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Cross Viewpoint Consistency in Open Distributed Processing (Intra language Consistency)

Prepared by Eerke Boiten, Howard Bowman, John Derrick, Maarten Steen

Computing Laboratory
University of Kent at Canterbury
Canterbury
Kent CT2 7NF

Phone +44-1227-764000
Fax +44-1227-762811
Email (E.A.Boiten,H.Bowman,J.Derrick,mwas)@ukc.ac.uk
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Chapter 1

Introduction

1.1 Introduction

Open Distributed Processing (ODP) is recognised as an important standardisation activity. The ODP model seeks to provide an architecture for building potentially global distributed systems with components from many vendors. Thus, ODP will realise the open systems ethos in the distributed systems domain.

A central concept in ODP is that of a viewpoint. Distributed systems are viewed to be so complex that a process of separation of concerns must be employed when describing such systems. Viewpoints provide such a separation of concerns by presenting five distinct views of a single system; these are the enterprise viewpoint, information viewpoint, computational viewpoint, engineering viewpoint and technology viewpoint.

It should be clear that in such viewpoint models it is essential that specifications in different viewpoints are related in order to determine whether the multiple specifications impose conflicting requirements. The project being reported here responds to these needs by investigating how to check that multiple viewpoint specifications are in some sense consistent.

The objective of the project is to perform a verification of the concept of cross viewpoint consistency checking. We will determine the feasibility of performing such checks by developing prototype techniques and tools for relating viewpoint specifications written in Z and LOTOS. These two languages have been chosen because they are formal; enabling formal reasoning to be applied, which we argue is an essential prerequisite for successful consistency checking. In addition, Z and LOTOS represent two of the most different specification techniques being advocated as viewpoint languages, thus, consistency checking between these two languages bounds the difficulty of the problem.

This deliverable describes the initial phase of our work; it focuses on consistency checking methods for individual FDTs. A second deliverable will be produced which extends this work by developing cross-language consistency checks between Z and LOTOS.

1.2 Overview of Project

Figure 1.1 depicts the work plan for the project. Depicted is the full three year trajectory of the project. The two milestone deliverables for BT are shown as bold tasks.

The project is divided into 5 workpackages:

1. WP1: Consistency Framework

The objective of this workpackage is to develop a general framework for consistency in ODP. This will define the basic concepts and mechanisms involved in consistency in a general, FDT independent, fashion. Particular instantiations of the framework can be made by substituting specific FDTs and their correctness properties into the general framework. An important
role of the framework is to ensure that consistency is treated uniformly in each FDT. The framework will provide an interpretation of the main consistency concepts: consistency, unification, translation etc and will define general strategies for consistency checking. The properties of these concepts and strategies will be determined and classified. These properties will highlight the character of the consistency checking problem in ODP. The framework will necessarily be formal in nature.

2. **WP2: Intra Language Consistency**

   Consistency checking between specifications in the same language will be investigated in this workpackage. Specifically, the consistency checking concepts and strategies developed in the previous workpackage will be instantiated for the two formal specification languages, Z and LOTOS. These strategies will be realised by the development of consistency checking tool support for both languages.

3. **WP3: Inter Language Consistency**

   Consistency checking between specifications in different languages will be investigated in this workpackage. As a verification of concept, cross language consistency checking between Z and LOTOS will be explored. This is an extremely demanding area for which there are currently few positive research results. It is intended that tool support for such Z and LOTOS inter language consistency checking will be developed.

4. **WP4: Case Study**

   In this workpackage the suitability of the techniques and tools developed in the previous workpackage will be investigated.
workpackages will be assessed against a number of case studies. A significant part of this workpackage will involve locating suitable examples of ODP systems which can act as case studies. These example ODP specifications must be in multiple viewpoints and have both Z and LOTOS viewpoint specifications.

5. WP5: Dissemination and Collaboration

Dissemination of the results of the project will be targeted at three main groups: the ODP standardisation community, interested industrial parties and the distributed systems research community. Dissemination to the first two of these groups will be facilitated by our collaboration with B.T. under the Formosa project. Two deliverables will be produced for B.T.; the first describing our results on intra language consistency and the second describing our results on inter language consistency. In addition, contributions will be made directly to the ODP standardisation forum through the BSI and to the research community through papers at major conferences and in learned journals.

Each of the five workpackages is divided into tasks. We list these here. The role of each task should be evident from the tasks title.

WP1, CF.1: Initial formulation of Consistency Framework
WP1, CF.2: Study of RM-ODP Definitions
WP1, CF.3: Refined Consistency Framework
WP1, CF.4: Location of Correspondence Rules
WP1, CF.5: Consistency Framework - Final Revision

WP2, INTRA.1: Z Consistency Techniques
WP2, INTRA.2: LOTOS Consistency Techniques
WP2, INTRA.3: Z Consistency Tool
WP2, INTRA.4: LOTOS Consistency Tool

WP3, INTER.1: Preliminary Study of Potential Approaches
WP3, INTER.2: Z and LOTOS Consistency Techniques
WP3, INTER.3: Z and LOTOS Consistency Tool

WP4, CS.1: Location of Possible Case Studies
WP4, CS.2: Z to Z Case Studies
WP4, CS.3: LOTOS to LOTOS Case Studies
WP4, CS.4: Z and LOTOS Inter Lang. Consistency Case Studies

WP5, DC.1: First BT Deliverable
WP5, DC.2: Input to RM-ODP Part 1
WP5, DC.3: Second BT Deliverable
WP5, DC.4: Final Deliverable and Recommendation to ODP

1.3 Overview of deliverable

This deliverable has the following structure:-

- **Chapter 1: Introduction.** This first chapter introduces the project and describes the structure of the deliverable.

- **Chapter 2: Background.** Background on the ODP initiative and the role of viewpoints within this work is presented in this chapter. Three aspects of ODP are considered with reference to the viewpoints model: the architectural semantics, system development for ODP and conformance assessment.
• **Chapter 3: Definitions of Consistency.** The consistency framework developed during the first phase of the project is presented in this chapter. Central to this framework is a precise definition of consistency. We argue that this definition is general enough to embrace all the interpretations of consistency that have already been proposed within ODP. In particular, we show how our definition of consistency can be related to all the interpretations of consistency in the RM-ODP.

• **Chapter 4: Consistency in LOTOS.** Consistency checking within LOTOS is investigated in this chapter. We present a number of possible LOTOS instantiations of the framework concepts and then relate these instantiations. In addition, we present specific mechanisms for consistency checking in LOTOS and give an example of a LOTOS consistency check.

• **Chapter 5: Consistency in Z.** This chapter describes the work on consistency checking in Z. A general algorithm for unifying two Z specifications is presented. Consistency checking by validating the implementability of the derived unification is also explored. An example of unification and consistency checking of two Z specifications is presented.

• **Chapter 6: Conclusions.** The deliverable is summarised and concluded in this chapter.
Chapter 2

Background

In this chapter, some of the developments in the light of which the research project described in this deliverable should be seen, are high-lighted. Our research was mainly triggered by the progressing standardisation of the Reference Model for Open Distributed Processing (RM-ODP). In section 2.1, a brief overview is given of the ODP concepts most relevant to this deliverable. This project is also closely related to the DTI and EPSRC funded FORMOSA project between BT and the University of Stirling. The major aim of the FORMOSA project is to advance the formulation of architectural semantics for the ODP standards using the Formal Description Techniques (FDTs) LOTOS and Z. In section 2.2, we identify the relationships between our work on consistency checking and the work on defining an architectural semantics for ODP, such as in the FORMOSA project. Another project that influenced our research on consistency checking is the PROST project.

2.1 The Reference Model for Open Distributed Processing

The standardisation of a Reference Model for Open Distributed Processing [30] is a joint effort of the International Standardisation Organisation (ISO) and the International Telecommunication Union (ITU-T). The objective is to enable the construction of distributed systems in a multi-vendor environment through the provision of a general architectural framework that such systems must conform to. One of the cornerstones of this framework is a model of multiple viewpoints which enables different participants each to observe a system from a suitable perspective and at a suitable level of abstraction [37]. Section 2.1.1 deals in more detail with the ODP viewpoints.

2.1.1 Viewpoints

The complete specification of any non-trivial distributed system involves a very large amount of information. Attempting to capture all aspects of the design in a single description is generally unworkable. Most design methodologies aim to establish a coordinated, interlocking set of models each aimed at capturing one facet of the design, satisfying the requirements which are the concern of some particular group involved in the design process.

In ODP, this separation of concerns is established by identification of five viewpoints, each with an associated viewpoint language which expresses the concepts and the rules relevant to a particular area of concern.

The viewpoints defined in the Reference Model for ODP are: Enterprise viewpoint, Information viewpoint, Computational viewpoint, Engineering viewpoint and Technological viewpoint, see figure 2.1.

The enterprise viewpoint is concerned with business policies, management policies and human user roles with respect to the systems and the environment with which they interact. The use of the word enterprise here does not imply a limitation to a single organisation. The model constructed may well describe the constraints placed on the interaction of a number of distinct
organisations. The enterprise language introduces concepts to support the expression of policy, particularly with regard to agreements and responsibilities between parts of the enterprise. Agents perform actions and artifacts represent resources. Agents can be assigned roles in a contract which expresses permissions or prohibitions. Groupings of agents are considered as communities which may be administered by a particular agent.

The information viewpoint is concerned with information modelling. By factoring an information model out of the individual components, it provides a consistent common view which can be referenced by the specifications of information sources and sinks and the information flows between them. The information language defines concepts for information schema definition. The language distinguishes between an instantaneous view of information (a static schema), a statement of information which is necessarily unchanged by the system (an invariant schema) and a description of information reflecting the behaviour and evolution of the system (a dynamic schema).

The computational viewpoint is concerned with the algorithms and data flows which provide the distributed system function. This viewpoint specifies the individual components which are the sources and sinks of information flows. The computational language enables the representation of the system and its environment in terms of objects which interact by transfer of information via interfaces. Interfaces are given types and rules are defined for the matching of these types, so that object interfaces need not be specified in an identical way in order to enable interaction between the objects. Objects are chosen to achieve a functional decomposition, but also identify the candidate boundaries for physical distribution.

The engineering viewpoint is concerned with the distribution mechanisms and the provision of the various transparencies needed to support distribution. The engineering language defines a number of functional building blocks which can be combined together to provide the requested transparencies.

The technology viewpoint is concerned with the hardware and software components from which the distributed system is constructed.

Requirements and specifications of an ODP system can be made from any of these viewpoints (as depicted in figure 2.1). However, these viewpoints are not independent. They are each partial views of the complete system specification. Some entities can, therefore, occur in more than one viewpoint, and there are a set of consistency constraints arising from the correspondences between
terms in two viewpoint languages and the statements relating the various terms within each language. The checking of such consistency is an important part of demonstrating the correctness of the full set of specifications.

2.2 Architectural Semantics for ODP

Each viewpoint language consists of a set of definitions and a set of rules which constrain the ways in which the definitions can be related. The notion of language used is an abstract one; the rules are, in effect, the foundations for the grammar of a set of possible detailed languages or notations.

The reference model is not prescriptive in the choice of a specific notation; rather the intention is that a number of existing notations will be used as viewpoint languages, by supporting the concepts and rules defined in the RM-ODP. To this end, a clear interpretation of the architectural concepts of the reference model should be available for those formal notations.

The need for an architectural semantics was recognised from the start of the work on the ODP reference model and is reflected by the inclusion of the architectural semantics as Part 4 of the standard. RM-ODP Part 4 provides an interpretation of the ODP modelling and specification concepts in LOTOS, Estelle, SDL and Z. Thus, this work will act as a bridge between the ODP model and the semantic models of the FDTs and will enable formal descriptions of standards for ODP systems to be developed in a sound and uniform way.

In order to achieve these requirements, the architectural semantics should be consistent in two ways. Firstly, it is necessary to demonstrate that the interpretations of the same architectural entity in different FDTs are consistent. Secondly, the architectural semantics of different viewpoints are related and should therefore be checked for consistency.

The architectural semantics will also provide the basis for uniform and consistent comparison between formal descriptions of the same system or standard in different FDTs. It is, therefore, of great significance to realistic consistency checking techniques.

2.3 ODP System development and conformance assessment

The RM-ODP provides a general architectural framework for the specification of open distributed systems. It does not prescribe a particular system development methodology. The PROST project has investigated a general system development strategy for ODP, which is outlined in section 2.3.1.

Within the framework, domain specific ODP standards can be formulated. When particular distributed systems conform to ODP standards, this will guarantee interoperability and exchange-ability of components. Conformance assessment for ODP is discussed in section 2.3.2.

2.3.1 System development

A number of development strategies could be envisaged for ODP. The multiple viewpoint approach to specification puts particular requirements on the system development methodology. Each viewpoint specification is, at least potentially, at the same level of abstraction; suggesting that viewpoints are related horizontally relative to a vertical system development. This is in contrast to classic waterfall development methodologies. In the PROST project such a, fully general, system development methodology has been investigated. The general development scenario is depicted in figure 2.2. The methodology promotes a number of specification to specification transformations, such as translation, refinement and unification, with the aim to derive a composite ‘implementation’ specification from the multiple viewpoint specifications.

Translation maps a specification from one language onto another while maintaining semantic equivalence. Refinement has the usual meaning of making a specification less abstract, and thus bringing it closer to an implementation while maintaining the captured requirements. Unification is a transformation which enables specifications to be combined.
Consistency is implicit in such a system development methodology. For example, two specifications would be viewed as inconsistent if a common unified specification did not exist. Thus, consistency arises during unification of specifications in models of ODP system development.

### 2.3.2 Conformance assessment

Conformance assessment has been considered from early on in the work on ODP. The meaning of conformance has been built into the RM-ODP. An ODP system conforms to an ODP standard if it satisfies the conformance requirements of that standard.

Conformance assessment of distributed systems is potentially more complex and costly than in traditional communications protocols. The complexity inherent in a distributed system infrastructure may require that in addition to direct testing, techniques such as verification and validation may be necessary to achieve sufficient confidence in the conformance of a particular product.

Results from the PROST project indicate that conformance assessment for distributed systems can be divided into two categories of activities: conformance testing and specification checking. Conformance testing is the activity that relates real implementations to specifications by applying a series of tests to the implementation that were derived from the specification. Specification checking involves activities such as validation, verification and consistency checking. Both validation and verification apply to specifications within one viewpoint. Validation checks a specification against its requirements. Verification checks a refinement against its specification. Consistency checking, potentially, relates specifications across viewpoints. Obviously, no product can conform to a set of viewpoint specifications that are inconsistent.
2.4 Summary

The research project on consistency checking, described in the rest of this deliverable, was mainly triggered by the progression of the standardisation of the Reference Model for Open Distributed Processing. One of the cornerstones of the RM-ODP is a model of multiple viewpoints from which ODP systems can be specified. One of the consequences of adopting a multiple viewpoint approach is that descriptions of the same or related objects can appear in different viewpoints and must co-exist. Consistency of specifications across viewpoints thus becomes a central issue.

The ongoing work on the definition of architectural semantics for ODP manifests important input to the work on consistency checking as it provides a formalisation of the ODP viewpoint languages. Conversely, techniques for consistency checking can also prove valuable in validating the architectural semantics.

Although the RM-ODP does not prescribe a particular system development methodology, it is clear that the multiple viewpoint approach to specification requires a horizontal relating of viewpoint specifications with respect to a vertical development. Several transformations of specifications may be required, such as translation and unification, which also play a role in consistency checking.

Conformance assessment for ODP encompasses both conformance testing (i.e. relating real implementations to specifications) and specification checking (i.e. relating specifications to specifications). Verification of cross viewpoint consistency is an important example of specification checking.
Chapter 3

Definitions of Consistency

In order for consistency to be treated uniformly, a single basic definition of the concept must be given. We seek a fully general interpretation that can be instantiated for different languages as appropriate; e.g. can be instantiated for the correctness properties of both Z and LOTOS.

3.1 Three Possible Definitions of Consistency

This section highlights three possible interpretations of consistency. These definitions all appear in the RM-ODP, the first two appear in part 1 (clause 12.2) and the third appears in part 3 [30] (clause 10). Although, the first of these definitions is only alluded to and the IS version of the reference model is considerably less prescriptive about the definition to use than previous drafts of the standard.

Definition 1
(1.1) Two specifications are consistent iff they do not impose contradictory requirements.
(1.2) Two specifications are consistent iff it is possible for at least one example of a product (or implementation) to exist that can conform to both of the specifications.
(1.3) Two specifications are consistent iff they are both behaviourally compatible with the other.

This last interpretation is a rewording of the RM-ODP definition. This is because the RM-ODP definition is expressed in terms of relating specific viewpoints. We are considering more generalised notions of consistency, thus, we have brought the definition into line with the other definitions in order to facilitate a direct comparison. In addition note that all these definitions are symmetric, i.e. if a specification S is consistent with a specification R then R is consistent with S. This is a reasonable intuitive requirement for a large class of consistency problems (see section 3.2.3).

Definition 2 (Behavioural Compatibility) A specification is behaviourally compatible with a second specification, with respect to a set of criteria, if the first specification can replace the second specification without the environment being able to notice the difference in the specification’s behaviour on the basis of the set of criteria.

This definition slightly adapts the RM-ODP presentation of this concept. Specifically, the RM-ODP definition is expressed in terms of objects. However, we would like to be more general than this and hence we have presented the concept in terms of the notion of a specification. In addition, these three consistency interpretations blur over the fact that specifications may be in different FDTs and that it may not be possible to relate specifications directly without some element of translation.

Each of these notions of consistency is intuitively reasonable. However, the question arises: what is the relationship between the interpretations and, in particular, are these definitions of
consistency themselves consistent? In fact, the different interpretations are likely to be applicable in different settings. For example, definition 1 is relevant to consistency checking in a logical setting, e.g., in an FDT such as Z which is based on first order logic.

We seek to reconcile these interpretations through formalisation. We formalise the first notion of consistency as follows,

**Definition 3** \( S_1 \rightarrow C \rightarrow S_2 \iff \neg (\exists \psi \text{ s.t. } S_1 \models \psi \land S_2 \models \neg \psi) \)

where \( \models \) is the satisfaction relation of the specification’s logic. This definition states that two specifications are consistent if and only if there is no property that holds over one of the specifications and its negation holds over the other specification.

To interpret consistency 1.2 we need a formal interpretation of conformance. There is a difficulty here because conformance relates real physical implementations to specifications and implementations are not amenable to formal interpretation. The classical approach to handling this difficulty is to only consider conformance up to a, so called, *implementation specification*. This is a specification that describes a real implementation in as much detail that a direct mapping from the implementation specification to the real implementation can be found. Thus, it is normal just to consider conformance relations between specifications, see [10] [12] [36] for typical approaches. However, implementation specifications relate to real implementations in different ways for different FDTs and, in particular, for some FDTs not all implementation specifications are implementable. This would, for example, be the case for Z, see discussion in section 3.2.2.

Our approach then is to divide conformance testing into two parts. Firstly, we consider conformance up to implementation specifications, using a relation \( \text{conf} \subseteq \text{SPEC} \times \text{SPEC} \), and then we consider conformance of implementation specifications to real implementations, using a relation \( \text{Conf} \subseteq \text{IMP} \times \text{SPEC} \), where \( \text{SPEC} \) is the set of possible ODP specifications and \( \text{IMP} \) is the set of possible ODP implementations.

By way of clarification, \( S_1 \rightarrow \text{conf} \rightarrow S_1 \) expresses the property that specification \( S_2 \) conforms to specification \( S_1 \), i.e., according to tests derived from \( S_1 \), \( S_2 \) cannot be distinguished from \( S_1 \). It should be noted that we have not specified how and what form of tests are derived from \( S_1 \); there are many options for such derivation [10] [12]. In a similar way \( I \rightarrow \text{Conf} \rightarrow S \) expresses the property that \( I \) conforms to \( S \). Interpretation 1.2 is now formalized as:-

**Definition 4** \( S \rightarrow C_{2,1} \rightarrow S_2 \iff \exists S \in \text{SPEC}, I \in \text{IMP} \text{ s.t. } S \rightarrow \text{conf} \rightarrow S_1 \land S \rightarrow \text{conf} \rightarrow S_2 \land I \rightarrow \text{Conf} \rightarrow S \).

i.e., two specifications are consistent iff an implementation specification which conforms to both and a real implementation of the implementation specification can be found. This definition is correct, but is not very useful since it uses \( \text{Conf} \), which is not subject to formal interpretation. In order to resolve this difficulty we introduce the concept of *internal validity* which holds whenever a specification is implementable:-

**Definition 5** \( S \) is *internally valid*, denoted \( \Psi(S) \), iff \( \exists I \in \text{IMP} \text{ s.t. } I \rightarrow \text{Conf} \rightarrow S \)

We will return to this notion of implementability in section 3.2.2. For example, a Z specification which contains contradictions would not be internally valid. Now we can redefine \( C_2 \) in a more usable way:

**Definition 6** \( S \rightarrow C_{2,2} \rightarrow S_2 \iff \exists S \in \text{SPEC} \text{ s.t. } S \rightarrow \text{conf} \rightarrow S_1 \land S \rightarrow \text{conf} \rightarrow S_2 \land \Psi(S) \).

The third and final consistency interpretation hinges on the notion of behavioural compatibility which is defined in terms of an environment and unspecified criteria. We will consider specific instantiations of behavioural compatibility when we look at specific FDTs; at this stage we formulate the interpretation completely generally, for \( bc \) a particular instantiation of behavioural compatibility.

