecessor, return nil.

succ -> (next: SelfType) is
  == Return the successor in this enumeration type. If there is no
  == successor, return nil.

ord -> (position: Integer) is
  == Return the ordinal position of this element in the enumeration list.
  == The first element has the ordinal 1, the next is 2, and so on.

toString -> (s: String) is
  == Conversion to String. Returns the name of this enumeration value
  == as a string
  == Alias: str ()

end class

Note that "SelfType" is no legal type in Blue. A Blue class like this cannot normally be
written. The meaning of the word "SelfType" in the routine interfaces above indicates that
for every concrete subclass of Enumeration that routine returns the type of the subclass.
For an Enumeration "Colour", for example, the routine pred returns a result of type Colour.

D.7 TextTerminal

class interface TextTerminal is

=====================================================================

== Author: Michael Kölling
== Version: 1.1
== Date: December 1997
== Short: Standard text terminal for the Blue environment
==
== The TextTerminal provides a simple standard terminal for text I/O.
Input and output is buffered by default. The output buffer is flushed at the end of each line (when a "\n" is written) and on a call of an input routine. If single characters have to be written, buffering can be switched off (see below).

---

**creation**

- Create the TextTerminal

**routines**

*inherited from IO_Channel:*

- **write** (s: String) is redefined
  - Write `s` to the terminal

- **readstr** -> (s: String) is redefined
  - Read a string from the terminal. Reads characters until a line break ("\n") is read. The resulting string does not include the line break character.

- **readChar** -> (s: String) is redefined
  - Read the next character from the terminal. Every character (including line break characters) are returned as entered.
  - The line break character is returned as "\n".
  - The input is buffered by default. (This means that editing the input line during input is possible for the user, and the program will return from this function only after a whole line was entered.) For unbuffered input, see **getChar** and **askChar** below.

  - When expecting single character input, note that "readChar" does not discard the NewLine character. If the input is, for instance, "A<Enter>", then "readChar" will return the "A", while the <Enter> remains in the input queue (as a NewLine character). Thus the next input operation will immediately read a NewLine. To deal with this situation, use **readStr** instead or call "endLine" after "readChar".

- **readInt** -> (i: Integer) is redefined
  - Read an integer from the terminal. Skips white space (spaces, tabs and newlines) before the integer. Returns "nil" if the next non-white characters do not represent a number.
  - Numbers are written with an optional - or + sign and digits.

- **readReal** -> (r: Real) is redefined
  - Read a real number from the terminal. Skips white space (spaces, tabs and newlines) before the number. Returns "nil" if the next non-white characters do not represent a real number.
  - Numbers are written with an optional - or + sign, digits and a decimal point.

- **endLine** is redefined
  - Read and discard all characters up to (and including) the next NewLine character.

- **atEnd** -> (isAtEnd: Boolean)
  - Return true if the end of input has been reached. While input from a terminal does not have a natural end, end end-of-input
== condition can be generated by pressing \textsc{CTRL}-D.
== If the next character entered is a \textsc{CTRL}-D, \texttt{atEnd} removes the
== character from the input and returns true. If it is any other
== character, \texttt{atEnd} leaves it pending in the input queue and
== returns false.

\textit{new routines:}

\texttt{getChar} -> (s: String) is redefined
== Get the next character from the terminal (unbuffered).
== This function is similar to \texttt{"readChar"} (see above), but the
== input is not buffered. The function returns as soon as a
== character is entered - no line editing is provided.

\texttt{askChar} -> (s: String)
== Check whether there is a character to be read from the terminal.
== If so, read it.
== This function is similar to \texttt{"getChar"} (see above), but the
== function always returns immediately. If a character has been
== entered, that character is returned, otherwise \texttt{"nil"} is returned.

\texttt{show}
== Show the terminal window.
== This routine opens the window if it was not open, de-iconifies
== the window if it was iconified and brings it to the top of the
== window stack.

\texttt{hide}
== Hide the terminal window.  The window is closed.

\texttt{clear}
== Clear the terminal window and set the cursor to 0,0 (the upper
== left corner)

\texttt{width} -> (columns: Integer)
== Return the current width of the terminal window in characters.
== (Note that the terminal currently does not support resizing.)

\texttt{height} -> (lines: Integer)
== Return the current height of the terminal window in text lines.
== (Note that the terminal currently does not support resizing.)

\texttt{cursorTo} (x: Integer, y: Integer)
== Set the cursor to position (x,y).  The legal ranges for
== x and y are:
== x: 0 .. width-1
== y: 0 .. \texttt{height}-1
== If x or y is outside its range, it is set to the nearest
== legal value.
== This routine works only after the terminal window has been
== exposed at least once (otherwise the terminal size is not known).
== \texttt{pre}
== x <> \texttt{nil} \ and \ y <> \texttt{nil}

\texttt{cursorOn}
== Switch the screen cursor on.  (This is the default.)
**D.8 OutputFormat**

class interface **OutputFormat** is

```
creation
  == Create format object. No parameters needed.
```

routines
setWidth (n : Integer)
   == Set field width for output. Default is 0. If the string
   == representation of a value consists of more characters than
   == the field width, the width is ignored.
   ==
   == Used by Integer, Real, Boolean and Enumeration.
   ==
   == Note that the class String does not use the format object for
   == output. String output is unaffected by the field width. To
   == align strings, use the 'fill' function in class String.
   pre
      n >= 0

alignRight (right : Boolean)
    == Specify alignment of output in the field defined by 'setWidth'.
    == The default is 'left' (spaces will be added to the right of the
    == output as appropriate). If 'right' is set to true, output will
    == be right-aligned within the specified field width.
    ==
    == Used by Integer, Real, Boolean and Enumeration.
    pre
       right <> nil

scientific (useScientific : Boolean)
    == If useScientific is true, scientific notation is used for output
    == of real numbers (e.g. 1.23456e+02). Otherwise fixed point
    == notation is used (e.g. 123.456). The default is useScientific=false.
    ==
    == Used by Real.
    pre
       useScientific <> nil

roundTo (n : Integer)
   == Round output of fixed point real numbers to n digits after the
   == decimal point. Default is 6. Used by Real.
   pre
      n >= 0

getWidth -> (n : Integer)
   == Return the current field width.

isRightAligned -> (right : Boolean)
   == True if right alignment is currently on.

isScientific -> (scientific : Boolean)
   == True if scientific notation is currently on.

getRound -> (n : Integer)
   == Return the current round value.

reset
   == Reset all values to their defaults

end class
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