---

1The order of our relations is in accordance with LOTOS conventions and is opposed to Z conventions.
3.2. A GENERAL DEFINITION OF CONSISTENCY

This subsection presents a general definition of consistency. Our work with the three RM-ODP definitions has shown that each is a specialized notion of consistency that is applicable in a certain setting, e.g., \( C_1 \) to consistency in Z, but none of the definitions gives the “big picture” and is general enough to be instantiated reasonably for many FDTs and many notions of ODP consistency. We will give general definitions of the consistency checking relationships: consistency, both intra and inter language, and unification. First though we will present the notation that we will work with. Importantly, this notation reflects the search for a general interpretation of consistency by defining very general notational conventions. These conventions will be specialized for particular FDTs and particular forms of consistency.

3.2.1 Notation

We begin by assuming a set \( \text{COMP} \) of all possible computations. Subsets of \( \text{COMP} \) include \( \text{IMP} \) the set of physical implementations and \( \text{DES} \) the set of formal descriptions. The latter of these is the domain that we will be working in; \( \text{DES} \) contains both formal specifications in languages such as LOTOS and Z and semantic descriptions in notations such as labelled transition systems and ZF set theory.

We assume a set \( \text{DR} \) of description relations. Members of this set relate pairs of descriptions in \( \text{DES} \). \( \text{DR} \) embraces all possible ways of relating descriptions, e.g., refinement relations or semantic maps. For a particular relation \( r \in \text{DR} \), where \( r \subseteq \text{DES} \times \text{DES} \), we define the left and right projections of \( r \) as:
\[
\text{pr}_L(r) = \{ D : \exists D'.s.t. (D, D') \in r \} \quad \text{and} \quad \text{pr}_R(r) = \{ D' : \exists D.s.t. (D, D') \in r \}
\]

\( \text{DR} \) is subdivided into \( \text{DEV} \) the set of development relations and \( \text{SEM} \) the set of semantic maps. Importantly, although members of \( \text{DEV} \) and of \( \text{SEM} \) have very different functions, both can be viewed as relations between pairs of descriptions, possibly in different languages.

Development relations are written \( d \in \text{dev} \) and if \( X \text{ dev } X' \) then, in some sense, \( X \) is a valid development of \( X' \). Our concept of a development relation generalises all notions of evolving a formal description towards an implementation and thus embraces the many such notions that have been proposed. In particular, \( \text{DEV} \) contains refinement relations, equivalences and relations which can broadly be classed as implementation relations [36] such as the LOTOS conformance relation.
Consistency.

These different classes of development are best distinguished by their basic properties. Refinement is typically reflexive and transitive (i.e., a preorder); equivalences are reflexive, symmetric and transitive; and implementation relations are only reflexive. The distinction between refinement and implementation relations is particularly significant; transitivity is a crucial property in enabling incremental development of specifications towards realizations and implementation relations are typically lacking in this respect.

In general, though, we do not require that development relations support any specific properties. In particular, we cannot even assume reflexivity in the general case. This is because, in order to support inter language consistency checking, we allow development relations to relate descriptions in different notations. In these circumstances reflexivity is not a sensible concept.

Members of SEM are semantic maps between descriptions in formal techniques. Typically they map descriptions from one formal technique to a second formal technique. Elements of SEM will usually be denoted $\square$.

Descriptions are written in formal techniques. The set of all such techniques is denoted $FT$. Formal techniques are triplets; they are elements of $P_{DES} \times P_{DEV} \times P_{SEM}$. Thus, every formal technique is characterized by the set of possible descriptions in the notation, a set of associated development relations and a set of semantic maps. We require that the left projection of all elements of $DEV$ and $SEM$ contains a subset of $DES$. For a particular formal technique $ft$ we denote the set of all descriptions in $ft$ as $DES_{ft}$, the set of all development relations as $DEV_{ft}$ and the set of all semantic maps as $SEM_{ft}$.

### 3.2.2 Consistency

**Basic Definition.** In its general form consistency is a check which takes any number of descriptions and returns true if all the descriptions are consistent and false otherwise. This check will be performed according to a list of development relations, one per description, and is denoted, $C[dv_1, dv_2, ..., dv_n](X_1, X_2, ..., X_n)$. The validity of the check has two elements: type correctness and consistency.

**Definition 8 (Type Correctness)** $C[dv_1, dv_2, ..., dv_n](X_1, X_2, ..., X_n)$ is type correct iff $X_1 \in p_r(dv_1) \land X_2 \in p_r(dv_2) \land ... \land X_n \in p_r(dv_n) \land (p_l(dv_1) \cap p_l(dv_2) \cap ... \cap p_l(dv_n) \neq \emptyset)$.

Type correctness ensures, firstly, that for every description the corresponding development relation, i.e., $dv_i$ for $X_i$, is correctly typed with regard to the description. In addition, type correctness ensures that the target types of the relations have some intersection. This check has the function of determining that the consistency check being attempted is sensible. Type correctness will not be an issue for intra language consistency, but will be necessary when determining an appropriate inter language consistency check to apply. When writing $C[dv_1, dv_2, ..., dv_n](X_1, X_2, ..., X_n)$ unless otherwise stated we will assume the check has already been shown to be type correct.

Once type correctness has been determined we can investigate consistency. Intuitively we view $n$ specifications $X_1, X_2, ..., X_n$ as consistent if and only if there exists a physical implementation which is a realization of all the specifications, i.e., $X_1, X_2$ through to $X_n$ can be implemented in a single system. However, we can only work in the formal setting, so we express consistency in terms of a common (formal) description, $X$, and a list of development relations, $dv_1, dv_2, ..., dv_n$. The definition states that $n$ descriptions are consistent if and only if a description can be found which is a development of $X_1$ according to $dv_1$, $X_2$ according to $dv_2$, through to $X_n$ according to $dv_n$, and the third description is internally valid, written $\Psi(X)$. The structure of the consistency check is depicted in figure 3.1 and is formalized in definition 9.

**Definition 9 (Consistency)** $X_1, X_2, ..., X_n$ are consistent by $dv_1, dv_2, ..., dv_n$, i.e., $C[dv_1, dv_2, ..., dv_n](X_1, X_2, ..., X_n)$ holds, iff $\exists X \in DES$ s.t. $(X \ land dv_1 \ X_1 \ land dv_2 \ X_2 \ land ... \ land X \ land dv_n \ X_n) \land \Psi(X)$.

For $n$ descriptions to be consistent this definition requires that $X$ is a common development of $X_i$ for all $i$ between 1 and $n$. Notice that we allow the descriptions to be related to their
common development in different ways, i.e. if \( dv_i \neq dv_j \). This is important in order to support the full generality of ODP viewpoints. A particular specialization of the viewpoints may for example require a viewpoint to be related to a second viewpoint directly by refinement. In order to reflect alternative classes of specialization we will distinguish between balanced and unbalanced consistency. These two alternatives will be discussed in sections 3.2.3 and 3.2.4.

The internal validity check in the above definition formalizes the notion of implementability. It is required because descriptions relate to physical implementations in different ways for different languages and, in particular, for some FDTs not all specifications are implementable. For example, a Z specification that contains an operation \([n! : N | n! = 5 \land n! = 3]\) has no real implementation. Thus, for some FDTs it is possible to find a description which is a common development of a pair of specifications, but is not itself implementable. The property \( \Psi(X) \) is true if and only if the description \( X \) has a real implementation. Thus, \( \Psi \) acts as a receptacle for properties of particular languages that make descriptions in that language unimplementable. For example, a Z specification which contains contradictions would not be internally valid.

In most cases \( X_1, X_2, ..., X_n \) in the above definition will all be specifications, however, \( X \) will commonly be a semantic representation. In particular, if some of \( X_1, X_2, ..., X_n \) are in different languages then \( X \) is almost certain to be in a common semantic notation. The properties that enable a semantic notation to be suitable for representing common developments of specifications in different formal techniques will be discussed in section 3.4. If \( X_1, X_2, ..., X_n \) are in the same formal technique then \( C[dv_1, dv_2, ..., dv_n](X_1, X_2, ..., X_n) \) is called an intra language consistency check and if for some \( i \) and \( j \) between 1 and \( n \), \( X_i \) and \( X_j \) are in different formal techniques then \( C[dv_1, dv_2, ..., dv_n](X_1, X_2, ..., X_n) \) is called an inter language consistency check. We will denote this interpretation of consistency as \( C \) when we refer to it in text.

**Binary Consistency.** An important special case is binary consistency, i.e. the consistency check \( C[dv_1, dv_2](X_1, X_2) \) is performed. Binary consistency is a binary relation and is often written, \( X_1 C_{dv_1, dv_2} X_2 \). The possibility of inter language consistency make some of the standard properties of binary relations problematic. For example, we can consider reflexivity, symmetry and transitivity.

**Proposition 1**

*Binary consistency is neither (i) reflexive, (ii) symmetric or (iii) transitive.*

**Proof**

(i) Reflexivity is the case \( C[dv](X) \), and this will be false whenever \( X \) is not internally valid.

(ii) Assuming \( C[dv_1, dv_2](X_1, X_2) \) is true, in the case of inter language consistency \( C[dv_1, dv_2](X_2, X_1) \) is likely not even to be type correct. Thus, in its most general form symmetry of consistency does
not even yield a type correct consistency check.

(iii) Assuming $C[dv_1, dv_2](X_1, X_2)$ and $C[dv_3, dv_4](X_2, X_3)$ hold then transitivity requires us to show that $X_1$ and $X_3$ are consistent, however, according to what development relations will we check consistency? The transitivity variant that we would like is that $C[dv_1, dv_4](X_1, X_3)$ follows from the assumptions. However, nothing in our assumption guarantees that, $p_1(dv_1) \cap p_1(dv_4) \neq \emptyset$, thus, $C[dv_1, dv_4](X_1, X_3)$ may not be type correct. Furthermore, even if we assume type correctness of $C[dv_1, dv_4](X_1, X_3)$, consistency will not always hold, since $C[dv_1, dv_2](X_1, X_2)$ and $C[dv_3, dv_4](X_2, X_3)$ are likely to have different common developments that cannot be related. □

**Implementation Complete.** There are a number of languages in which all specifications are internally valid. This is for example the case for LOTOS, for which even a deadlock has an implementation equivalent. Thus, we introduce the following notation:-

**Notation 1 (Implementation Complete)** A formal technique $ft$ is called implementation complete iff $\forall X \in DES_{ft}, \Psi(X)$.

The following result is almost trivial, but it enables us to start characterising reflexivity of consistency.

**Proposition 2**
If $X \in D$ for $D \subseteq DES$, then $C[dv_1, dv_2](X, X)$ holds, i.e. reflexivity of consistency, iff $\forall X \in D \exists X' \in DES$ s.t. $X' (dv_1 \cap dv_2) X \land \Psi(X')$.

**Proof**
\((\Rightarrow) X C_{dv_1, dv_2} X \implies \exists X' s.t. X' dv_1 X \land X' dv_2 X \land \Psi(X') \implies X' (dv_1 \cap dv_2) X \land \Psi(X').
\((\Leftarrow) \exists X' s.t. X' (dv_1 \cap dv_2) X \land \Psi(X') \implies X' is the required common development. □

Note that $C[dv_1, dv_2](X, X)$ can also be written as $C[dv_1 \cap dv_2](X)$. Proposition 2 has the following immediate corollary.

**Corollary 1**
If $ft$ is implementation complete and $dv_1$ and $dv_2$ are reflexive, then $\forall X \in DES_{ft}, C[dv_1, dv_2](X, X)$ holds, i.e. reflexivity of consistency.

**Proof**
This result follows from the reflexivity of the development relations, which implies that $X$ is a common development, and from the right to left implication in proposition 2. □

This corollary implies that consistency is reflexive for a language such as LOTOS in which all specifications are internally valid and development is at least reflexive.

**Pairwise Consistency.** An important issue is in what way we can determine consistency, for example, can we assert consistency between three or more descriptions by performing a series of binary consistency checks. In order to determine this we consider the notion of a pairwise consistency check:-

**Definition 10 (Pairwise Consistency)** Descriptions $X_1, X_2, ..., X_n$ are pairwise consistent according to development relations $dv_1, dv_2, ..., dv_n$ iff $\forall X_i, X_j$ s.t. $1 \leq i, j \leq n$, $X_i C_{dv_i, dv_j} X_j$.

The following result characterizes the broad relationship between pairwise and normal consistency.

**Proposition 3**
(i) Consistency implies pairwise consistency.
(ii) Pairwise consistency of three or more specifications does not imply consistency.
3.2. A GENERAL DEFINITION OF CONSISTENCY

Proof
(i) Assume \( \exists X \in \text{DES} \text{ s.t.} (X \ dv_1 \ X_1 \land X \ dv_2 \ X_2 \land \ldots \land X \ dv_n \ X_n) \land \Psi(X) \). Now clearly \( X_i, C_{dv_j}, X_j \) for any \( 1 \leq i, j \leq n \) since \( X \) can act as the internally valid common development.
(ii) We demonstrate this by counterexample. Consider the three specifications: \( S_1 = [x!, y! : \mathbb{N} \mid x! = y!] \), \( S_2 = [z!, z! : \mathbb{N} \mid z! = z!] \) and \( S_3 = [z!, y! : \mathbb{N} \mid z! \neq y!] \) Intuitively these are balanced pairwise consistent, i.e. \( S_1 \ CS \ S_2 \), \( S_2 \ CS \ S_3 \), \( S_1 \ CS \ S_3 \), but, they are not globally consistent. \( \square \)

Intuitively, the second part of the above proposition arises because pairwise consistency only requires the existence of a common development. Thus, many pairwise consistency results may exist each of which focuses on a different common development. This is not sufficient to induce “global” consistency which requires the existence of a single common development. We should also emphasise that the generalized consistency that the ODP viewpoints model induces is our normal consistency, not pairwise consistency, since a single development of all the viewpoints (i.e. the realization) is required. In later sections we will characterize circumstances in which pairwise and normal consistency are the same.

3.2.3 Balanced Consistency

Balanced consistency reflects the situation in which the specifications being checked for consistency are at the same level of abstraction; balanced consistency is written: \( C[dv](X_1, X_2, \ldots, X_n) \). It should be noted that some of our previous papers have only considered balanced consistency, e.g. [9] and presented this as consistency in its entirety. This report presents a generalization of that work.

Definition 11 (Balanced Consistency)
\( C[dv_1, dv_2, \ldots, dv_n](X_1, X_2, \ldots, X_n) \) is balanced iff \( dv_i = dv_j \), \( \forall dv_i, dv_j \text{ s.t.} 1 \leq i, j \leq n \).

We have the following result:

Proposition 4
\( C[dv](X_1, X_2, \ldots, X_n) \) is well founded, i.e. \( C[dv](X_1, \ldots, X_n) = C[dv](Y) \) where \( Y \) is any possible permutation of \( X_1, \ldots, X_n \).

Proof
Immediate from definition of consistency and balanced consistency. \( \square \)

This proposition states that the ordering of \( X_1, \ldots, X_n \) in the argument list of \( C \) is not important in balanced consistency. This is once again a very obvious result, but does contrast with the situation for unbalanced consistency, where ordering is crucial and permuting the order of descriptions may invalidate type correctness.

Once again we can consider the special case of binary balanced consistency, \( C[dv](X_1, X_2) \), which is often written as \( C_{dv} \). The next result follows naturally from the previous result:-

Proposition 5
\( C_{dv} \) is symmetric.

Proof
From proposition 4. \( \square \)

This proposition characterizes symmetry of binary balanced consistency and proposition 2 characterizes reflexivity of all binary consistency; transitivity of \( C_{dv} \) is, however, more difficult to characterize. We have the following partial characterization.

Proposition 6
\( dv \) is transitive and symmetric \( \Rightarrow C_{dv} \) is transitive
Proof Assume \( \exists X, X' \text{ s.t. } X \vdash X_1 \wedge X' \vdash X_2 \wedge \Psi(X) \) and \( X' \vdash X_2 \wedge X' \vdash X_3 \wedge \Psi(X') \); then from symmetry of \( \vdash \) we get \( X_2 \vdash X \) and transitivity can be applied twice to get \( X' \vdash X_1 \).

Thus, \( X' \) is the required common development of \( X_1 \) and \( X_2 \).

The following results which relate the characteristics of the development relation used to the induced balanced consistency are also easily obtained:

**Proposition 7**

(i) If \( \vdash \) is reflexive and \( \Psi(X_1) \), then \( X_1 \vdash X_2 \implies X_1 \ C_{\vdash} X_2 \).

(ii) If \( \vdash \) is symmetric and transitive then \( X_1 \ C_{\vdash} X_2 \implies X_1 \vdash X_2 \).

Proof

(i) Assume \( X_1 \vdash X_2 \) and \( \Psi(X_1) \); from reflexivity of \( X_1 \) we get \( X_1 \) is the required common development.

(ii) Assume \( \exists X \ s.t. \ X \vdash X_1 \wedge X \vdash X_2 \wedge \Psi(X) \); then from symmetry \( X_1 \vdash X \) and from transitivity \( X_1 \vdash X_2 \) as required.

**Corollary 2**

If \( ft \) is implementation complete and \( \vdash \) is an equivalence relation, then for all descriptions in \( ft \),

\( \vdash = C_{\vdash} \).

This result will be valuable when we seek to relate behavioural compatibility to our interpretation of consistency. See section 4.2.2 for a discussion of this.

In addition, as the following result shows, if we impose some strong requirements on the development relation we can relate pairwise consistency to consistency in the balanced case:

**Proposition 8**

\( \vdash \) is transitive and symmetric \( \implies \) (balanced pairwise consistency \( \iff \) balanced consistency).

Proof Assume that \( \vdash \) is transitive and symmetric, then we can prove the equivalence of pairwise consistency and consistency as follows:-

(\( \iff \)) We already have this from 3.

(\( \implies \)) Assume \( \forall X_i, X_j \ s.t. \ 1 \leq i, j \leq n \wedge i \neq j \ \exists Y_k \ s.t. \ 1 \leq k \leq n(n-1)/2 \wedge \Psi(Y_k) \wedge Y_k \vdash X_i \wedge Y_k \vdash X_j \). (In this proof we assume that all \( Y_k \)'s are distinct. This is the worst case situation to prove, if different pairs have the same common development the situation only becomes easier.)

We will show that any of the \( Y_k \)'s could act as a common development for all the descriptions. So, pick \( Y_r \) s.t. \( 1 \leq r \leq n(n-1)/2 \). Firstly, note that \( \Psi(Y_r) \) by the assumption. Now pick \( X_i \ s.t. \ 1 \leq i \leq n \). Clearly, \( \exists X_i, s.t. \ Y_r \vdash X_i \) and \( \exists Y_m \ s.t. \ Y_m \vdash X_t \wedge Y_m \vdash X_i \). From symmetry of \( \vdash \) we get \( X_r \vdash Y_m \) and we can apply transitivity twice to \( Y_r \vdash X_i, X_s \vdash Y_m \) and \( Y_m \vdash X_t \) to get \( Y_r \vdash X_t \) as required.

\[ 3.2.4 \] Unbalanced Consistency

Unbalanced consistency reflects the situation in which the specifications being checked for consistency are at different levels of abstraction or have different granularities. Such relationships are easy to imagine for particular specializations of the ODP viewpoints. In circumstances in which confusion cannot arise unbalanced consistency is denoted \( C[\vdash_{v_1}, \vdash_{v_2}, \ldots, \vdash_{v_n}] (X_1, X_2, \ldots, X_n) \), although if we wish to specifically distinguish the unbalanced case from the general case we will write \( C^u[\vdash_{v_1}, \vdash_{v_2}, \ldots, \vdash_{v_n}] (X_1, X_2, \ldots, X_n) \). In general unbalanced consistency is considerably more difficult to work with than balanced consistency. We have the following general definition:

**Definition 12** (Unbalanced Consistency)

\( C[\vdash_{v_1}, \vdash_{v_2}, \ldots, \vdash_{v_n}] (X_1, X_2, \ldots, X_n) \), is unbalanced iff \( \exists \vdash_{v_i}, \vdash_{v_j} \ s.t. \ 1 \leq i, j \leq n \wedge \vdash_{v_i} \neq \vdash_{v_j} \).

Some results can be derived on binary unbalanced consistency, which is often denoted as \( C_{\vdash_{v_1}, \vdash_{v_2}} \) or as \( C^u_{\vdash_{v_1}, \vdash_{v_2}} \) if we wish to distinguish this class of consistency from the general binary case.
Proposition 9
If \( ft \) is implementation complete and \( dv \) is a preorder (i.e. reflexive and transitive) then \( \forall X_1, X_2 \in DES_{ft}, X_1 \ dv X_2 \iff X_1 \ C_{dv^{-1}, dv} X_2. \)

Proof
(\( \Rightarrow \)) Firstly, \( X_1 \ dv X_2 \) by assumption, but also \( X_1 \ dv^{-1} X_1 \) by reflexivity of \( dv \). So, \( X_1 \) is the required common development.

(\( \Leftarrow \)) Assume \( \exists X \) s.t. \( X_1 \ dv X \land X \ dv X_2 \) then by transitivity of \( dv, X_1 \ dv X_2. \)

Interestingly, we also have:

Proposition 10
If \( ft \) is implementation complete and \( dv \) is a preorder then \( \forall X_1, X_2 \in DES_{ft}, X_1 \ dv X_2 \iff X_1 \ C_{(dv \cap dv^{-1})}, dv X_2. \)

Proof
(\( \Rightarrow \)) \( X_1 \ dv X_2 \) by assumption, also \( X_1 \ dv \cap dv^{-1} X_1 \) by reflexivity.

(\( \Leftarrow \)) Since \( X_1 \ C_{(dv \cap dv^{-1})}, dv X_2 \ \Rightarrow X_1 \ C_{dv^{-1}, dv} X_2 \) and using previous proposition

This final result characterizes the relationship between \( dv \) and \( C_{\omega, dv} \) where \( \omega \) is the equivalence defined by \( \omega = dv \cap dv^{-1} \). These results will be important when we characterize the consistency arising from the LOTOS refinement preorders, e.g. reduction and extension, and their equivalence, testing equivalence.

Corollary 3
For implementation complete formal techniques and \( dv \) a preorder, \( dv = C_{dv^{-1}, dv} = C_{(dv \cap dv^{-1})}, dv. \)

3.3 Unification

Unification is the mechanism by which descriptions are composed in such a way that the composition is a development of all the descriptions.

Definition 13 (Unification Set) \( \mathcal{U}[dv_1, dv_2, \ldots, dv_n](X_1, X_2, \ldots, X_n) = \{ X : X \in DES \land (X \ dv_1 X_1 \land X \ dv_2 X_2 \land \ldots \land X \ dv_n X_n) \}. \)

The unification set is the set of all common developments of a list of descriptions, i.e. the set of all unifications. Clearly, \( C[dv_1, dv_2, \ldots, dv_n](X_1, X_2, \ldots, X_n) \) holds if and only if \( \exists X \in \mathcal{U} \) such that \( \Psi(X) \). In fact, one approach to consistency checking is to perform a unification and then to show that this unification is internally valid. This will be the approach taken for consistency checking in \( Z \).

In the same way as for consistency we can consider binary unification and distinguish between balanced and unbalanced unification. We begin by considering balanced unification.

3.3.1 Balanced Unification

Definition 14 \( \mathcal{U}[dv](X_1, X_2, \ldots, X_n) \) is a balanced unification set iff \( \forall dv_i, dv_j \) s.t. \( 1 \leq i, j \leq n, dv_i = dv_j. \)

The balanced unification set is denoted \( \mathcal{U}[dv](X_1, X_2, \ldots, X_n) \) and the following results are easily obtained:

Proposition 11
(i) \( \mathcal{U}[dv](X_1, X_2, \ldots, X_n) \) is well founded, i.e. any permutation of the order of \( X_1, X_2, \ldots, X_n \) will give the same unification set.

(ii) \( \mathcal{U}[dv](X_1, X_2) = \mathcal{U}[dv](X_2, X_1) - \text{commutativity} \)

(iii) \( \forall X \in \mathcal{U}[dv](X_1, X_2, \ldots, X_n), X \ dv X_1, X_2, \ldots, X_n - \text{common development} \)

Proof Straightforward.
3.3.2 Unbalanced Unification

**Definition 15** \( U[dv_1, dv_2, \ldots, dv_n](X_1, X_2, \ldots, X_n) \) is an unbalanced unification set iff \( \exists \) \( dv_i, dv_j \) s.t. \( 1 \leq i, j \leq n \land dv_i \neq dv_j \).

Wellfoundedness and commutativity will not hold for unbalanced unification, common development will hold.

3.3.3 Representative Unification

A particular unification algorithm will construct just one member of the unification set. Importantly, we need to know that the unification that we construct is internally valid if and only if an internally valid unification exists; otherwise we may construct an internally invalid unification despite the fact that an alternative unification may be internally valid.

Thus, we introduce the concept of a representative unification, which is defined as follows:-

**Definition 16** \( X \in U[dv_1, dv_2, \ldots, dv_n](X_1, X_2, \ldots, X_n) \) is a representative unification iff \( (\exists X' \in U[dv_1, dv_2, \ldots, dv_n](X_1, X_2, \ldots, X_n) \text{ s.t. } \Psi(X')) \implies \Psi(X) \).

We denote a representative unification as \( U[dv_1, dv_2, \ldots, dv_n](X_1, X_2, \ldots, X_n) \). The following result is very straightforward:-

**Proposition 12**

It is implementation complete and \( X_1, \ldots, X_n \in DES_{\mu} \implies \forall X \in U[dv_1, \ldots, dv_n](X_1, \ldots, X_n), X \) is a representative unification.

We believe that the least unification will, in general, generate a representative unification. The importance of taking the least unification is that the contradictions contained in the unification will reflect contradictions occurring in the original specifications and will not have been introduced during development.

Unfortunately, for inter language consistency the least of the set of unifications is a problematic concept. Specifically, descriptions in the unification set, \( U[dv_1, dv_2, \ldots, dv_n](X_1, X_2, \ldots, X_n) \), are likely to be in a different notation from \( X_1, X_2, \ldots, X_n \), thus it is unlikely that the unifications can be related in a type correct manner using \( dv_1, dv_2, \ldots, dv_n \). Thus, we will consider the least unification in the intra language case.

**Definition 17 (Least Unification (Intra Language))** \( X \in U[dv_1, dv_2, \ldots, dv_n](X_1, X_2, \ldots, X_n) \) is the least unification iff \( \neg \exists X' \in U[dv_1, dv_2, \ldots, dv_n](X_1, X_2, \ldots, X_n), \text{ s.t. } X \text{ (} dv_1 \cap dv_2 \cap \ldots \cap dv_n \text{) } X' \land \neg (X' \text{ (} dv_1 \cap dv_2 \cap \ldots \cap dv_n \text{) } X) \).

We denote the least unification \( LU[dv_1, dv_2, \ldots, dv_n](X_1, X_2, \ldots, X_n) \).

3.4 Inter Language Consistency

The basic definition of consistency that we presented in section 3.2.2 is able to relate descriptions in different formal techniques and thus to support inter language consistency. This would be the case if the unification sought is a description in a common semantics, i.e. a semantic notation that can represent the formal techniques of both the original descriptions. However, it is important that we consider what constitutes a satisfactory common semantic representation. This section makes a first effort to define the criteria that such a common semantics should satisfy.

We use the notion of a development relation in different notations correlating under certain conditions. These conditions amount to a full abstraction requirement.

**Definition 18 (Correlation between relations)** A development relation \( dv_2 \) correlates to a development relation \( dv_1 \), written \( dv_1 \rightarrow dv_2 \) with regard to a semantics \([[X]]\) for a formal technique \( ft \) iff \( \forall X_1, X_2 \in DES_{\mu}, X_1 \vdash dv_1 X_2 \iff [[X_1]] \vdash dv_2 [[X_2]] \).
**3.5 Conclusion**

This section concludes this chapter with a discussion of the properties of our consistency definition. In particular, the next section considers the generality of $C$.

### 3.5.1 Generality of the Definition

**Reconciling the RM-ODP Definitions**

As we have indicated previously all of the three RM-ODP definitions are reasonable, but each is relevant to consistency checking in a particular setting. What we have sought in this chapter is a single general definition, which can be instantiated as appropriate for different notations. We believe we have obtained such a general interpretation. In particular, our definition embraces one of the reference models definitions directly and the other two through imposition of certain constraints on the development relation used. We will consider these results here.

**Reconciling $C_{2.2}$.** The following proposition represents a straightforward instantiation into $C$.

**Proposition 13**

If $dv$ is instantiated as conformance then $C_{dv} = C_{2.2}$.

**Reconciling $C_1$.** Our approach is to define a development relation with the required characteristics and instantiate this into $C$. We define $dev$ as:

1. $(dev.1) \ X_1 \ dev \ X_2 \iff (X_1 \models \gamma \iff X_2 \models \gamma)$ and
2. $(dev.2)$ $dev$ is a reflexive development relation.

These represent strong constraints on development. However, neither are unreasonable and could be defined in specific settings, e.g. the first constraint could be defined for $Z$ and the second constraint is a standard requirement of refinement relations. In addition, we define internal validity as,

$\Psi(X) \iff \neg(\exists \gamma \ such\ that\ X_1 \models \gamma, \neg \gamma)$

which is a natural instantiation.

We need a simple lemma.

**Lemma 1**

$X_1 \ C_1 \ X_2 \implies (\neg(\exists \gamma \ such\ that\ X_1 \models \gamma, \neg \gamma) \land \neg(\exists \gamma' \ such\ that\ X_2 \models \gamma', \neg \gamma'))$.

**Proof**

We will show that $X_1 \ C_1 \ X_2 \implies \neg(\exists \gamma \ such\ that\ X_1 \models \gamma, \neg \gamma)$. We will use contradiction, so assume $\exists \gamma \ such\ that\ X_1 \models \gamma, \neg \gamma$. Now if we consider $X_2$ it is clear that either $X_2 \models \gamma \lor X_2 \models \neg \gamma$. However, if either of these hold then $C_1$, i.e. $\neg(\exists \gamma'' \ such\ that\ X_1 \models \neg \gamma'' \land X_2 \models \gamma'')$, is contradicted. This gives us the required contradiction. We can make a similar argument to show that $\neg(\exists \gamma' \ such\ that\ X_1 \models \gamma', \neg \gamma')$ follows from $C_1$. □

Now we can prove the equality that we want.

**Proposition 14**

$C_1 = C_{dev}$.

**Proof**

$(C_1 \implies C_{dev})$

Assume $C_1$, i.e. $\neg(\exists \gamma \ such\ that\ X_1 \models \neg \gamma \land X_2 \models \gamma)$. From this we can draw the following implications:-
Thus $X_1 \text{ dev } X_2$, by (dev.i). Now by (dev.ii) we also have that $X_1 \text{ dev } X_1$. So, $X_1$ is a common development and from lemma 1 we have that $\Psi(X_1)$. Thus, $X_1 C_{\text{dev}} X_2$ as required.

(C$_1$ $\iff$ C$_{\text{dev}}$) We will use contradiction. Thus, assume $C_{\text{dev}}$ and the negation of $C_1$:

$\exists X \text{ s.t. } X \text{ dev } X_1 \land X \text{ dev } X_2 \land \exists (\exists \gamma \text{ s.t. } X \models \gamma, \neg \gamma) \land \exists \gamma' \text{ s.t. } X_1 \models \gamma' \land X_2 \models \neg \gamma'$.

but from (dev.i) these assumptions imply $X \models \gamma', \neg \gamma'$ which is the required contradiction. $\square$

Reconciling C$_3$. As a concept, behavioural compatibility is extremely general; the notion is, firstly, FDT dependent and, secondly, can be interpreted a number of ways for each FDT, thus, a direct relating of C$_3$ and C is not possible. However, we can give strong evidence that C$_3$ can be fully embraced. In particular, the following results give a general relationship for implementation complete formal techniques.

**Proposition 15**

For implementation complete languages and $\omega$ an equivalence $C_{\omega}^\omega = C_\omega$.

**Proof**

Directly from 2. $\square$

Thus, if ft is implementation complete and behavioural compatibility induces an equivalence on C$_3$ we can make a straightforward instantiation of behavioural compatibility in the development relation and obtain an equivalent definition. Furthermore, the restriction to implementation complete formal techniques is not overly restrictive, since the target of C$_3$ is the behavioural portion of notations such as, LOTOS, Estelle and SDL, which can be viewed to be inherently implementation complete.

We will further justify that C$_3$ can be embraced by C by showing, in section 4.2.2, that all the obvious LOTOS instantiations of behavioural compatibility can be given an equivalent C interpretation. This is strong evidence since LOTOS is a main target for the behavioural compatibility concept. We will summarise these results here.

Firstly, using proposition 15 above we can reconcile any LOTOS instantiation that interprets C$_3$ as an equivalence, e.g. testing equivalence or weak or strong bisimulation. In addition, we will show that the single remaining interpretation can also be embraced. Under this interpretation behavioural compatibility is viewed as the LOTOS $\text{conf}$ relation, which is a realistic interpretation of conformance. By using a relation based on $\text{conf}$, denoted $\text{xc}$, as development relation we can get the required relationship between C$_3$ and C.

**Proposition 16** For LOTOS specifications and be = $\text{conf}$, $C_3 = C_{\text{xc}}$.

**Proof**

See section 4.2.2.

We will explain the relation $\text{xc}$ and prove this result in section 4.2.2.

**Aspects of Generality**

The previous subsection has indicated that the three RM-ODP definitions of consistency can be embraced by C. This suggests that our interpretation of consistency is general relative to the RM-ODP definitions. It is also worth highlighting the particular aspects of our interpretation that make it general:-

- Different development relations can be instantiated which are appropriate both to different FDTs and to assessing different forms of consistency.
• Notions of internal validity relevant to different languages can be employed.

• Both intra and inter language consistency are incorporated.

• Consistency checking between an arbitrary number of descriptions can be supported and checked according to a list of different development relations. Binary consistency is just a special case of this global consistency.

• Both balanced and unbalanced consistency are incorporated.

• Both logical and behavioural notions of consistency are embraced.

3.5.2 Discussion

We will ultimately need global consistency, however, there are a number of outstanding issues to consider on this topic. In particular, we need to determine which unification to choose when doing incremental consistency checking in order that we can obtain global consistency between a group of n (larger than 2) specifications. Taking representative and least unifications is clearly relevant to this issue, but more work is required in order to fully characterise the problem. Thus, in the remainder of this report we will concentrate on consistency between pairs of specifications.
CHAPTER 3. DEFINITIONS OF CONSISTENCY
Chapter 4

Consistency in LOTOS

In this chapter we present an overview of the work done on consistency checking and unification within the formal specification language LOTOS. The approach to consistency checking and unification outlined in this chapter is general: ODP correspondence rules or the object based nature of ODP specifications, for example, are not yet taken into account. As a consequence, the unification methods given in this chapter also apply to composition of partial specifications beyond the scope of Open Distributed Processing.

LOTOS is a process algebra based specification language which is used for the formal specification of distributed, concurrent information processing systems (see [6] for a general introduction). In particular, LOTOS [27] was adopted by the International Standardisation Organisation (ISO) to formally describe the services and protocols of the Open Systems Interconnection Reference Model (OSI/RM) [26, 29, 28]. Currently, LOTOS is also being used for the specification of ODP systems and standards.

The LOTOS language has two parts: a behavioural part and a data part. The former of these is a process algebraic language, related to CCS, CSP and ACP, in which systems are described in terms of the temporal relationship between externally observable actions. The subset of LOTOS that consists solely of the behavioural part is usually referred to as Basic LOTOS. The latter is an abstract data typing language, ACT-ONE. In this chapter, we will largely focus on Basic LOTOS, i.e. the behavioural part. Basic LOTOS is a natural point of focus since it is the subset of LOTOS that is most fundamentally different to Z and ultimately we are interested in addressing the hardest consistency checking between Z and LOTOS, which will arise in this circumstance. Nevertheless, since the semantics of a Full LOTOS specifications can be expressed in Basic LOTOS, the consistency checking mechanisms developed here for Basic LOTOS apply equally well to Full LOTOS behaviour specifications.

Structure of Chapter. The general framework for consistency checking outlined in chapter 3 relies on the existence of development relations for the formal techniques used. In section 4.1, some of the existing development relations for LOTOS are briefly reviewed. This is followed in section 4.2 by an investigation of the generality of our definition of consistency, definition 9, in the LOTOS setting. The three RM-ODP consistency definitions for LOTOS are interpreted and then related to the general definition of consistency given in chapter 3. Section 4.3 considers part of the spectrum of possible LOTOS instantiations of definition 9 and relates these instantiations in both the unbalanced and balanced case. This is followed in section 4.4 by a presentation of some specific consistency checking techniques for balanced consistency. Operational semantics for a number of classes of unification are presented. Some of the developed techniques are then applied to two example viewpoint specifications of a shared memory system in section 4.5. Finally, we present some concluding remarks in section 4.6.
CHAPTER 4. CONSISTENCY IN LOTOS

4.1 Development Relations

In this section, some well-known development relations for LOTOS are recapitulated. We assume the reader has some familiarity with LOTOS and its semantics. Before definitions of these relations are given, we introduce some notation that will allow us to reason about processes.

**Notation.** LOTOS has a well-defined operational semantics which maps LOTOS behaviour expressions onto Labelled Transition Systems (LTSs). As a result of the existence of such a mapping, and the possibility to express any LTS in LOTOS, we can use behaviour expressions and their corresponding LTSs interchangeably. In particular, relations defined on transition systems are likewise applicable to behaviour expressions and the processes they define.

A labelled transition system is a tuple, \( \langle S, \mathcal{L}, T, s_0 \rangle \). \( S \) is a set of states which ranges over the possible process behaviours that the system can reach; \( \mathcal{L} \) is a set of action labels; \( T \) is a set of transitions of the form \( P \xrightarrow{a} P' \); and \( s_0 \) is a starting state.

In the following \( P, P', Q, Q' \) stand for processes. \( \mathcal{L} \) is the alphabet of observable actions associated with a certain process, while \( i \) is the invisible or internal action. We use the variables \( a, a_i \) to range over \( \mathcal{L} \), and the variables \( \mu, \mu_i \) to range over \( \mathcal{L} \cup \{i\} \). Furthermore, \( \mathcal{L}^* \) denotes strings over \( \mathcal{L} \). The constant \( e \in \mathcal{L}^* \) denotes the empty string, and the variables \( \sigma, \sigma_i \) are used to range over \( \mathcal{L}^* \). Elements of \( \mathcal{L}^* \) are also called traces. In table 4.1 the notion of transition is generalised to traces. The definitions of \( \xrightarrow{a} \) and \( \overset{\sigma}{\rightarrow} \) are inductive on the length of \( \sigma \).

Using the notation derived in table 4.1, we can now define some other useful concepts:

- \( \text{Tr}(P) = \{ \sigma \mid P \xrightarrow{\sigma} \} \), denotes the set of traces of a process \( P \).
- \( \text{out}(P, \sigma) = \{ a \mid \sigma a \in \text{Tr}(P) \} \), denotes the set of possible observable actions after the trace \( \sigma \).
- \( P \text{ after } \sigma = \{ P' \mid P \xrightarrow{\sigma} P' \} \), denotes the set of all states reachable from \( P \) by the trace \( \sigma \).
- \( P \text{ after } \sigma = \{ P' \mid P \xrightarrow{\sigma} P' \} \subseteq P \text{ after } \sigma \), denotes the set of states that \( P \) leads to under \( \sigma \neq e \).
- \( \text{Ref}(P, \sigma) = \{ X \mid \exists P' \in (P \text{ after } \sigma), \text{ such that } \forall a \in X : P' \xrightarrow{\sigma} \} \), denotes the refusal set of \( P \) after the trace \( \sigma \).

### 4.1.1 Trace preorder

An important category of system properties that one would like have satisfied by a LOTOS specification, are safety properties. Safety properties state that something bad should not happen, where something bad can be interpreted as a certain trace of the specification. Observe that if \( S \) is a safety property, then \( \forall \sigma_1, \sigma_2 \) we have if \( \sigma_1 \leq \sigma_2 \) then \( S(\sigma_1) \Rightarrow S(\sigma_2) \), i.e. if \( S \) holds for the trace \( \sigma_2 \), it also holds for all its prefixes. In particular, all safety properties hold for the empty trace \( e \).

<table>
<thead>
<tr>
<th>Notation</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>( P \xrightarrow{a} P' )</td>
<td>denotes a transition, i.e. ( P ) can do ( a ) and consequently behaves as ( P' ).</td>
</tr>
<tr>
<td>( P \xrightarrow{\sigma} )</td>
<td>( \exists P' ) such that ( P \xrightarrow{\sigma} P' )</td>
</tr>
<tr>
<td>( P \overset{i}{\rightarrow} )</td>
<td>( \exists P' ) such that ( P \overset{i}{\rightarrow} P' )</td>
</tr>
<tr>
<td>( \xrightarrow{a} )</td>
<td>reflexive and transitive closure of ( \xrightarrow{a} )</td>
</tr>
<tr>
<td>( P \xrightarrow{\sigma} P' )</td>
<td>( \exists Q, Q' \cdot P \xrightarrow{\sigma} Q \overset{\sigma}{\rightarrow} Q' \xrightarrow{\sigma'} P' )</td>
</tr>
<tr>
<td>( P \xrightarrow{\sigma} P' )</td>
<td>( \exists P' \cdot P \xrightarrow{\sigma} P' )</td>
</tr>
<tr>
<td>( P \overset{\alpha}{\rightarrow} P' )</td>
<td>( \exists P' \cdot P \overset{\alpha}{\rightarrow} P' )</td>
</tr>
</tbody>
</table>

Table 4.1: Derived transition denotations
When a specification is refined, it seems reasonable to require that the refinement is at least as safe as the specification. This intuition is reflected by the trace preorder.

**Definition 20 (trace preorder)**

Given two process specifications $P$ and $Q$, then $P$ refines $Q$ by reducing the possible traces, denoted $P \leq_{tr} Q$, iff:

- $\text{Tr}(P) \subseteq \text{Tr}(Q)$, or equivalently
- $\forall \sigma \in \mathcal{L}^* \cdot P \xrightarrow{\sigma} \Rightarrow Q \xrightarrow{\sigma}$.

### 4.1.2 Conformance

In addition to safety properties we are sometimes also interested in the liveness (or deadlock) properties of a system specification. A liveness property states that something good must eventually happen. It may be required that a development of a specification does not deadlock in a situation where the specification would not deadlock, in other words, every trace that the specification *must* do, the development *must* do as well. This requirement is formalised by the $\text{conf}$ relation [11] [12], which has been adopted as the primary interpretation of conformance in LOTOS.

**Definition 21 (conformance)**

Given two process specifications $P$ and $Q$, then $P$ conforms to $Q$, denoted $P \text{ conf } Q$, iff:

- $\forall \sigma \in \text{Tr}(Q)$ and $\forall A \subseteq \mathcal{L}$ we have
  - if $\exists P' \in (P \text{ after } \sigma)$ such that $\forall \alpha \in A \cdot P' \xrightarrow{\alpha}$,
  - then $\exists Q' \in (Q \text{ after } \sigma)$ such that $\forall \alpha \in A \cdot Q' \xrightarrow{\alpha}$, or equivalently
- $\forall \sigma \in \text{Tr}(Q) \cdot \text{Ref}(P, \sigma) \subseteq \text{Ref}(Q, \sigma)$

We will also use a development relation which is a symmetric subset of $\text{conf}$. This relation is called $\text{conf symmetric}$ and is denoted $\text{cs}$; it will play a central role in instantiations of $C_3$, the RM-ODP definition of consistency based on behavioural compatibility. In particular, since the ODP architectural semantics adopt $\text{conf}$ as their interpretation of behavioural compatibility $\text{cs}$ is the obvious interpretation of $C_3$.

**Definition 22 (conf symmetric)**

Given two process specifications $P$ and $Q$, then $P \text{ cs } Q$ iff $P \text{ conf } Q \land Q \text{ conf } P$.

### 4.1.3 Reduction

A refinement relation that combines both the preservation of safety and liveness properties is the reduction relation, $\text{red}$, defined in [11].

**Definition 23 (reduction)**

Given two process specifications $P$ and $Q$, then $P$ (deterministically) reduces $Q$, denoted $P \text{ red } Q$, iff:

1. $P \leq_{tr} Q$, and
2. $P \text{ conf } Q$
4.1.4 Extension

Another refinement relation proposed in [11] is the extension relation. This relation allows for the introduction of new possible traces in a refinement, while preserving the liveness properties of the specification. Extension seems particularly relevant in the context of partial specification.

**Definition 24 (extension)**

Given two process specifications \( P \) and \( Q \), then \( P \) extends \( Q \), denoted \( P \text{ ext} \ Q \), iff:

1. \( \text{Tr}(P) \supseteq \text{Tr}(Q) \), and
2. \( P \text{ conf} Q \)

4.1.5 Structural Refinement - Testing Equivalence

Another view of refinement is that the refinement provides more detail on the subdivision of the system into smaller components than the specification. The specification and its refinement are semantically equivalent, i.e. they express the same external or observable behaviour. The intension of both descriptions is not the same though, as the refinement gives more detail about the internal structure of the system under consideration. The weakest interpretation of observational equivalence is captured by the testing equivalence relation.

**Definition 25 (testing equivalence)**

Given two process specifications \( P \) and \( Q \), then \( P \) is testing equivalent to \( Q \), denoted \( P \text{ te} Q \), iff:

- \( P \text{ red} Q \) and \( Q \text{ red} P \), or equivalently
- \( P \text{ ext} Q \) and \( Q \text{ ext} P \), or equivalently
- \( \text{Tr}(P) = \text{Tr}(Q) \land \forall \sigma \in \text{Tr}(P) \cdot \text{Ref}(P, \sigma) = \text{Ref}(Q, \sigma) \).

4.1.6 Structural Refinement - Bisimulation Equivalences

An alternative interpretation of observational identity is given by the bisimulation equivalences, strong and weak bisimulation [40]. The observational identity that they induce is in some circumstances seen to be too strong [35], for example, when the observer is viewed as a tester for the process.

The definition of weak bisimulation equivalence, \( \approx \), of LOTOS processes is given below.

**Definition 26 (weak bisimulation)**

Given two LOTOS process specifications \( P_1 \) and \( P_2 \), we define \( P_1 \approx P_2 \iff \forall \sigma \in \mathcal{L}^* \)

1. if \( \exists P'_1 \cdot P_1 \xrightarrow{\sigma} P'_1 \) then \( \exists P'_2 \cdot P_2 \xrightarrow{\sigma} P'_2 \) and \( P'_1 \approx P'_2 \); and
2. if \( \exists P'_2 \cdot P_2 \xrightarrow{\sigma} P'_2 \) then \( \exists P'_1 \cdot P_1 \xrightarrow{\sigma} P'_1 \) and \( P'_1 \approx P'_2 \).

Strong bisimulation, denoted \( \sim \), is defined in a similar manner to weak bisimulation, except \( i \) actions are matched in addition to observable actions. Hence strong bisimulation is an even stronger notion of observational identity than weak bisimulation.
4.2. RELATING THE RM-ODP DEFINITIONS

4.1.7 Discussion: Properties of the Development Relations

Apart from cs the properties of the development relations presented above have been well documented in the literature. We will review some of these properties here.

**Proposition 17**

(i) $\leq_{tr}$, red and ext are preorders.
(ii) te, $\sim$, and $\approx$ are equivalences.
(iii) conf is reflexive, but neither symmetric or transitive.
(iv) cs is (a) reflexive and (b) symmetric, but (c) not transitive.

**Proof**

(i), (ii) and (iii) are all standard results from the theory of LOTOS and process algebra in general, see for instance [35], [24] and [40]. However, (iv) needs some justification:

(iva) This is a consequence of conf being reflexive.

(ivb) This is immediate from the definition of cs.

(ivc) The following counterexample justifies this. Let $P := b; \text{stop}[]i; a; \text{stop}; Q := i; a; \text{stop}$ and $R := b; c; \text{stop}$; then $P \text{ cs } Q, Q \text{ cs } R$, but $\neg(P \text{ cs } R)$. This is because $\neg(P \text{ conf } R)$ as $P$ refuses $c$ after the trace $b$, but $R$ cannot refuse $c$ after the same trace. \qed

Thus, $\leq_{tr}$, red and ext can be classed together as preorder refinement relations; te, $\sim$, and $\approx$ can be classed together as equivalences; conf and cs are weaker implementation relations.

4.2 Relating the RM-ODP Definitions

Section 3.5.1 has related the three RM-ODP definitions of consistency to $C$ in a very general way, in this section we will specialize this to LOTOS instantiations of the RM-ODP definitions. This will give us even more evidence that $C$ is fully general. This is particularly an issue since $C_3$ is dependent upon the FDT specific interpretation of behavioural compatibility adopted and thus only a language specific relating of $C_3$ and $C$ can be given.

The section begins by giving a set of LOTOS instantiations of relevant definitions and, in particular, the RM-ODP definitions, and then these instantiations are related to $C$ in the following subsection.

4.2.1 RM-ODP Instantiations

Firstly, we have the following:-

**Proposition 18**

$\forall P \in DES_{LOTOS}, \Psi(P)$.

This follows intuitively from considering the nature of LOTOS specifications. In particular, at least theoretically, we can view all LOTOS specifications as implementable. Even degenerate specifications, such as those containing deadlocks, for example, have a physical implementation equivalent. This is a fundamental characteristic of behavioural languages that distinguishes them from logically based specification notations.

**Corollary 4**

LOTOS is implementation complete and all unifications are representative.
This corollary is important as it considerably simplifies the class of consistency that must be considered for LOTOS. Furthermore, we assume that all consistency checks are type correct. This is reasonable since we are only considering consistency intra the LOTOS language.

Of the specific RM-ODP definitions, we could relate $C_1$ via an interpretation of LOTOS in logic, for example, the temporal logic interpretation in [20], however, this is a complex interpretation with a number of subtle issues. Thus, we will view this as beyond the immediate scope of our work and we will not consider $C_1$ further in the context of LOTOS. In contrast, $C_{2.2}$ and $C_3$ are immediately appropriate to LOTOS. We will consider these in turn.

**Instantiation of $C_{2.2}$**. This is very straightforward, we give the following definition:

**Definition 27** For $P_1, P_2, P \in DES_{LOTOS}$ $P_1 \sim C_{2.2} P_2$ iff $\exists P \cdot P \text{conf} P_1 \land P \text{conf} P_2$.

Although it should be noted that this instantiation is dependent on the interpretation of conformance adopted. conf is a weak interpretation, in particular, it does not enforce the preservation of safety properties. However, conf is a realistic reflection of the capabilities of conformance testing and is now accepted as the basis of work on test case generation for LOTOS [18] [12].

**Instantiation of $C_3$**. Consistency definition $C_3$ is dependent upon the interpretation of behavioural compatibility, which in turn hinges on the interpretation of a specification’s environment and the criteria imposed on that environment. The looseness of the definition of behavioural compatibility implies that one of a number of interpretations of $C_3$ could be made. It is our view that $C_3$ could be interpreted as any of the following:

**Definition 28**

(i) $P_1 C_3 L P_2$ iff $P_1 \sim P_2$ - Strong Bisimulation

(ii) $P_1 C_3 L P_2$ iff $P_1 \approx P_2$ - Weak Bisimulation

(iii) $P_1 C_3 L P_2$ iff $P_1 \text{te} P_2$ - Testing Equivalence

(iv) $P_1 C_3 L P_2$ iff $P_1 \text{cs} P_2$ - Conf symmetric

Definitions 28(i) and 28(ii) view the environment as an unconstrained observer, in the sense of bisimulation equivalences. Note definition 28(i) is particularly interesting because $C_3 L$ has the advantage of being a congruence for LOTOS.

In contrast, 28(iii) and 28(iv) view the environment as a tester for the specifications. The distinction between 28(iii) and 28(iv) is that 28(iii) implies robustness testing and 28(iv) implies restricted testing, see [18] [12] for a discussion of these alternatives. Amongst these definitions $C_3 L$ is particularly important for a number of reasons. Firstly, this interpretation agrees with the LOTOS definition of behavioural compatibility in the RM-ODP architectural semantics [30]. In addition, as indicated in the following proposition, $C_3 L$ is the weakest of the LOTOS interpretations of $C_3$.

**Proposition 19**

$C_3 L \subset C_3^{te} \subset C_3^{cs} \subset C_3^{cs}$.

**Proof**

$C_3 L \subset C_3^{te} \subset C_3^{cs}$ are standard process algebra results. $C_3^{te} \subset C_3^{cs}$ requires some justification. Firstly, it is straightforward to see that $te \subseteq cs$. In addition, we can provide the two processes $P := a; \text{stop}[i]; b; \text{stop}$ and $Q := i; b; \text{stop}$ as counterexamples to justify that $cs \not\subseteq te$, since $P cs Q$, but $\lnot(P \text{te} Q)$ as the trace sets of the two processes are not equal.

Furthermore, [9] has shown that $C_3$ is the strongest of the RM-ODP interpretations of consistency, thus, $C_3^{cs}$ bounds the relationship between $C_3$ and the other RM-ODP consistency definitions and warrants particular attention.
4.2.2 Relating Definitions

This subsection specializes the results of section 3.5.1 to LOTOS.

Reconciling $C_{2.2}$. This interpretation of consistency is very straightforward:

**Proposition 20**

For LOTOS $C_{2.2} = C_{\text{conf}}$.

**Proof**

Immediate from a comparison of instantiations.

Reconciling $C_3$. Three of the interpretations made in section 4.2.1 can be related to our general definition easily.

**Corollary 5**

(i) $C_3^{\sim} = C_{\sim}$

(ii) $C_3^{=} = C_{=}$

(iii) $C_3^{\text{te}} = C_{\text{te}}$

**Proof**

Immediate from corollary 2.

Thus, interpretations of behavioural compatibility in LOTOS which are based on one of the language’s equivalences are easily reflected in our general definition of consistency. But, $C_3^{=} \not\subseteq C_3^{\sim}$ is not transitive, proposition 17, so corollary 2 does not get us a relationship between $C_3^{=} \not\subseteq C_3^{\sim}$ and $C_3^{\text{te}}$.

In fact, we have the following result.

**Proposition 21**

$C_3^{=} \subseteq C_3^{\text{cs}}$.

**Proof**

Firstly, $P_1 \xrightarrow{C_3^{=}} P_2 \implies P_1 C_3^{\text{cs}} P_2$, follows immediately from the reflexivity of $C_3^{\text{cs}}$, i.e. either of $P_1$ or $P_2$ could act as the required common $C_3^{=} - \text{development}$.

In addition, we can provide a counterexample to show that, $C_3^{=}$ does not get us a relationship between $C_3^{=} \not\subseteq C_3^{\sim}$ and $C_3^{\text{te}}$. Consider, $P_1 := i; a; \text{stop}$ and $P_2 := i; a; \text{stop}$ and $P := i; a; \text{stop}$. Now, $P \xrightarrow{C_3^{=}} P_1$ and $P \xrightarrow{C_3^{\text{cs}}}$ $P_2$, but $\neg(P_1 \xrightarrow{C_3^{=} \text{cs}} P_2)$. This is because $\neg(P_2 \xrightarrow{C_3^{=} \text{conf}} P_1)$ as $P_2$ refuses $c$ after the trace $b$, but $P_1$ cannot refuse $c$ after the same trace.

This result is disappointing, but interesting. The counterexample provided is one of the few situations in which the unification has a smaller trace set than both the original specifications and furthermore a unification with a larger trace set does not seem to exist for this example. This observation motivates the following, which considers a development relation in which the trace set increases. Thus, we define extended $C_3^{=} \text{symmetric}$, denoted $C_3^{\text{cs}}$ as:

**Definition 29**

$P_1 \xrightarrow{C_3^{\text{cs}}} P_2$ iff $P_1 \xrightarrow{C_3^{=}} P_2 \wedge \text{Tr}(P_1) \supseteq \text{Tr}(P_2)$.

It should be clear that an alternative derivation of $C_3^{\text{cs}}$ is: $P_1 \xrightarrow{C_3^{\text{cs}}} P_2$ iff $P_1 \xrightarrow{C_3^{=}} P_2 \wedge P_2 \text{conf} P_1$. So, we have added a trace extension constraint on the development. Note in particular that using $C_3^{\text{cs}}$ as development relation in $C$ will rule out the counterexample used in the previous proposition.

So let us try to relate $C_3^{\text{cs}}$ and $C_3^{=}$.  

**Proposition 22**

$C_3^{\text{cs}} \subseteq C_3^{=}$

**Proof**

Assume $P_1 \xrightarrow{C_3^{=}} P_2$, i.e., $\exists P : P \text{ conf} P_2 \wedge P \text{ conf} P_1 \wedge \text{Tr}(P) \supseteq \text{Tr}(P_2)$ and $P \text{ conf} P_1 \wedge P_1 \text{ conf} P_2 \wedge \text{Tr}(P) \supseteq \text{Tr}(P_1)$.
which expands to:-

(i) \( \forall \sigma \in Tr(P_2), \ Ref(P, \sigma) \subseteq Ref(P_2, \sigma) \land \)
(ii) \( \forall \sigma' \in Tr(P), \ Ref(P_2, \sigma') \subseteq Ref(P, \sigma') \land \)
(iii) \( \forall \sigma'' \in Tr(P_1), \ Ref(P, \sigma'') \subseteq Ref(P_1, \sigma'') \land \)
(iv) \( \forall \sigma' \in Tr(P), \ Ref(P_1, \sigma') \subseteq Ref(P, \sigma') \land \)
(v) \( \forall \sigma \in Tr(P) \supseteq Tr(P_1), Tr(P_2) \)

From properties (i), (iv) and (v) we get:-

\( \forall \sigma_1 \in Tr(P_2), \ Ref(P_1, \sigma_1) \subseteq Ref(P_1, \sigma) \leq Ref(P_2, \sigma_1) \)

i.e. \( P_1 \ conf P_2 \). Similarly, \( P_2 \ conf P_1 \) since properties (iii), (ii) and (v) give us:

\( \forall \sigma_2 \in Tr(P_1), \ Ref(P_2, \sigma_2) \subseteq Ref(P_1, \sigma_2) \leq Ref(P_2, \sigma_1) \)

Notice these relationships can only be derived because \( Tr(P) \supseteq Tr(P_1), Tr(P_2) \).

So, we have the direction of implication that we could not get with \( C_{cs} \), but now the other implication direction is more difficult as we need to show a unification with trace extension exists. Before we consider this we need a simple result.

**Proposition 23**

\( P_1 \ conf P_2 \implies \forall \sigma \in Tr(P_1) \cap Tr(P_2), \ Ref(P_1, \sigma) = Ref(P_2, \sigma) \).

**Proof**

A straightforward consequence of the definition of \( cs \).

We will use the following unification construction:-

Denote \( U_{cs}(P_1, P_2) \) as the set of all LOTOS specifications characterised by the following constraints:-

\[
\begin{align*}
Tr(U_{cs}(P_1, P_2)) &= Tr(P_1) \cup Tr(P_2) \quad & \text{(a)} \\
\forall \sigma \in Tr(U_{cs}(P_1, P_2)), &
\begin{align*}
\sigma \in Tr(P_1) \cap Tr(P_2) &\implies Ref(U_{cs}(P_1, P_2), \sigma) = Ref(P_1, \sigma) = Ref(P_2, \sigma) \quad & \text{(b)} \\
\sigma \in Tr(P_1) \setminus Tr(P_2) &\implies Ref(U_{cs}(P_1, P_2), \sigma) = Ref(P_1, \sigma) \quad & \text{(c)} \\
\sigma \in Tr(P_2) \setminus Tr(P_1) &\implies Ref(U_{cs}(P_1, P_2), \sigma) = Ref(P_2, \sigma) \quad & \text{(d)}
\end{align*}
\end{align*}
\]

Notice that (b) is only possible because of proposition 23. It should also be noted that this construction is well founded and will always yield a LOTOS specification. One justification for this is that [35] performs the same construction with his rooted failure tree model (definition 6.4.1 on page 153) and shows that the resulting tree is well-formed, i.e. can be mapped to a labelled transition system.

**Proposition 24**

\( C_{cs}^s \subseteq C_{cs} \).

**Proof**

Assume \( P_1 \ conf P_2, \) i.e. \( P_1 \ conf P_2 \), then take \( X \in U_{cs}(P_1, P_2) \), we suggest that \( X \) is a common \( xcs \) development of \( P_1 \) and \( P_2 \), as required by \( C_{cs} \). Let us show that \( X \ xcs P_1 \). We will show first that \( X \ conf P_1 \), then that \( P_1 \ conf X \) and then that \( Tr(X) \supseteq Tr(P_1) \).

(i) \( (X \ conf P_1) \).

Take \( \sigma \in Tr(P_1) \). Now, if \( \sigma \) is also a trace of \( P_2 \), by (b), \( Ref(X, \sigma) = Ref(P_1, \sigma) \), however, if \( \sigma \notin Tr(P_2) \), by (c), \( Ref(X, \sigma) = Ref(P_1, \sigma) \).

(ii) \( (P_1 \ conf X) \).

Take \( \sigma \in Tr(X) \). We have the following cases:-

(a) \( \sigma \in Tr(P_1) \cap Tr(P_2) \implies Ref(P_1, \sigma) = Ref(X, \sigma) \), by (b).

(b) \( \sigma \in Tr(P_1) \setminus Tr(P_2) \implies Ref(P_1, \sigma) = Ref(X, \sigma) \), by (c).

(c) \( \sigma \in Tr(P_2) \setminus Tr(P_1) \implies Ref(P_1, \sigma) = \emptyset \) as \( \sigma \notin Tr(P_1) \), hence \( Ref(P_1, \sigma) \subseteq Ref(X, \sigma) \).

(iii) \( (Tr(X) \supseteq Tr(P_1)) \).

This is immediate from (a).
Thus, $X \xcs P_1$ and it can be similarly verified that $X \xcs P_2$.

Corollary 6

$C_{xcs}^a = C_{xcs}$.

This result completes our relating of $C_3$ to $C$ and shows that all, obvious LOTOS instantiations of behavioural compatibility in $C_3$ can be given an equivalent formulation in $C$ and justifies proposition 16 quoted in chapter 3.

4.3 General LOTOS Instantiations of Consistency

The previous section has shown how $C$ is general with regard to LOTOS interpretations of the RM-ODP definitions. In this section we will consider the broad properties of LOTOS instantiations of $C$. We will provide categorisations of a number of the definitions. The section is divided into two subsections, the first considers unbalanced consistency and the second balanced consistency.

4.3.1 Unbalanced Consistency

The main motivation for considering unbalanced consistency is to enable us to address situations in which a viewpoint is a direct development of a second viewpoint, or in which the viewpoint specifications have a different level of granularity. Thus, we would like to give LOTOS instantiations of consistency that model the LOTOS development relations, $conf$, $red$, $ext$, etc. Firstly, it is clear that notions of development based on equivalence, e.g. $\approx$, $\sim$ and $te$, can be easily embraced. In addition, development using one of the LOTOS preorders can be easily embraced as corollary 3 suggests. We have the following instantiations of this result for LOTOS preorders:

**Proposition 25**

(i) $\leq_{tr} = C_{\leq_{tr}^{-1}} = C_{\leq_{tr}^{-1} \cap \leq_{tr} : \leq_{tr}}$.
(ii) $red = C_{red^{-1}, \red} = C_{te, red}$.
(iii) $ext = C_{ext^{-1}, \ext} = C_{te, ext}$.

**Proof**

Follows immediately from corollary 3.

However, since $conf$ is not transitive we have to work a bit harder to relate this notion of development. Firstly, we note the following negative result:

**Proposition 26**

$conf \not\supseteq C_{conf^{-1}, conf}$

**Proof**

We can use our $conf$ transitivity counterexample again here. Specifically, let $P_1 := k; \ stop[i]; a; \ stop$, $P_2 := k; c; \ stop[i]; a; \ stop$ and $P := i; a; \ stop$ then, $P$ is the required common development to give $P_1 C_{conf^{-1}, conf} P_2$, but $\neg(P_1 conf P_2)$.

However, the following stronger result gives us the necessary relationship:

**Proposition 27**

$conf = C_{te, conf}$.

**Proof**

$\ (conf \subseteq C_{te, conf})$

Assume $P_1 conf P_2$, but in addition from reflexivity of $te$, $P_1 te P_1$ and, thus, $P_1$ is the required common development.
Take \( P \) such that \( P \cong P_1 \) and \( P \conf P_2 \). If we expand these out we get:

\[
\begin{align*}
\Tr(P) &= \Tr(P_1) \land \forall \sigma \in \Tr(P), \Ref(P, \sigma) = \Ref(P_1, \sigma) \land \\
\forall \sigma \in \Tr(P_2), \Ref(P, \sigma) &\subseteq \Ref(P_2, \sigma)
\end{align*}
\]

Equality of the traces of \( P_1 \) and \( P \) implies that there are no traces of \( P_2 \) that \( P_1 \) could do, but \( P \) could not do, thus,

\[
\forall \sigma \in \Tr(P_2), \Ref(P, \sigma) = \Ref(P_1, \sigma) \text{ and thus, } \forall \sigma \in \Tr(P_2), \Ref(P_1, \sigma) \subseteq \Ref(P_2, \sigma).
\]

So, \( P_1 \conf P_2 \) as required.

This completes our relating of LOTOS development relations to \( C \).

### 4.3.2 Balanced Consistency

In this section balanced consistency is instantiated with the LOTOS development relations. This class of consistency is the easiest to work with and thus we will be able to obtain a complete categorisation of the relationship between the different instantiations. Our presentation works from the weakest interpretations of consistency to the strongest.

We begin with a surprising result:

**Proposition 28**

Take \( P_1, P_2 \in DES_{\text{LOTOS}} \), then:

(i) \( P_1 \conf P_2 \equiv \text{true} \).

(ii) \( P_1 \ext P_2 \equiv \text{true} \).

(iii) \( P_1 \conf P_2 \equiv \text{true} \).

**Proof**

We justify these results in turn:

(i) For any two processes \( P \) and \( Q \), we have \( \text{stop} \leq_{tr} P \) and \( \text{stop} \leq_{tr} Q \), i.e. the empty process is always a common development of any two specifications.

(ii) In [32] it is shown that for any two LTSs a third LTS can be found that is an extension of both. Since LTSs provide the semantics for processes, this result extends to LOTOS specifications.

(iii) Since \( \ext \implies \conf \) this is an easy consequence of (ii).

**Corollary 7**

\[
\conf = \ext = \conf = \text{true}
\]

The implication of these results is that all pairs of LOTOS specifications will be found to be consistent by \( \conf \), \( \ext \) and \( \conf \). Thus, these instantiations of consistency are very weak and are unable to distinguish any specifications. In other words, when \( \leq_{tr} \), \( \ext \) or \( \conf \) is the appropriate development relation, there is no need for a consistency check.

**Example 1** We illustrate the second case, (ii), of proposition 28. \( \ext \) supports functionality extension and any pair of LOTOS specifications can be reconciled according to such extension. Figure 4.1 illustrates this fact with some examples. The following properties hold:

\[
\begin{align*}
P &\in U[\ext](P_1, P_2) \\
Q &\in U[\ext](Q_1, Q_2) \\
Q' &\notin U[\ext](Q_1, Q_2) \\
R &\in U[\ext](R_1, R_2) \\
R' &\notin U[\ext](R_1, R_2)
\end{align*}
\]

Notice that (b) shows that \( U[\ext] \) must not introduce new non-determinism, e.g. \( Q \) is a unification, but \( Q' \) is not as it may refuse either \( d \) or \( e \) after performing \( a \) and \( Q_1 \) cannot refuse \( d \) after \( a \) and \( Q_2 \) cannot refuse \( e \) after \( a \). Additionally, (c) shows that unification may limit non-determinism. Specifically, \( R \) is a unification, but \( R' \) is not as it can refuse everything after the empty trace, while \( R_2 \) must offer either \( a \) or \( c \).
However, other instantiations of consistency are distinguishing. We begin with the following result:

**Proposition 29**  
\( C_{\text{red}} \subset \text{true}. \)

**Proof**  
We simply need to exhibit a pair of specifications that are not consistent. Consider \( P_1 := a; \text{stop} \) and \( P_2 := b; \text{stop} \). The only process which is a trace subset of both \( P_1 \) and \( P_2 \) is \( P := \text{stop} \), but \( P \) cannot be a common reduction of either \( P_1 \) or \( P_2 \) since it will refuse \( a \) and \( b \) immediately. \( \square \)

Thus, \( C_{\text{red}} \) is stronger than \( C_{\leq \text{tr}} \), \( C_{\text{ext}} \) and \( C_{\text{conf}} \).

**Example 2** Consider the examples in figure 4.2. The following properties hold:-  
\( U[\text{red}](P_1, P_2) = \emptyset \)  
\( Q \in U[\text{red}](Q_1, Q_2) \)  
\( R_2 \in U[\text{red}](R_1, R_2) \)

The consistency relation \( C_{\text{cs}} \) was introduced in section 4.2.2 in an attempt to reconcile \( C_{3}^{\text{cs}} \) with \( C \), our results there showed that \( C_{\text{cs}} \) is weaker than \( C_{3}^{\text{cs}} \), however, we are interested to determine how much weaker. If we can show that \( C_{\text{cs}} \) is stronger than \( C_{\text{red}} \) then we will have an upper and lower bound on the strength of \( C_{\text{cs}} \). Thus, we will consider the relationship between \( C_{\text{cs}} \) and \( C_{\text{red}} \). We will use the following result:-  

**Lemma 2**  
\( P_1 \ C_{\text{cs}} \ P_2 \implies \forall P \in U[\text{cs}](P_1, P_2), \forall \sigma \in \text{Tr}(P) \cap \text{Tr}(P_1) \cap \text{Tr}(P_2), \text{Ref}(P, \sigma) = \text{Ref}(P_1, \sigma) = \text{Ref}(P_2, \sigma). \)

**Proof**  
Assume \( \sigma \in \text{Tr}(P) \cap \text{Tr}(P_1) \cap \text{Tr}(P_2) \) then from \( P \ C_{\text{cs}} P_1 \) we apply proposition 23 to get \( \text{Ref}(P, \sigma) = \text{Ref}(P_1, \sigma) \) and from \( P \ C_{\text{cs}} P_2 \) we get \( \text{Ref}(P, \sigma) = \text{Ref}(P_2, \sigma) \) and the result follows directly. \( \square \)
Using this result we can obtain the following:

**Proposition 30**

\[ C_{\text{red}} \not\subseteq C_{\text{cs}} \]

**Proof**

We show that \( P_1 \) and \( P_2 \) exist such that \( P_1 \ C_{\text{red}} \ P_2 \), but \( \neg(P_1 \ C_{\text{cs}} \ P_2) \). Consider \( P_1 := a; b; \text{stop} \) and \( P_2 := a; b; \text{stop} \). Now \( P_1 \ C_{\text{red}} P_2 \), because \( P_2 \) can act as the required common reduction.

We argue by contradiction that \( \neg(P_1 \ C_{\text{cs}} P_2) \). So, assume \( P \) is such that \( P \ C_{\text{cs}} P_1 \) and \( P \ C_{\text{cs}} P_2 \). Firstly, \( P \) must be able to perform the trace \( a \), because if \( a \notin \text{Tr}(P) \) then \( \{a\} \not\subseteq \text{Ref}(P, \epsilon) \), but since \( \{a\} \not\subseteq \text{Ref}(P_2, \epsilon) \) this implies that \( \text{Ref}(P, \epsilon) \not\subseteq \text{Ref}(P_2, \epsilon) \) and \( \neg(P \ \text{conf} \ P_2) \).

So, we assume \( a \in \text{Tr}(P) \), hence \( a \in \text{Tr}(P) \cap \text{Tr}(P_1) \cap \text{Tr}(P_2) \) and we can apply lemma 2 which implies that \( \text{Ref}(P, a) = \text{Ref}(P_1, a) = \text{Ref}(P_2, a) \). But this cannot be the case as \( \{b\} \not\subseteq \text{Ref}(P_1, a) \) and \( \{b\} \not\subseteq \text{Ref}(P_2, a) \), so \( \text{Ref}(P_1, a) \neq \text{Ref}(P_2, a) \). Which gives us the required contradiction and implies that such \( P \) does not exist.

So, \( C_{\text{cs}} \) is not weaker than \( C_{\text{red}} \). Using the following small result we will be able to further clarify the relationship between \( C_{\text{cs}} \) and \( C_{\text{red}} \).

**Lemma 3**

\( P \) is fully deterministic (in the usual sense) \( \implies (\forall \sigma \in \text{Tr}(P), a \in \text{out}(P, \sigma) \iff \neg \exists X \in \text{Ref}(P, \sigma) \cdot a \in X) \).

**Proof**

Standard from theory of LOTOS.

This result states that for a deterministic process an action cannot be both offered and refused.

**Proposition 31**

\( C_{\text{cs}} \subseteq C_{\text{red}} \).

**Proof**

Assume \( P_1 \ C_{\text{cs}} P_2 \), i.e. \( \exists P : P \ C_{\text{cs}} P_1 \land P \ C_{\text{cs}} P_2 \). Now construct \( P' \) as the fully deterministic
process characterised by:
\[ Tr(P') = Tr(P) \cap Tr(P_1) \cap Tr(P_2) \]
Noting from this construction that \( Tr(P') \subseteq Tr(P_1), Tr(P_2) \) and from lemma 2 that \( \forall \sigma \in Tr(P') \), \( \text{Ref}(P, \sigma) = \text{Ref}(P_1, \sigma) = \text{Ref}(P_2, \sigma) \). In order to show that \( P' \) is the required common reduction of \( P_1 \) and \( P_2 \) we need to show that \( \forall \sigma \in Tr(P'), \text{Ref}(P', \sigma) \subseteq \text{Ref}(P_1, \sigma), \text{Ref}(P_2, \sigma) \).
This is enough because any \( \sigma' \in Tr(P_1) \cup Tr(P_2) \cdot \sigma' \notin Tr(P') \) will give \( \text{Ref}(P', \sigma') = \emptyset \), which trivially gives us the required refusal relationship.

We argue by contradiction that \( \forall \sigma \in Tr(P'), \text{Ref}(P', \sigma) \subseteq \text{Ref}(P_1, \sigma), \text{Ref}(P_2, \sigma) \). So, assume \( \exists \sigma \in Tr(P') \cdot \text{Ref}(P', \sigma) \supset (\text{Ref}(P_1, \sigma) = \text{Ref}(P_2, \sigma) = \text{Ref}(P, \sigma)) \). Thus, \( \exists \{ \sigma \} \notin (\text{Ref}(P_1, \sigma) = \text{Ref}(P_2, \sigma) = \text{Ref}(P, \sigma)) \) such that \( \{ \sigma \} \in \text{Ref}(P', \sigma) \). From here we can use lemma 3 to get \( a \notin \text{out}(P', \sigma) \), but it must also be the case that \( a \in \text{out}(P_1, \sigma), \text{out}(P_2, \sigma), \text{out}(P, \sigma) \) and thus we have a contradiction as the trace \( \sigma, a \) is in \( Tr(P) \cap Tr(P_1) \cap Tr(P_2) \). So, it must be the case that \( \text{Ref}(P', \sigma) \subseteq \text{Ref}(P_1, \sigma), \text{Ref}(P_2, \sigma) \) and \( P' \text{ red} P_1 \) and \( P' \text{ red} P_2 \) as required.

\textbf{Corollary 8}

\( C_{cs} \subseteq C_{cs} \).

\textbf{Proof}

From propositions 30 and 31.

Thus, \( C_{cs} \) is strictly stronger than \( C_{cs} \).

When the above results are combined with proposition 21 and corollary 6 we obtain a precise classification of \( C_{cs} \), as follows:

\textbf{Corollary 9}

\( C_{cs} \subseteq C_{cs} \subseteq C_{cs} \).

\textbf{Proof}

Immediate from proposition 21, corollary 6 and corollary 8.

In addition, we can put an upper bound on the strength of these relationships using the following proposition:

\textbf{Proposition 32}

\( C_{cs} \subseteq C_{cs} \).

\textbf{Proof}

This follows immediately from proposition 19, which relates \( C_{cs} \) to \( C_{te} \), corollary 5 and corollary 6, which shows that \( C_{cs} \leq C_{cs} \).

The relationship between the different interpretations of consistency are shown in figure 4.3. In addition, consistency based on behavioural compatibility can be incorporated into this categorisation through the following properties:

\[ C_{cs} \subseteq C_{cs}; C_{te} \subseteq C_{te}; C_{\sim} \subseteq C_{\sim}; \text{ and } C_{\sim} \subseteq C_{\sim}. \]

These instantiations present us with a number of possible interpretations of consistency in LOTOS. This situation reflects our view that consistency checking must be performed selectively, this issue was discussed in some depth in [9]. In particular, it is inappropriate to view consistency checking as a single mechanism which can be applied to any pair of specifications. For example, it would be inappropriate to check two specifications which express exactly corresponding functionality with \( C_{ext} \). An implication of this is that, in order to apply suitable consistency checks the relationship of the specifications being checked must be made available by the specifier(s).
4.4 Consistency Checking and Unification Techniques

The previous sections have classified the different LOTOS instantiations of consistency. In this section we will provide some results on how actually to check for consistency. Thus, we will provide mechanisms by which certain classes of consistency can be checked. Specifically, we provide LOTOS unification strategies for balanced binary consistency.

Although consistency implies that the set of unifications is not empty and we can determine the denotational semantics of the least unification in some cases, it would be useful for system development purposes to have a method to construct a unification within the specification language. This unification can then be used for further refinement or as the implementation specification.

The purpose of this section is to find operators for LOTOS that can be used to unify specifications. In case the original specifications are consistent (with respect to some notion of development), the unification should obviously be a common development (with respect to that notion of development). Such operational definitions of unification have several advantages over the unification algorithms that are applied to a denotational model of the specifications. The operational semantics of the unification operators give unification a constructive character which is useful for simulation and implementation purposes.

We will consider in order trace preorder, reduction, extension and then testing equivalence.

4.4.1 Trace preorder preserving unification

Parallel composition of specifications preserves the safety properties.

**Proposition 33 (Unification w.r.t. \( \leq_{tr} \))**

If \( P_1 \) and \( P_2 \) are two arbitrary LOTOS process specifications, then the process \( S := P_1 \parallel P_2 \) is a unification of \( P_1 \) and \( P_2 \) with respect to \( \leq_{tr} \).

**Proof**

We only prove \( P \parallel Q \leq_{tr} P \) as the other case is symmetric.
We derive that $\forall \alpha \in \mathcal{L} \cdot (P \parallel Q \xrightarrow{\alpha} P \xrightarrow{\alpha})$, and therefore $Tr(P \parallel Q) \subseteq Tr(P)$:
\[
P \parallel Q \xrightarrow{\alpha} \iff (\text{from definition of } \xrightarrow{\alpha})
\]
$\exists P', Q' \cdot P \parallel Q \xrightarrow{\alpha} P' \parallel Q' \land P' \parallel Q' \xrightarrow{\alpha} \iff (\text{from definition of } \parallel)$
$\exists P' \cdot P \parallel Q \xrightarrow{\alpha} P' \parallel P' \xrightarrow{\alpha} \iff (\text{by definition of } \xrightarrow{\alpha})$
$P \xrightarrow{\alpha} . \quad \square$

### 4.4.2 Reduction preserving unification

The following theorem gives a necessary and sufficient condition on two specifications for them to be consistent with respect to reduction. The condition requires that $P_1$ and $P_2$ can at least refuse all the actions they may not both do after a certain trace.

**Theorem 1 (consistency w.r.t. reduction)**

Let $P_1$, $P_2$ be two LOTOS specifications using the alphabet $\mathcal{L}$, then:

$P_1 \equivred P_2 \iff \forall \sigma \in Tr(P_1) \cap Tr(P_2) \cdot \mathcal{L} - \mathit{out}(P_1, \sigma) \cap \mathit{out}(P_2, \sigma) \in \mathit{Ref}(P_1, \sigma) \cap \mathit{Ref}(P_2, \sigma)$

**Proof**

"$\Leftarrow$" We need to prove that
\[
\forall \sigma \in Tr(P_1) \cap Tr(P_2) \cdot \mathcal{L} - \mathit{out}(P_1, \sigma) \cap \mathit{out}(P_2, \sigma) \in \mathit{Ref}(P_1, \sigma) \cap \mathit{Ref}(P_2, \sigma)
\]
implies $\exists P \cdot P \equivred P_1$ and $P \equivred P_2$.

Now, take $P$ to be the fully deterministic process that is completely determined (modulo strong bisimulation) by the intersection of the traces of $P_1$ and $P_2$, i.e. $Tr(P) = Tr(P_1) \cap Tr(P_2)$. For such a deterministic process, we have
\[
\forall \sigma \in Tr(P) \cdot \mathit{Ref}(P, \sigma) = \mathcal{P}(\mathcal{L} - \mathit{out}(P, \sigma)),
\]
where $\mathcal{P}(X)$ denotes the powerset of the set $X$.

Next, we prove that $P \equivred P_1$ and $P \equivred P_2$:

1. From $Tr(P) = Tr(P_1) \cap Tr(P_2)$, it follows that $Tr(P) \subseteq Tr(P_1)$ and $Tr(P) \subseteq Tr(P_2)$
2. From the definition of $P$, we derive that
\[
\forall \sigma \in Tr(P_1) \cap Tr(P_2), \forall X \in \mathit{Ref}(P, \sigma) \cdot X \subseteq (\mathcal{L} - \mathit{out}(P, \sigma)).
\]

As $Tr(P) = Tr(P_1) \cap Tr(P_2)$, it follows, by the definition of $\mathit{out}(P, \sigma)$ and the prefix-closedness of tracesets, that
\[
\forall \sigma \in Tr(P_1) \cap Tr(P_2) \cdot \mathit{out}(P, \sigma) = \mathit{out}(P_1, \sigma) \cap \mathit{out}(P_2, \sigma).
\]

From the condition of the proposition, we can now derive that
\[
\forall \sigma \in Tr(P_1) \cap Tr(P_2) \cdot X \in \mathit{Ref}(P, \sigma) \iff X \subseteq \mathit{Ref}(P_1, \sigma) \cap \mathit{Ref}(P_2, \sigma).
\]

Finally, we have $\forall \sigma \in Tr(P_1) - Tr(P) \cdot \mathit{Ref}(P, \sigma) = \emptyset \subseteq \mathit{Ref}(P_1, \sigma)$ and similarly for $P_2$, by $\mathit{Ref}(P, \sigma) = \emptyset \iff \sigma \notin Tr(P)$.

"$\Rightarrow$" Assume that there exists an $P$ such that $P \equivred P_1$ and $P \equivred P_2$. By contradiction:

Suppose $\mathcal{L} - \mathit{out}(P_1, \sigma) \cap \mathit{out}(P_2, \sigma) \notin \mathit{Ref}(P_1, \sigma) \cap \mathit{Ref}(P_2, \sigma)$, for a certain $\sigma \in Tr(P_1) \cap Tr(P_2)$.

Then there exists an $\alpha \in \mathcal{L} - \mathit{out}(P_1, \sigma) \cap \mathit{out}(P_2, \sigma)$ such that $\{\alpha\} \notin \mathit{Ref}(P_1, \sigma) \cap \mathit{Ref}(P_2, \sigma)$. It follows that either
\[
\alpha \in \mathit{out}(P_1, \sigma) \land \alpha \notin \mathit{out}(P_2, \sigma) \land \{\alpha\} \notin \mathit{Ref}(P_1, \sigma) \cap \mathit{Ref}(P_2, \sigma), \quad (4.1)
\]
\[
\alpha \notin \mathit{out}(P_1, \sigma) \land \alpha \in \mathit{out}(P_2, \sigma) \land \{\alpha\} \notin \mathit{Ref}(P_1, \sigma) \cap \mathit{Ref}(P_2, \sigma). \quad (4.2)
\]
As both cases are symmetric, we only prove case (1). Since \( a \notin \text{out}(P_2, \sigma) \) implies \( \{a\} \in \text{Ref}(P_2, \sigma) \) and \( \{a\} \notin \text{Ref}(P_1, \sigma) \cap \text{Ref}(P_2, \sigma) \) \( \land \{a\} \in \text{Ref}(P_1, \sigma) \) implies \( \{a\} \notin \text{Ref}(P_1, \sigma) \), we derive:

\[
\alpha \in \text{out}(P_1, \sigma) \land \alpha \notin \text{out}(P_2, \sigma) \land \{a\} \notin \text{Ref}(P_1, \sigma) \land \{a\} \in \text{Ref}(P_2, \sigma)
\]

For \( P \text{ red } P_1 \) and \( P \text{ red } P_2 \) to hold, it is necessary that \( (a) \sigma\alpha \notin \text{Tr}(P) \), because \( \sigma.a \notin \text{Tr}(P_2) \), and that \( (b) \{a\} \notin \text{Ref}(P, \sigma) \), because \( \{a\} \notin \text{Ref}(P_1, \sigma) \). Clearly, (a) and (b) can never be satisfied by any process. Hence, there does not exist a reduction of both \( P_1 \) and \( P_2 \). This contradicts with the assumption that such a common reduction exists. \( \square \)

Unfortunately, none of the existing LOTOS operators yields a common reduction when applied to two behaviour expressions. Therefore, we define a new operator. This binary operator, coined the conjunction operator, resolves all non-determinism present and removes all internal actions in either operand. The conjunction of two processes can only perform a trace if both processes are able to perform that trace.

**Definition 30 (conjunction operator)**

Let \( P \) and \( Q \) be LOTOS behaviour expressions. We define their conjunction, \( P \circ Q \), by the following inference rule (\( \Sigma(X) \) represents generalised choice):

\[
P \xrightarrow{\alpha} , Q \xrightarrow{\alpha} \implies P \circ Q \xrightarrow{\alpha} \Sigma(P \xrightarrow{\alpha} Q) \circ \Sigma(Q \xrightarrow{\alpha})
\]

**Proposition 34 (correctness of conjunction)**

If \( P_1, C \text{ red } P_2, \) then

\( P_1 \circ P_2 \text{ red } P_1 \) and \( P_1 \circ P_2 \text{ red } P_2 \)

**Proof**

We will prove only that \( P_1 \circ P_2 \text{ red } P_1 \). The other case is symmetric. The proof consists of two parts:

1. \( \text{Tr}(P_1 \circ P_2) \subseteq \text{Tr}(P_1) \):

   Inspecting the only inference rule, we see that \( P_1 \circ P_2 \xrightarrow{\alpha} \implies P_1 \xrightarrow{\alpha} \).

2. \( P_1 \circ P_2 \text{ conf } P_1 \):

   Suppose there exists a \( P \) such that \( P_1 \circ P_2 \xrightarrow{\alpha} P \xrightarrow{\alpha} \), where \( \alpha \in \text{Tr}(P_1), \alpha \in \mathcal{L} \). From the inference rule, it follows that \( P = P'_1 \circ P'_2 \), where \( P'_1 = \Sigma(P_1 \xrightarrow{-|\alpha}) \) and \( P'_2 = \Sigma(P_2 \xrightarrow{-|\alpha}) \). For \( P \) to refuse \( \alpha \) either \( P'_1 \xrightarrow{\alpha} \) or \( P'_2 \xrightarrow{\alpha} \). In the first case, we are ready, because \( P_1 \xrightarrow{\alpha} P'_1 \xrightarrow{\alpha} \).

In the case where \( P'_2 \xrightarrow{\alpha} \), we take the consistency condition \( \mathcal{L} \xrightarrow{-|\alpha} \text{out}(P_1, \sigma) \cap \text{out}(P_2, \sigma) \) into consideration. \( P'_2 \xrightarrow{\alpha} \) implies \( \{a\} \notin \text{out}(P_2, \sigma) \). By the condition, we derive that \( \{a\} \in \text{Ref}(P_1, \sigma) \cap \text{Ref}(P_2, \sigma) \), and therefore \( \{a\} \in \text{Ref}(P_1, \sigma) \). Hence there exists a \( P'' \in (P_1 \text{ after } \sigma) \) such that \( P_1 \xrightarrow{\alpha} P'' \xrightarrow{\alpha} \).

Although the conjunction operator yields a common reduction, it is usually not the least reduction of its two operands. This is due to the fact that all non-determinism is resolved. When both original specifications contain non-determinism it may not be necessary to resolve all non-determinism, as is shown in the following example.

**Example 3** Take \( P_1 := i; a; \text{stop}[i]; b; \text{stop}[i]; c; \text{stop} \) and \( P_2 := i; a; \text{stop}[i]; b; \text{stop} \). Then \( P_1 \circ P_2 = a; \text{stop}[b]; \text{stop} \), which is a common reduction, but not the least common reduction, which is \( i; a; \text{stop}[i]; b; \text{stop} \).
4.4.3 Extension preserving unification

From proposition 28 we know that any two specifications are consistent with respect to extension. However, none of the existing LOTOS operators will always yield a common extension. Therefore, we define a new operator. This binary operator, coined the join operator, merges as it were those patterns of behaviour that the two operand specifications have in common, and then provides a choice between the two behaviours when they start to differ. Note that the composition will be completely deterministic until a choice for either behaviour has been made.

**Definition 31 (join operator)**

Let $P$ and $Q$ be LOTOS behaviour expressions. We define their join, $P \uplus Q$, by the following inference rules:

$$
\begin{align*}
(1) & \quad P \overset{\alpha}{\rightarrow}, Q \overset{\alpha}{\rightarrow} \quad \Rightarrow \quad P \uplus Q \overset{\alpha}{\rightarrow} \Sigma(P \rightarrow_{\alpha}) \uplus \Sigma(Q \rightarrow_{\alpha}) \\
(2) & \quad P \overset{\alpha}{\rightarrow}, Q \overset{\beta}{\rightarrow} \quad \Rightarrow \quad P \uplus Q \overset{\alpha}{\rightarrow} \Sigma(P \rightarrow_{\alpha}) \\
(3) & \quad P \overset{\alpha}{\rightarrow}, Q \overset{\alpha}{\rightarrow} \quad \Rightarrow \quad P \uplus Q \overset{\alpha}{\rightarrow} \Sigma(Q \rightarrow_{\alpha})
\end{align*}
$$

**Proposition 35 (correctness of join)**

If $P_1$ and $P_2$ are two arbitrary LOTOS process specifications, then the process $P := P_1 \uplus P_2$ is a unification of $P_1$ and $P_2$ with respect to $\text{ext}$.

**Proof**

We will prove only that $P_1 \uplus P_2 \equiv \text{ext} P_1$. The other case is symmetric. The proof consists of two parts:

1. $\text{Tr}(P_1 \uplus P_2) \supseteq \text{Tr}(P_1)$:
   - Inspecting the inference rules, we see that $P_1 \overset{\alpha}{\rightarrow}$ implies $P_1 \uplus P_2 \overset{\alpha}{\rightarrow}$.

2. $P_1 \uplus P_2 \equiv \text{conf} P_1$
   - Suppose there exists a $P$ such that $P_1 \uplus P_2 \overset{\alpha}{\rightarrow} P \not\overset{\alpha}{\rightarrow}$, where $\sigma \in \text{Tr}(P_1)$, $\alpha \in \mathcal{L}$. From the inference rules, we derive the following three possible cases:
     
     (a) $P = P'_1 \uplus P'_2$, where $P'_1 = \Sigma(P_1 \rightarrow_{\beta})$ and $P'_2 = \Sigma(P_2 \rightarrow_{\beta})$, and $P'_1 \not\overset{\alpha}{\rightarrow}$ and $P'_2 \not\overset{\alpha}{\rightarrow}$.
     - In this case, it follows directly that there exists a $P' \equiv P'_1$ such that $P_1 \overset{\sigma'}{\rightarrow} P' \not\overset{\alpha}{\rightarrow}$.

     (b) $P = P'_1$, where $P'_1 = \Sigma(P_1 \rightarrow_{\beta})$, and $P'_1 \not\overset{\alpha}{\rightarrow}$. Clearly, $P_1 \overset{\sigma}{\rightarrow} P'_1 \not\overset{\alpha}{\rightarrow}$.

     (c) $P = P'_2$, where $P'_2 = \Sigma(P_2 \rightarrow_{\beta})$, and $P'_2 \not\overset{\alpha}{\rightarrow}$. We argue that this case will not occur. For it to exist there should exist traces $\sigma'$ and $\sigma''$ such that $\sigma' \sigma'' = \sigma$, process $P'' \in \Sigma(P_1 \rightarrow_{\beta'})$ and a $Q'' \in \Sigma(P_2 \rightarrow_{\beta''})$ such that $P_1 \overset{\sigma'}{\rightarrow} P'' \mathbin{\uplus} Q' = P'' \mathbin{\uplus} Q'$ while $Q'' \overset{\sigma''}{\rightarrow}$. However, $\sigma' \sigma'' \in \text{Tr}(P_2)$ and because $P'' \equiv \Sigma(P_1 \rightarrow_{\beta'})$ we have $P'' \overset{\sigma''}{\rightarrow}$. \qed

In the definition of the operational semantics of the join operator, we make use of a so-called negative premise (see [21]). This is potentially dangerous, because the transition relation may not be uniquely defined by the collection of all the inference rules of the language. Indeed, using unguarded recursion, which is allowed in LOTOS, we can show that the use of the negative premise here is not safe. Consider, for example, the following recursive process definition: $P := \alpha; \text{stop} \mathbin{\uplus} i; P$. Since the $\uplus$-operator abstracts from internal actions, the associated LTS is non image-finite, which makes it impossible to determine the next state. Even if unguarded recursion would be forbidden, then it would still be possible to make guarded recursion unguarded by applying the hide operator. Despite this, the $\uplus$ operator is meant to compose process specifications, with the aim to yield a new composed specification. A system will not be composed of itself and some other process, i.e. specifications of the shape $S_1 := S_1 \mathbin{\uplus} S_2$ do not make sense. Therefore, if the join operator is used for composition of specifications, its usage is safe.

Another drawback of the join operator is the fact that it does not yield the least common extension as is shown in the example below.
Example 4 Consider the following two gambling machine specifications: $P_1 = \text{coin; lose; stop} \parallel \text{coin; win; stop}$, $P_2 = \text{coin; lose; stop} \parallel \text{coin; win; (coin; jackpot; stop} \parallel \text{coin; lose; stop})$. Although the unification $P_1 \Join P_2 = \text{coin; (lose; stop} \parallel \text{win; coin; (jackpot; stop} \parallel \text{lose; stop})$ is an extension of both specifications, the $\Join$-operator makes it completely deterministic. This is clearly not the desired effect here. The least possible unification is given by $P_2$.

4.4.4 Testing equivalence preserving unification

With respect to testing equivalence, either of the two original specifications will do as the composition, because both specifications are testing equivalent to each other. However, this may result in the loss of the intention of the other specification.

4.5 Example of Unification in LOTOS

In order to demonstrate consistency checking and unification of partial specifications of a distributed information system, a simple Shared Memory system is introduced. We give two partial specifications of the system. One specification focuses on the computational aspects of the system, i.e. it describes the functional components of the system and the possible communication patterns between them. The second specification focuses on the information flow within the whole of the system without identifying the components it is composed of: it describes an invariant that should be satisfied by the data contained in the system. We thus obtain a nice separation of concerns.

Note, that the given specifications are not intended to be ODP compliant. They merely serve as a means to demonstrate the techniques for consistency checking and unification developed above.

The Shared Memory system is depicted in figure 4.4. It consists of a memory and two users. The users can access the memory through the read and write interfaces, but cannot communicate directly to one another. For simplicity we assume that the memory only contains one data value at a time.

4.5.1 Computational specification

In the computational specification the Shared Memory system is viewed as a collection of communicating processes. Two types of components are identified: a Memory component and a collection of User components. For this example, there are only two instantiations of the User process, but this could easily be extended to an arbitrary number. The User components do not communicate directly with each other. Therefore the two instantiations of User in process Users are placed in parallel ($\parallel$). There is communication between the Memory and the Users, which is represented by the synchronisation operator ($\|$) between the processes Memory and Users.

The computational specification is not concerned with the specific data values that are exchanged between communicating components. We see this reflected in the definition of the behaviour of a User. It specifies that a user can, at any time, either perform a read or a write operation, but this viewpoint does not care what the read or written data value is, represented by a non-deterministic choice.
process ComputationalSpec[read, write] : noexit :=
Memory [read, write] || Users [read, write]

where

process Users[read, write] : noexit :=
User [read, write] (0) ||| User [read, write] (1)

where

process User[read, write](uid : UserId) : noexit :=
choice x : Data []
i;
  (read!uid !x ; User [read, write] (uid)
   []
   write!uid ?x : Data ; User [read, write] (uid) )
endproc (* User *)
endproc (* Users *)

process Memory[read, write] : noexit :=
ConcurrentReads [read]
>
write?uid : UserId ?input : Data ; Memory [read, write]

where

process ConcurrentReads[read] : noexit :=
SequentialReads [read] ||| SequentialReads [read]

where

process SequentialReads[read] : noexit :=
choice uid : UserId, output : Data []
i; read!uid !output ; SequentialReads [read]
endproc (* SequentialReads *)
endproc (* ConcurrentReads *)
endproc (* Memory *)
endproc (* ComputationalSpec *)

The specification of the Memory component expresses that read operations can take place concurrently (process ConcurrentReads), but that these can at any time be disabled (>) by a write operation. This disabling ensures that write operations take place atomically in order to avoid data inconsistencies. Taking a closer look at process ConcurrentReads, we see that it actually only allows two concurrent read operations at one time. Obviously, this can easily be extended to a higher degree of concurrency. The individual threads of ConcurrentReads, process SequentialReads, allow for read operations by arbitrary users to happen sequentially. Again the data value is non-deterministically chosen.

4.5.2 Information specification

In the information specification, only the information flow within the Shared Memory system is considered. It specifies that the Memory is initialised with an arbitrarily chosen data value first (InitMem). From then on the following invariant must hold: the data value associated with each consecutive read operation is equal to the last written value. This is ensured by the process MemoryInvariant, which has one parameter to pass the last written value to it. The process MemoryInvariant is defined using a process ReadIn, which allows arbitrary users to do read operations while ensuring that mem is the data value read. Process ReadIn can at any time be disabled by a write operation taking place. The UserId and the data value associated with the write operation can be randomly chosen, but the written data value will consequently be passed to a new invocation of MemoryInvariant.

process InformationSpec[read, write] : noexit :=
InitMem >> accept mem : Data in MemInvariant [read, write] (mem)
where

\begin{verbatim}
process InitMem : exit(Data) :=
  choice mem : Data [] ; exit(mem)
endproc (* InitMem *)

process MemInvariant[read, write](mem : Data) : noexit :=
  Read_n [read] (mem)
|>
  (choice uid : UserId, x : Data []
    write!uid !x ; MemInvariant [read, write] (x))

endproc (* MemInvariant *)
endproc (* InformationSpec *)
\end{verbatim}

Note that we have generally avoided variable declarations of the form \( \text{read } ?\text{uid}:\text{UserId} \ ?\text{x}:\text{Data} \) as a shorthand for the set of actions \( \{ \text{read}<\text{uid},\text{x}> \mid \text{uid}\in\text{UserId}, \text{x}\in\text{Data} \} \). Instead we have made the choice more abstract in terms of non-determinism by using the construct:

\begin{verbatim}
choice uid : UserId, x : Data [] ; read!uid !x.
\end{verbatim}

Although this 'style' of specification is not required in a constraint oriented style, it is necessary to make the specifications consistent by reduction.

**4.5.3 Consistency check and unification**

First, we need to identify which instantiation of consistency applies here. As both viewpoint specifications in this example use the same event structure and the intention is for them to work together and not extend each others functionality, \( C_{\text{red}} \) seems more applicable than \( C_{\text{ext}} \). As it is unlikely the two specifications describe exactly the same behaviour, \( C_{\text{te}} \) is not applicable here. As for \( C_{\text{str}} \), this form of consistency is also covered by \( C_{\text{red}} \).

In order to show that the two specifications are consistent by \( C_{\text{red}} \) we have to verify the consistency condition of Theorem 1. Unfortunately, the presence of data variables in the specifications leads to a state explosion, which makes it hard to verify this condition.

Fortunately, it is possible to assess the consistency of the two specifications in an alternative way. If the specifications are consistent, the conjunction operator will yield a common reduction of both specifications. Thus, we first apply the conjunction operator, and then verify whether the unification is a reduction of either specification. Part of the LTS representing the conjunction, ComputationalSpec \( \circ \) InformationSpec, is shown in figure 4.5. In order to represent infinitely branching transitions caused by variable declarations, we have represented such transitions symbolically. It can be verified that this composition is indeed a reduction of both original specifications.

**4.6 Summary and Discussion**

This chapter has investigated consistency in LOTOS. We explored instantiations of consistency with a number of the LOTOS development relations. We have given appropriate LOTOS instantiations of the RM-ODP definitions of consistency and related these to our definition. This work supports the view presented in chapter 3 that our definition is general and can embrace all the RM-ODP definitions.

We have also characterised the relative strengths of different LOTOS instantiations of \( C \). The results of this are summarised in figure 4.3. Further, necessary and sufficient conditions were investigated for specifications to be consistent with respect to particular notions of development.

Several forms of unification are already supported in standard LOTOS. In fact, it can be argued that all binary operators enable some form of unification. The parallel operator has proved
especially useful for the unification of constraints in the constraint-oriented specification style [52]. However, the $||$ operator only supports unification w.r.t. trace preorder ($\leq_{tr}$), which is but a weak notion of development.

We have proposed two new operators for LOTOS to support unification with respect to refinement by reduction and extension. The latter operator, $\mathcal{W}$, was inspired by the specification merge operator, $\oplus$, in a pioneering paper on incremental specification [25]. Our $\mathcal{W}$ operator is an improvement of the $\oplus$-operator, in the sense that it can deal with non-deterministic specifications.

The problem of composing specifications with respect to reduction and extension have been reported on before. For processes modelled by acceptance trees, an algorithm to unify such processes with respect to reduction is given in [31]. In order to apply this algorithm to LOTOS specifications a denotational semantics in terms of acceptance trees must be provided. In [32] and in [35] algorithms are given that can be used to compose two specifications, such that the composition is an extension of both. The first algorithm applies to LTSs, but uses acceptance trees as an intermediate model. The second algorithm applies to rooted failure trees (RFTs), which can provide a denotational semantics for a subset of LOTOS. In contrast to these composition methods, we have given an operational semantics for unification. This enables us to unify specifications on the specification language level, rather than on the semantic level. This is useful for simulation purposes.

The general framework for consistency and unification will allow us to investigate more instantiations of consistency (for example with implementation relations based on action refinement) in future research. Further, we intend to develop software tools for simulation and construction of unification, and for assessing the consistency of specifications.
Chapter 5

Consistency Checking
Mechanisms in Z

In this part we describe a general strategy for unifying two Z specifications. In order to increase its applicability, it is not specific to any particular ODP viewpoint, nor is it tied to any particular instantiation of the architectural semantics. We show how to use this unification to check the consistency of two viewpoints written in Z. This is illustrated with a number of examples, including an information viewpoint specification from OSI Management. The link into the general consistency checking framework is made by using the logical definition of consistency, which can be integrated into the framework as previously discussed.

5.1 Unifying Viewpoint Specifications in Z

In this section we describe a general strategy for unifying two Z specifications. As described above we would like unification of two specifications to yield the least common refinement of both viewpoints. Such least unification is what we will investigate in this chapter. Unification of Z specifications will therefore depend upon the Z refinement relation, which is given in terms of two separate components - data refinement and operation refinement, [41]. Two specifications will thus be consistent if their unification is internally valid, and for Z this holds when the unification is free from contradictions (assuming the specifications that were unified were both internally valid). Thus to check the consistency of two specifications, we check for contradictions within the unified specification.

Z is a state based FDT, and Z specifications consist of informal English text interspersed with formal mathematical text. The formal part describes the abstract state of the system (including a description of the initial state of the system), together with the collection of available operations, which manipulate the state. One Z specification refines another if the state schemas are data refinements and the operation schemas are operation refinements of the original specifications state and operation schemas. Details of the language and its refinement relation are contained in introductory texts, for example [41, 42, 50].

The unification algorithm we describe is divided into three stages: normalization, common refinement (which we usually term unification itself), and re-structuring. This algorithm can be shown to be the least refinement of both viewpoints, [8]. Related work on the combination of Z specifications includes [2, 1].

Normalization identifies commonality between two different viewpoint specifications, and re-writes each specifications into a normal form suitable for unification in the following manner. Clearly, the two specifications that are to be unified have to represent the world in the same way within them (e.g. if an operation is represented by a schema in one viewpoint, then the other viewpoint has to use the same name for its (possibly more complex) schema too), and that the correspondences between the specifications have to have been identified by the specifiers involved.
These will be given by mappings that describe the naming, and other, conventions in force. Once the commonality has been identified, normalization re-names the appropriate elements of the specifications. Normalization will also expand data-type and schema definitions into a normal form. Examples of normalization are given in [41, 42].

Unification itself takes two normal forms and produces the least refinement of both. Because normalization will hide some of the specification structure introduced via the schema calculus, it is necessary to perform some re-structuring after unification to re-introduce the structure chosen by the specifier. We do not discuss re-structuring here.

5.1.1 State Unification

The purpose of state unification is to find a common state to represent both viewpoints. The state of the unification must be the least data refinement of the states of both viewpoints, since viewpoints represent partial views of an overall system description.

The essence of all constructions will be as follows. We unify declarations rather than types, so non-identical types with name clashes are resolved by re-naming, then we unify declarations as follows. If an element $x$ is declared in both viewpoints as $x : S$ and $x : T$ respectively, then the unification will include a declaration $x : U$ where $U$ is the least refinement of $S$ and $T$. The type $U$ will be the smallest type which contains a copy of both $S$ and $T$. For example, if $S$ and $T$ can be embedded in some maximal type then $U$ is just the union $S \cup T$.

Given two viewpoint specifications both containing the following fragment of state description given by a schema $D$:

\[
\begin{align*}
D \\
x : S \\
\text{pred}_S \\
\end{align*}
\]

\[
\begin{align*}
D \\
x : T \\
\text{pred}_T \\
\end{align*}
\]

we unify as follows

\[
\begin{align*}
D \\
x : S \cup T \\
x \in S \implies \text{pred}_S \\
x \in T \implies \text{pred}_T \\
\end{align*}
\]

whenever $S \cup T$ is well founded. If $S$ and $T$ cannot be embedded in a single type then the unification will declare $x$ to be a member of the disjoint union of $S$ and $T$, and the mechanism to describe disjoint unions has to be included in the unification. In these circumstances we again achieve the least refinement of both viewpoints.

Axiomatic descriptions are unified in exactly the same manner.

5.1.2 Operation Unification

Once the data descriptions have been unified, the operations from each viewpoint need to be defined in the unified specification. We assume all renaming of names visible to the environment has taken place. Unification of schemas then depends upon whether there are duplicate definitions. If an operation is defined in just one viewpoint, then it is included in the unification (with appropriate adjustments to take account of the unified state).

For operations which are defined in both viewpoint specifications, the unified specification should contain an operation which is the least refinement of both, w.r.t. the unified representation of state. The unification algorithm first adjusts each operation to take account of the unified state in the obvious manner, then combines the two operations to produce an operation which is a refinement of both viewpoint operations.
5.1. UNIFYING VIEWPOINT SPECIFICATIONS IN Z

The unification of two operations is defined via their pre- and post-conditions. Given a schema it is always possible to derive its pre- and post-conditions, [33]. Given two schemas \( A \) and \( B \) representing operations, both applicable on some unified state, then the unification of \( A \) and \( B \) is:

\[
U(A, B) = \begin{align*}
\text{pre } A & \lor \text{pre } B \\
\text{pre } A & \Rightarrow \text{post } A \\
\text{pre } B & \Rightarrow \text{post } B
\end{align*}
\]

where the declarations are unified in the manner of the preceding subsection. This definition ensures that if both pre-conditions are true, then the unification will satisfy both post-conditions. Whereas if just one pre-condition is true, only the relevant post-condition has to be satisfied. This provides the basis of the consistency checking method for object behaviour which we discuss below.

We show, in [8], that this construction is the least refinement of the two viewpoint specifications. It is also associative, allowing the natural extension of unification to an arbitrary finite number of viewpoints.

5.1.3 Example 1 - A classroom

As an illustrative example we perform state and operation unification on a simple specification of a classroom. The example consists of the state represented by the schema \( \text{Class} \), and operation \( \text{Leave} \). The two viewpoint specifications to be unified are:

\[
\begin{align*}
\text{Max} & : \mathbb{N} \\
\text{Class} & \\
\quad d & : \mathbb{P}[1, 2] \\
\quad \#d & \leq \text{Max} \\
\text{Leave} & \\
\quad \Delta\text{Class} & \\
\quad p? & : \{1, 2\} \\
\quad p? & \in d \\
\quad d' & = d \setminus \{p?\}
\end{align*}
\]

\[
\begin{align*}
\text{Min} & : \mathbb{N} \\
\text{Class} & \\
\quad d & : \mathbb{P}[2, 3, 4] \\
\quad \#d & \geq \text{Min} \\
\text{Leave} & \\
\quad \Delta\text{Class} & \\
\quad p? & : \{2, 3, 4\} \\
\quad \#d & > \text{Min} + 1 \\
\quad p? & \in d \\
\quad d' & = d \setminus \{p?, 2\}
\end{align*}
\]

As described above, we first unify the state model, i.e. the schema \( \text{Class} \) in this example, which becomes:

\[
\begin{align*}
\text{Class} & \\
\quad d & : \mathbb{P}[1, 2] \cup \mathbb{P}[2, 3, 4] \\
\quad d & \in \mathbb{P}[1, 2] \Rightarrow \#d \leq \text{Max} \\
\quad d & \in \mathbb{P}[2, 3, 4] \Rightarrow \#d \geq \text{Min}
\end{align*}
\]

With this unified state model we can unify the operation \( \text{Leave} \) on this state. To do so we calculate the pre and post-conditions in the usual manner, and for this we need to expand the schema \( \text{Leave} \) into normal form in each viewpoint. This will involve, for example, declaring \( p? : \mathbb{Z} \) and containing \( p? \in \{1, 2\} \) as part of the predicate for the description of \( \text{Leave} \) in the first viewpoint. The pre-condition of \( \text{Leave} \) in the first viewpoint is then \( p? \in d \cap \{1, 2\} \) (in fact this is the part of the pre-condition which is distinct from the pre-condition in the second viewpoint, the rest acting as a state invariant). Hence, the unified \( \text{Leave} \) becomes:
To show that the unified Leave is indeed a refinement of Leave in viewpoint one we will decorate elements in viewpoint one with a subscript one. We can then use the retrieve relation

$$R_1$$

$$\begin{array}{c}
\text{Class} \\
\text{Class}_1 \\
\end{array}$$

$$d_1 \in \{d\} \cap P\{1,2\}$$

to describe the refinement between the unified state and the state in the first viewpoint. To demonstrate the refinement is correct, we make the following deductions. Suppose \(\text{pre} \ \text{Leave}_1 \land \Delta R_1 \land \text{Leave}\), we have to show the result of this schema is compatible with \(\text{post} \ \text{Leave}_1\). Now if \(\text{pre} \ \text{Leave}_1\), then \(p^? \in d_1 \in \{d\} \cap P\{1,2\}\), and hence \(d' = d \setminus \{p^?\}\). Then \(d'_1 \in \{d'\} \cap P\{1,2\} = \{d \setminus \{p^?\}\} \cap P\{1,2\}\). So \(d'_1 = d' \cap \{1,2\} = (d \setminus \{p^?\}) \cap \{1,2\} = d_1 \setminus \{p^?\}\), since by \(\text{pre} \ \text{Leave}_1\), \(p^? \in \{1,2\}\). The deduction that \(\text{pre} \ \text{Leave}_1 \land R_1 \Rightarrow \text{pre} \ \text{Leave}\) is similar. These two deductions complete the proof that the unification is a refinement of viewpoint one. The case for viewpoint two is symmetrical.

5.1.4 Example 2 - Dining Philosophers

As a second illustration of unification in Z, we shall consider the following viewpoint specifications of the dining philosophers problem. The dining philosophers problem, [18], is a classic problem in synchronisation. A group of \(N\) philosophers sit round a table, laid with \(N\) forks. There is one fork between each adjacent pair of philosophers. Each philosopher alternates between thinking and eating. To eat, a philosopher must pick up its right-hand fork and then the left-hand fork. A philosopher cannot pick up a fork if its neighbour already holds it. To resume thinking, the philosopher returns both forks to the table.

The synchronisation needed is that no philosopher thinks or eats forever, and that the philosophers must not starve through deadlock. The three viewpoint specifications defined are the philosophers, forks and tables viewpoints. The philosophers and forks describe individual philosopher and fork objects and the operations available on those objects. The table viewpoint describes a system constructed from those objects and the synchronisation mechanism between operations upon them. We shall then describe the unification of the three viewpoints.

Although this example is not one of an ODP system, it provides a suitable illustration of the issues involved in viewpoint specification and consistency checking.

The Philosophers Viewpoint

This viewpoint considers the specification from the point of view of a philosopher. A philosopher either thinks, eats or holds her right fork. Note that since the latter is just a state of mind (for a philosopher!) there is no need to describe the operations from a forks point of view at all in this viewpoint.

\[
\text{PhilStatus} ::= \text{Thinking} \mid \text{HasRightFork} \mid \text{Eating}
\]

A philosopher object is just defined by the state of the philosopher.
Initially a philosopher is thinking.

\[ \text{InitPHIL} \]
\[ \text{PHIL}' \]
\[ \text{status}' = \text{Thinking} \]

We can now describe the operations available. A thinking philosopher can pick up its right-hand fork.

\[ \text{GetRightFork} \]
\[ \Delta \text{PHIL} \]
\[ \text{status} = \text{Thinking} \]
\[ \text{status}' = \text{HasRightFork} \]

Philosophers who hold their right fork can begin eating upon picking up their left-hand fork. Finally to resume thinking, a philosopher releases both forks.

\[ \text{GetLeftFork} \]
\[ \Delta \text{PHIL} \]
\[ \text{status} = \text{HasRightFork} \]
\[ \text{status}' = \text{Eating} \]

\[ \text{DropForks} \]
\[ \Delta \text{PHIL} \]
\[ \text{status} = \text{Eating} \]
\[ \text{status}' = \text{Thinking} \]

The Forks Viewpoint

This viewpoint specifies a fork object. Each fork is either free or busy. The fact that the philosopher might change state when a fork is picked up or dropped does not concern forks.

\[ \text{ForkStatus} ::= \text{Free} \mid \text{Busy} \]

The state of the fork is given by:

\[ \text{FORK} \]
\[ \text{fstatus} : \text{ForkStatus} \]

Initially a fork is free.

\[ \text{InitFORK} \]
\[ \text{FORK}' \]
\[ \text{fstatus}' = \text{Free} \]

We can now describe the operations available. A free fork can be picked up, and both forks can be released.

\[ \text{Acquire} \]
\[ \Delta \text{FORK} \]
\[ \text{fstatus} = \text{Free} \]
\[ \text{fstatus}' = \text{Busy} \]

\[ \text{Release} \]
\[ \Delta \text{FORK} \]
\[ \text{fstatus} = \text{Busy} \]
\[ \text{fstatus}' = \text{Free} \]
The Tables Viewpoint

This viewpoint has a number of schemas from the other viewpoints as parameters, these are given as empty schema definitions. Upon unification the non-determinism in this viewpoint will be resolved by the other viewpoint specifications, and thus unification will allow functionality extension of these parameters.

The parameters we require are:

<table>
<thead>
<tr>
<th>PHIL</th>
<th>InitPHIL</th>
<th>GetRightFork</th>
<th>ΔPHIL</th>
</tr>
</thead>
<tbody>
<tr>
<td>GetLeftFork</td>
<td>DropForks</td>
<td>FORK</td>
<td>ΔPHIL</td>
</tr>
<tr>
<td>InitFORK</td>
<td>Acquire</td>
<td>Release</td>
<td>ΔFORK</td>
</tr>
</tbody>
</table>

The system from the table viewpoint is defined by a collection of fork and philosopher objects:

\[
\begin{align*}
&\text{Table} \\
&\text{forks} : 1..N \rightarrow \text{FORK} \\
&\text{phils} : 1..N \rightarrow \text{PHIL}
\end{align*}
\]

Initially the table consists of forks and philosophers all in their respective initial states.

\[
\begin{align*}
&\text{InitTable} \\
&\exists \text{InitFORK, InitPHIL} \bullet \text{ran forks} = \{0\text{InitFORk}\} \land \text{ran phils} = \{0\text{InitPHIL}\}
\end{align*}
\]

Here we use promotion (ie the \(\theta\) operator) in the structuring of viewpoints, which allows an operation defined on an object in one viewpoint to be promoted up to an operation defined over that object in another viewpoint. As we can see, this can be used effectively to reference schemas in different viewpoints without their full definition.

In order to define operations on the table, we define a schema \(\Phi \text{Table}\) which will allow individual object operations to be defined in this viewpoint. See [42] for a discussion of the use of promotion.

\[
\begin{align*}
&\Phi \text{Table} \\
&\Delta \text{Table} \\
&\Delta \text{PHIL} \\
&\Delta \text{FORK} \\
&m? : 1..N \\
&n? : 1..N \\
&\text{phils}(n?) = \theta \text{PHIL} \\
&\text{phils}' = \text{phils} \uplus \{\text{phils}(n?) = \theta \text{PHIL}'\} \\
&\text{forks}(m?) = \theta \text{FORK} \\
&\text{forks}' = \text{forks} \uplus \{\text{forks}(m?) = \theta \text{FORK}'\}
\end{align*}
\]

Note that we use two inputs \(m?, n?\), because we want to control later the synchronisation between operations on forks and those on philosophers. System operations to get the left and right forks, and to drop both forks can now be defined.

\[
GLF \triangleq (\Phi \text{Table} \land \text{GetLeftFork} \land \text{Acquire} \land [n?, m? : 1..N \mid m? = n?]) \setminus (\Delta \text{FORK}, \Delta \text{PHIL})
\]
5.1. UNIFYING VIEWPOINT SPECIFICATIONS IN Z

\[
GRF \triangleq (\Phi_{\text{Table}} \land \text{GetRightFork} \land \text{Acquire} \land [n?, m? : 1..N \mid m? = (n? \mod N + 1)]) \setminus (\Delta\text{FORK}, \Delta\text{PHIL})
\]

\[
DF \triangleq (\Phi_{\text{Table}} \land \text{DropForks} \land \text{Release} \land [n?, m? : 1..N \mid m? = n?]) \setminus (\Delta\text{FORK}, \Delta\text{PHIL})
\]

The last schema in each conjunction performs the correct synchronisation between the individual object operations. For example, it forces operations \text{GetLeftFork} and \text{Acquire} to be performed on \text{phils}(n?) and \text{forks}(n?) in GLF; whilst in GRF, \text{GetRightFork} will be performed on \text{phils}(n?) and \text{Acquire} on \text{forks}(n? \mod N + 1).

**Unifying the Philosophers and Forks Viewpoints**

Since the fork and philosopher object descriptions are independent, i.e., there are no state or operation schemas in common, the unification of these two viewpoints is just the concatenation of the two specifications. We do not re-write that concatenation here.

**Unifying the Table, Philosophers and Forks Viewpoints**

The Table specification does have commonality with the other two viewpoints. For each state or operation schema defined in two viewpoints (i.e., the Table and one other), we build one schema in the unification. In fact, the separation and object-based nature (in a loose sense) of this example means that we will not make extensive use of unification by pre- and post-conditions. This is desirable, since it reduces the search for contradictions in the consistency checking phase. In fact, our experiences with viewpoint specifications confirms that such a viewpoint methodology is really only feasible if one adopts this object-based approach.

For example, the schema \text{FORK} defined in the Table viewpoint is just a parameter from the fork viewpoint, and consequently its unification will just be:

\[
\text{FORK}
\]

\[
fstatus : \text{ForkStatus}
\]

Similarly the unification of \text{GetLeftFork} from the Table and Philosophers viewpoint is

\[
\text{GetLeftFork}
\]

\[
\Delta\text{PHIL}
\]

\[
status = \text{HasRightFork}
\]

\[
status' = \text{Eating}
\]

since the pre-condition of \text{GetLeftFork} in Table is just false. Notice that this provides a mechanism in Z by which to achieve functionality extension across viewpoints in a manner previously not supported.

The remaining schemas can be unified in the obvious manner.

### 5.1.5 Example 3 - OSI Management

Our third example involves the application of Z in the ODP information viewpoint to the modelling of OSI Management, which has been investigated by a number of researchers [44, 53]. We show here how unification and consistency checking can be used with such modelling techniques by considering viewpoint specifications of sieve managed objects and their controlling CME agent.

To illustrate some of the techniques we consider two viewpoint specifications of an event reporting sieve object together with a third viewpoint which describes a CME agent and its manipulation of the sieve objects. In this simplified model we have not considered the relationships between managed objects, although a complete presentation would include their specification.

Within ODP, an information object template is modelled by a Z specification. An information object instance is then modelled as a Z specification instance (i.e., a specification complete with initialization of variables), and an ODP action is described by a Z operation.
The variable declarations in a state schema represent the attributes of a managed object. The state schema also describes the state invariant which constrains the values of the attributes. The initialization schema (e.g. \textit{InitSieve}) constrains the initial values of the state schema.

A managed object definition cannot include a \textit{Create} operation, since before it is created a managed object cannot perform any operation, including \textit{Create} itself. However, by including a \textit{Create} operation in the CME agent viewpoint as we do below, we can describe formally the interaction between \textit{Create} and the sieve managed object definition.

We have not considered any particular flavours, or design considerations, to differentiate between the first two viewpoints. Their purpose here is to represent to view of the system from similar standpoints.

\textbf{Viewpoint 1: Sieve object}

To describe the sieve object, we first declare the types. \textit{SieveConstruct} is used in the event reporting process, its internal structure is left unspecified at this stage, hence it is defined as a given set.

\begin{verbatim}
[SieveConstruct]
\end{verbatim}

The remaining types are declared as enumerated types.

\begin{verbatim}
Operational := disabled | active | enabled | busy
Admin := locked | unlocked | shuttingdown
Event := nothing | enrol | deenrol
Status := created | deleted
\end{verbatim}

\textit{Status} models the life-cycle of the sieve object, and is used as an internal mechanism to control which operations are applicable at a given point within an object’s existence. The state schema defines the attributes of the sieve object, here there are no constraints upon them; and the initialization describes their initial values.

\begin{verbatim}
Sieve
opstate : Operational
sico : SieveConstruct
adminstate : Admin
status : Status
\end{verbatim}

\begin{verbatim}
InitSieve
Sieve
opstate = active
adminstate = unlocked
\end{verbatim}

We describe two of the operations available within a sieve object (for a full description of operations see \cite{44}). The first is an operation to delete a sieve. Upon deletion a sieve sends a \textit{deenrol} notification to its environment, and moves into a state where no further operations can be applied.

\begin{verbatim}
Delete
\Delta Sieve
notification! : Event
status = created
notification! = deenrol
status' = deleted
\end{verbatim}

We define a relation \textit{filter} to represent criteria to decide which events to filter out and which to pass on

\begin{verbatim}
| filter : Event \rightarrow SieveConstruct
\end{verbatim}

and the \textit{Filter} schema represents the operation to perform the filtering.
5.1. UNIFYING VIEWPOINT SPECIFICATIONS IN Z

### Filter

```
Filter
| Sieve
| event? : Event
| notification! : Event
```

\[
\begin{align*}
\text{status} &\neq \text{deleted} \\
\text{opstate} = \text{active} &\land \text{adminstate} = \text{unlocked} \\
(\text{event}?, \text{sico}) &\in \text{filter} \Rightarrow \text{notification!} = \text{event}? \\
(\text{event}?, \text{sico}) &\notin \text{filter} \Rightarrow \text{notification!} = \text{nothing}
\end{align*}
\]

### Viewpoint 2: Sieve object

To illustrate some of the unification techniques we now describe a second view of the same sieve object. First of all we declare the types

```
[SieveConstruct]
```

```
\begin{align*}
\text{Operational} &::= \text{disabled} \mid \text{active} \mid \text{enabled} \mid \text{busy} \\
\text{Admin} &::= \text{locked} \mid \text{unlocked} \mid \text{shutting\textendash down} \\
\text{Event} &::= \text{nothing} \mid \text{enrol} \mid \text{deenrol} \\
\text{Status} &::= \text{being\_created} \mid \text{created} \mid \text{deleted}
\end{align*}
```

Notice that in this viewpoint `Status` includes an additional value, `being\_created`. The state schema and its initialization are then declared.

```
Sieve
| opstate : Operational \\
| sico : SieveConstruct \\
| adminstate : Admin \\
| status : Status
```

```
InitSieve
| Sieve
| status = being\_created \\
| opstate = active \\
| adminstate = unlocked
```

The change of state of a sieve object from `being\_created` to `created` is governed by an internal operation, which can occur spontaneously. This change in `status` allows other operations to be invoked subsequently apart from the `Enrol` operation itself.

```
Enrol
| ΔSieve
| notification! : Event
```

```
\begin{align*}
\text{status} & = \text{being\_created} \\
\text{status}' & = \text{created} \\
\text{notification!} & = \text{enrol}
\end{align*}
```

In this viewpoint a relation `newfilter` defines the filtering criteria, and the operation `Filter` performs the filtering.

```
| newfilter : Event ← SieveConstruct
```
**Viewpoint 3 : CME agent**

The final viewpoint is a description of a controlling CME agent. For our purposes here we present a very simplified version of an agent which consists of a number of sieve managed objects. We then show how we can promote the Delete operation defined on individual sieve objects, and define a Create operation to instantiate sieve objects as required.

This viewpoint has a number of schemas from the other viewpoints as parameters, these are given as empty schema definitions. Upon unification the under-specification of these parameters in this viewpoint will be resolved by the other viewpoint specifications, and thus unification will allow functionality extension of these parameters. The parameters we require are:

- **Sieve**
- **InitSieve**
- **Delete**
- **ΔSieve**

We declare types to represent the set of object classes and set of object identifiers respectively. A **CMEagent** is then modelled as a collection of sieve objects, and initially no sieve objects have been created, so the range of **sieves** cannot include a state described by **Sieve**

\[
\begin{align*}
\text{CMEagent} & : (\text{Class} \times \text{Id}) \\
\text{sieves} & : \text{Class} \times \text{Id} \rightarrow \text{Sieve} \\
\text{InitCMEagent} & : \text{CMEagent} \\
\text{CMEagent} & , \text{Sieve} \\
\neg \text{Sieve} & \in \text{ran} \text{sieves}
\end{align*}
\]

In order to define CME agent operation, we define a schema \(\Phi CMEagent\) which will allow individual object operations to be defined in this viewpoint.

\[
\begin{align*}
\Phi CMEagent & \\
\Delta CMEagent & \\
\Delta Sieve & \\
objectclass? & : \text{Class} \\
sieveid? & : \text{Id} \\
sieves(\text{objectclass}?, \text{sieveid}?) & = \text{0Sieve} \\
sieves' & = \text{sieves} \uplus \{\text{sieves(\text{objectclass}?, \text{sieveid}?) = 0Sieve'}
\end{align*}
\]

An agent operation to delete a specific sieve object can now be defined by promotion of the **Delete** parameter specified in another viewpoint. The other managed object operations are promoted in a similar fashion.

\[
\text{DeleteSieve} \equiv (\Phi CMEagent \land \text{Delete}) \setminus (\Delta Sieve)
\]
Finally the \textit{Create} operation can be defined. Notice this is not part of the sieve specification, so we have preserved the concept that \textit{Create} must occur before any operation in the sieve specification can be applied.

\begin{verbatim}
Create
ΔCMEagent
ΔSieve, ΔInitSieve
objectclass? : Class
sieveid? : Id

sieves(objectclass?, sieveid?) ≠ ∅Sieve
sieves' = sieves ⊕ {sieves(objectclass?, sieveid?) = ∅InitSieve'}
\end{verbatim}

\subsection*{Unification of Viewpoints}

To describe the unification of viewpoints, we decorate with subscripts, so for example \textit{Filter}_1 is the schema \textit{Filter} from the first viewpoint. To unify viewpoints 1 and 2 we first unify the state. The only conflict in the declarations are due to differing types \textit{Status}_1 and \textit{Status}_2. To resolve this conflict, the type \textit{Status} in the unification is taken as the least refinement of \textit{Status}_1 and \textit{Status}_2 (i.e. \textit{Status}_1 \cup \textit{Status}_2), and state unification is applied to the schema \textit{Sieve}. Hence, in addition to the declarations which are not in conflict, the unification will contain the following:

\[\text{Status} ::= \text{being-created} \mid \text{created} \mid \text{deleted}\]

\begin{verbatim}
Sieve
opstate : Operational
sico : SieveConstruct
adminstate : Admin
status : Status

status ∈ \{created, deleted\} ⇒ true
status ∈ \{being_created, created, deleted\} ⇒
(opstate ∈ \{active, disabled\} ∧ adminstate ∈ \{locked, unlocked\})
\end{verbatim}

Upon simplification the schema \textit{Sieve} becomes

\begin{verbatim}
Sieve
opstate : Operational
sico : SieveConstruct
adminstate : Admin
status : Status

opstate ∈ \{active, disabled\}
adminstate ∈ \{locked, unlocked\}
\end{verbatim}

In a similar fashion we unify \textit{InitSieve}_1 and \textit{InitSieve}_2, which simplifies to

\begin{verbatim}
InitSieve
Sieve

status = being_created
opstate = active
adminstate = unlocked
\end{verbatim}
The \textit{Delete} and \textit{Enrol} schemas are defined in just one viewpoint. Hence, both these schemas are included in the unification (with adjustments due to the unified state schema \textit{Sieve}). Similarly the unification contains both relations \textit{filter} and \textit{newfilter}.

To unify \textit{Filter}_1 and \textit{Filter}_2 we first adjust \textit{Filter}_1 due to the unified state schema. The predicate part of \textit{Filter}_1 is then

\[ \text{status} \in \{ \text{created, deleted} \} \Rightarrow (\text{status} \neq \text{deleted} \land \text{opstate} = \text{active} \land \text{adminstate} = \text{unlocked}) \]

Calculation of the pre-conditions \textit{preFilter}_1 \lor \textit{preFilter}_2 then simplifies to

\[ (\text{status} = \text{created} \land \text{opstate} = \text{active} \land \text{adminstate} = \text{unlocked}) \]

Thus the unification of \textit{Filter}_1 and \textit{Filter}_2 is then given by:

\[
\begin{array}{l}
\text{Filter} \\
\text{\textit{Sieve}} \\
\text{\textit{Event}} \\
\text{\textit{Notification}} \\
\text{\textit{status} = \text{created} \land \text{opstate} = \text{active} \land \text{adminstate} = \text{unlocked}} \\
(\text{\textit{event}}?, \text{\textit{sico}}) \in \text{\textit{filter}} \Rightarrow \text{\textit{notification}}! = \text{\textit{event}}? \\
(\text{\textit{event}}?, \text{\textit{sico}}) \notin \text{\textit{filter}} \Rightarrow \text{\textit{notification}}! = \text{nothing} \\
(\text{\textit{event}}?, \text{\textit{sico}}) \notin \text{\textit{newfilter}} \Rightarrow \text{\textit{notification}}! = \text{nothing} \\
\end{array}
\]

To complete the unification we must unify this specification with the third viewpoint which specified the CME agent. The parameters in the third viewpoint have their functionality extended upon unification. For example, the schema \textit{InitSieve} defined in the third viewpoint is just a parameter from the other viewpoints, and consequently its unification will just be:

\[
\begin{array}{l}
\text{InitSieve} \\
\text{\textit{Sieve}} \\
\text{\textit{status} = \text{being}_\text{created}} \\
\text{\textit{opstate} = \text{active}} \\
\text{\textit{adminstate} = \text{unlocked}} \\
\end{array}
\]

The complete unification is then achieved in the obvious manner, by expanding \textit{Sieve}, \textit{InitSieve} and \textit{Delete} and including \textit{Enrol} and \textit{Filter} along with the CME agent operations.

## 5.2 Consistency Checking of Viewpoint Specifications in Z

The mechanism for unifying two Z specification yields a consistency checking process. A specification is consistent if it does not contain specifications of entities which cannot possibly exist. That is, given a proof system for Z, with a validity relation \( \vdash \), a specification is said to be consistent if it is not possible to prove \( S \vdash \text{false} \). For example, a \textit{Z} specification of a function will be inconsistent if the predicate part of its axiomatic definition contradicts the fact that it was declared as a function. Another way in which inconsistencies can arise in \textit{Z} specifications is in the definition of free types. Examples of how such inconsistencies can occur are given in \cite{4,48,45}. In general, it is undecidable whether or not a set of axioms given in a \textit{Z} specification is consistent. \cite{4} discusses sufficient conditions for the consistency of certain combinations of \textit{Z} paragraphs, in particular axiomatic definitions, given sets and free types.

In addition, consistency usually refers to consistency of the state model, i.e. for a given state there exists at least one possible set of bindings that satisfies the state invariant, \cite{41,42}. With
this consistency condition comes a requirement to prove the Initialisation theorem (see below), which asserts there exists a state that satisfies the initial conditions of the model. Due to an ODP requirement associated with multiple viewpoints, we also require operation consistency, because an ODP conformance statement in Z corresponds to an operation schema(s), [47]. A conformance statement is behaviour one requires at the location that conformance is tested. Thus a given behaviour (i.e. occurrence of an operation schema) conforms if the post-conditions and invariant predicates are satisfied in the associated Z schema. That is, operations defined in two viewpoints are consistent if whenever both operations are applicable, their post-conditions agree. Hence, operations in a unification will be implementable whenever each operation has consistent post-conditions on the conjunction of its pre-conditions.

Thus a viewpoint consistency check in Z involves checking the unified specification for contradictions, and has 5 components: axiom, axiomatic, state and operation consistency in addition to the Initialisation theorem. Assuming the individual viewpoints themselves are consistent, the components then take the following form.

**Axiom Consistency**: Axioms constrain existing global constants. Hence, to check for consistency of the two viewpoints, axioms from one viewpoint have to be checked against the second viewpoint w.r.t. any terms appearing in the axioms which are defined in the second viewpoint. If an axiom contains no terms appearing in other viewpoints, its consistency checking requirements are discharged.

**State Consistency**: Consider the general form of state unification given earlier:

\[
\begin{align*}
D & \quad x : S \cup T \\
& \quad x \in S \implies \text{pred}_S \\
& \quad x \in T \implies \text{pred}_T
\end{align*}
\]

This state model is consistent as long as both \(\text{pred}_S\) and \(\text{pred}_T\) can be satisfied for \(x \in S \cap T\).

**Axiomatic Consistency**: Similar to state consistency.

**Operation Consistency**: Consistency checking also needs to be carried out on each operation in the unified specification. The definition of operation unification means that we have to check for consistency when both pre-conditions apply. That is, if the unification of \(A\) and \(B\) is denoted \(U_A(B, B)\), we have:

\[
\text{pre } U(A, B) = \text{pre } A \lor \text{pre } B, \quad \text{post } U(A, B) = (\text{pre } A \implies \text{post } A) \land (\text{pre } B \implies \text{post } B)
\]

So the unification is consistent whenever the post-conditions agree on the conjunction of the pre-conditions, \((\text{pre } A \land \text{pre } B)\).

**Initialisation Theorem**: The Initialisation Theorem is a consistency requirement of all Z specifications. It asserts that there exists a state of the general model that satisfies the initial state description, formally it takes the form:

\[\vdash \exists \text{State } \bullet \text{InitState}\]

For the unification of two viewpoints to be consistent, clearly the Initialisation Theorem must also be established for the unification.

The following result can simplify this requirement: Let \(\text{State}\) be the unification of \(\text{State}_1\) and \(\text{State}_2\), and \(\text{InitState}\) be the unification of \(\text{InitState}_1\) and \(\text{InitState}_2\). If the Initialisation Theorem holds for \(\text{State}\) and \(\text{State}_2\), then state consistency of \(\text{InitState}\) implies the Initialisation Theorem for \(\text{State}\). In other words, it suffices to look at the standard state consistency of \(\text{InitState}\).

If, however, \(\text{InitState}\) is a more complex description of initiality (possibly still in terms of \(\text{InitState}_1\) and \(\text{InitState}_2\)), the Initialisation Theorem expresses more than state consistency of \(\text{Initstate}\), and hence will need validating from scratch.
5.2.1 Example 1 - The classroom

State Consistency: The unified state in this example was given by

<table>
<thead>
<tr>
<th>Class</th>
</tr>
</thead>
<tbody>
<tr>
<td>( d : \mathbb{P} {1, 2} \cup \mathbb{P} {2, 3, 4} )</td>
</tr>
<tr>
<td>( d \in \mathbb{P} {1, 2} \Rightarrow #d \leq \text{Max} )</td>
</tr>
<tr>
<td>( d \in \mathbb{P} {2, 3, 4} \Rightarrow #d \geq \text{Min} )</td>
</tr>
</tbody>
</table>

To show consistency, we need to show that if \( d \in \mathbb{P} \{1, 2\} \cap \mathbb{P} \{2, 3, 4\} \), then both \( \#d \leq \text{Max} \) and \( \#d \geq \text{Min} \) hold. Suppose the class consisted of just the element 2, i.e. \( d = \{2\} \). Both pre-conditions in the unified state, \( d \in \mathbb{P} \{1, 2\} \) and \( d \in \mathbb{P} \{2, 3, 4\} \), now hold giving the state invariant \( \text{Min} \leq \#d \leq \text{Max} \). Thus the consistency of the viewpoint specifications of the classroom requires that \( \text{Min} \leq 1 \leq \text{Max} \). This type of consistency condition should probably fall under the heading of a correspondence rule in ODP, [30], that is a condition which is necessary but not necessarily sufficient to guarantee consistency.

Operation Consistency: In this example, this amounts to checking the operation \( \text{Leave} \) when

\[
(p \in d \cap \{1, 2\}) \land (p \in d \cap \{2, 3, 4\} \land \#d > \text{Min} + 1)
\]

In these circumstances, the two post-conditions are \( d' = d \setminus \{p\} \) and \( d' = d \setminus \{p', 2\} \). These two pre-conditions apply when \( p = 2 \) and \( 2 \in d \). A consistency check has to be applied for all possible values of \( d \). For example, let \( d = \{1, 2\} \), then \( d' = d \setminus \{p\} \). If further \( \#d > \text{Min} + 1 \), then in addition we have \( d' = d \setminus \{p', 2\} \). These two conditions are consistent (since \( p = 2 \)) regardless of Max or Min.

Let \( d = \{2\} \), then both pre-conditions apply iff \( \text{Min} < 0 \), in which case the post-conditions are \( d' = d \setminus \{2\} \) and \( d' = d \setminus \{2\} \), and thus consistent.

Hence the two viewpoint specifications are consistent whenever the correspondence rule \( \text{Min} \leq 1 \leq \text{Max} \) holds.

5.2.2 Example 2 - Dining Philosophers

Inspection of the unification in the Dining Philosophers example shows that both state and operation consistency is straightforward (note, however, that with non-object based viewpoint descriptions of this example, consistency checking is a non-trivial task, this points the need for further work on specification styles to support consistency checks). Hence, consistency will follow once we establish the Initialization Theorem for the unification.

The Initialization Theorem for the unification is

\[ \vdash \exists \text{Table}' \cdot \text{InitTable} \]

which expands to

\[ \vdash \exists \text{forks}' : 1..N \rightarrow \text{FORK}, \text{phil}' : 1..N \rightarrow \text{PHIL} \cdot \exists \text{InitFORK}, \text{InitPHIL} \cdot \text{ran forks}' = \{\emptyset \text{InitFORK}\} \land \text{ran phil}' = \{\emptyset \text{InitPHIL}\} \]

Upon simplification this becomes

\[ \vdash \exists \text{forks}' : 1..N \rightarrow \text{FORK}, \text{phil}' : 1..N \rightarrow \text{PHIL} \cdot \text{ran forks}' = \{\text{Free}\} \land \text{ran phil}' = \{\text{Thinking}\} \]

which clearly can be satisfied. Hence the viewpoint descriptions given for the dining philosophers are indeed consistent.
5.2.3 Example 3 - OSI Management

State Consistency: Consider the state schema Sieve. The unified schema across all three viewpoints was given by:

\[
\begin{align*}
\text{Sieve} & \equiv \\
\text{opstate} : \text{Operational} & \\
\text{sico} : \text{SieveConstruct} & \\
\text{adminstate} : \text{Admin} & \\
\text{status} : \text{Status} & \\
\text{status} \in \{ \text{created, deleted} \} & \Rightarrow \text{true} \\
\text{status} \in \{ \text{being\_created, created, deleted} \} & \Rightarrow \\
(\text{opstate} \in \{ \text{active, disabled} \} \land \text{adminstate} \in \{ \text{locked, unlocked} \}) & \\
\end{align*}
\]

From this it can be seen that both predicates true and \((\text{opstate} \in \{ \text{active, disabled} \} \land \text{adminstate} \in \{ \text{locked, unlocked} \})\) can be satisfied for \text{status} \in \{ \text{created, deleted} \} \cap \{ \text{being\_created, created, deleted} \}, which is the requirement for consistency for this state schema.

Operation Consistency: Consider the unification of the Filter operation schema. From the unification we found that

\[
\text{prefFilter}_1 = \text{prefFilter}_2 = (\text{status} = \text{created} \land \text{opstate} = \text{active} \land \text{adminstate} = \text{unlocked})
\]

Thus to show operation consistency we have to show that under this pre-condition, we have

\[
((\text{event?}, \text{sico}) \in \text{filter} \Rightarrow \text{notification!} = \text{event?}) \land ((\text{event?}, \text{sico}) \notin \text{filter} \Rightarrow \text{notification!} = \text{nothing}) \\
= ((\text{event?}, \text{sico}) \notin \text{newfilter} \Rightarrow \text{notification!} = \text{nothing})
\]

It is easy to show that a necessary, but not sufficient, condition for the consistency of this operation is

\[
\forall (\text{event?}, \text{sico}) \in \text{Event} \times \text{Event} \bullet (\text{event?}, \text{sico}) \in \text{filter} \land (\text{event?}, \text{sico}) \notin \text{newfilter} \Rightarrow \text{event?} = \text{nothing}
\]

Thus the consistency of Filter requires this condition to be maintained. By giving specifiers explicit notification of which relationships between objects in the viewpoints need preserving, this constraint can then be used by the individual viewpoint specifiers to ensure that any further refinement of the viewpoints do not violate consistency.

Inspection of the unification of the CME agent with the sieve objects shows that both state and operation consistency carry over from the state and operation consistency of the unification of viewpoints 1 and 2. Hence, consistency will follow once we establish the Initialization Theorem for the unification of all three viewpoints.

The Initialization Theorem for the unification is

\[
\vdash \exists \text{CMEagent} \bullet \text{InitCMEagent}
\]

which expands to

\[
\vdash \exists \text{sieves} : \text{Class} \times \text{Id} \to \text{Sieve} \bullet \exists \text{Sieve} \bullet \theta \text{Sieve} \in \text{ran sieves}
\]

Upon simplification this becomes

\[
\vdash \exists \text{sieves} : \text{Class} \times \text{Id} \to \text{Sieve} \bullet (\text{status}, \text{opstate}, \text{adminstate}, \text{sico}) \notin \text{ran sieves}
\]

The final term describes a set of bindings, and it is clear that such a function \text{sieves} exists. Hence, the viewpoint descriptions given for the CME agent and sieve objects are indeed consistent.

The consistency checking mechanism works well for small to medium sized Z specifications. For larger specifications additional structure is needed in order that the consistency checking strategy can be scaled up. [17] shows how support for this can be provided by using object oriented variants of Z. These object based methodologies look likely to provide sufficient structure for the consistency checking to remain feasible, this is an area of ongoing research and is discussed briefly below.
5.3 Software Engineering Issues

The previous section elucidated the consistency checking requirements for unified viewpoint specifications. Given these requirements, it is necessary to seek software engineering strategies that make viewpoint decomposition feasible. By feasible we mean it is possible to describe viewpoints which are consistent and that the effort involved in consistency checking is minimised. In this section we explore some of the software engineering consequences of the consistency checking requirements. There are two initial possibilities for how to write viewpoint specifications:

(a) viewpoint description and analysis will work with arbitrary specifications and specification styles;
(b) some style guidelines or further methodology is needed for the process to become feasible.

If one adopts position (a), then there are some issues which need to be addressed:

No encapsulation of state and operations When the state is unified, all operations acting on that state are adjusted to take account of the unified state. Hence, during unification of an operation, two adjustments are made: the first due to the unified state (declarations are updated to take account of the unified state) and the second due to the change in pre- and post-conditions. Therefore, to keep track of the consistency checking requirements the operations need to be encapsulated with the state they affect. Without this consistency checking is possible, but unrealistic for larger examples.

No operation set representation As Zave noted in [56], Z provides no means of representing the operation set of a specification (i.e. the set of operations visible by the environment). The consequences of this for unification is that if an operation schema appears in both viewpoints, then it has to be unified, since there is no means to tell whether it was defined (in either viewpoint) for internal structuring purposes only. If there was such information available, then internal structuring schemas could be re-named and just operations in the operational set unified.

Correspondence rules In order to describe the relation between viewpoints, the RM-ODP includes a notion of correspondence rule. Part of their purpose is to identify the commonality between the specifications, and describe any possible renamings between them. Any viewpoint methodology will need to include mappings such as these. The limited structure in an ordinary Z specification makes a succinct naming impossible for correspondences, since, for any non-trivial systems, it is likely that a correspondence will wish to name more than one state/operation.

Viewpoint encapsulation The work of [2] indicates that in a non-object approach a large number of re-namings and re-workings of the viewpoints have to be undertaken during the unification process. This appears to be because the boundaries of the problem are not well defined, leading to viewpoint specifiers referencing and defining similar aspects of the same entity. Again this is a manifestation of the lack of encapsulation when defining the area of concern for each viewpoint specifier.

From case studies undertaken and consideration of these issues it seems that viewpoint description without any style guidelines is unlikely to be practical for anything other than small examples. Encapsulation and identity are central to the practical realization of the viewpoints model. Both of these facets could be provided by a number of software engineering methodologies, however, object orientation is an obvious choice. Many of the problems identified above can be resolved if one adopts an object oriented approach.

Encapsulation of state and operation The over-riding advantage of object oriented methods is their encapsulation of state and operation. This will clearly delimit the consistency checking requirements within a unification, with each unified object generating local consistency checking requirements which do not escalate to global consistency checking problems.
**Operation set representation** Some, although not all, object oriented methodologies in Z provide the ability to specify an operation set, or visibility list, \([39, 19]\), which partitions all the defined operations into disjoint sets of visible and internal operations. If this is provided, then the issue of operation set representation is completely resolved. Even if such a visibility list is not provided by the language used, the encapsulation that comes with object orientation still provides the opportunity for partial resolution of the problem. In an object-based world it is likely that a viewpoint partitioning will include the internal specification of the behaviour of an object in only one of the viewpoints. The other viewpoints will then (possibly) reference objects from viewpoints as parameters, or place constraints on the use of those objects. Hence, in these circumstances the unification of two internal representations is unlikely to occur, and so the issue of operation set representation would not occur. Of course, if the internal specification of an object’s behaviour did occur in more than one viewpoint, the need for a visibility list would then arise again.

**Correspondence rules** Identity is a key property of an object, and will allow correspondence rules to relate suitably complex parts and combination of parts of the viewpoint descriptions in a manner which is not currently supported in Z.

These considerations naturally lead to a choice of an object-based or object oriented language for viewpoint decomposition, where each viewpoint specifies a number of interacting objects. Full OO is not necessarily needed, however, if it is available then OO facilities such as inheritance can be exploited. It is preferable that only one viewpoint specifies the internal representation of a given object, and references to objects from one viewpoint will appear as parameters, either as inheritance within another object, or as an abstraction purely in terms of object or method names. The next section investigates the support available in Z for this approach.

### 5.4 Using Object Oriented Techniques

The previous section indicated that an object oriented style of specification is particularly suitable for viewpoint descriptions, and indeed the RM-ODP has adopted such an approach. There have been a number of different approaches proposed for providing Z with object oriented facilities. These include the provision of object oriented style guidelines, and extensions to Z to allow fully object oriented specifications. Examples of using Z in an object oriented style include: Hall’s style \([22, 23]\); ZERO \([54]\); and the ODP architectural semantics \([30]\). Examples of object oriented extensions to Z include: Object-Z \([19]\); ZEST \([15]\); MooZ \([39]\); OOZE \([3]\); Schuman & Pitt \([46]\); and Smith \([49]\). See \([51]\) for a summary and comparison of several approaches.

Z itself is not object oriented because it does not provide sufficient support for either encapsulation or inheritance. However, it is also possible for Z to be used in an object based fashion, see \([42]\) for a discussion, although there is nothing to keep the specifier within an object based style in contrast to the style guidelines above.

The ODP standardisation initiative requires the use of (near) standardised formal methods, hence the architectural semantics uses Z as opposed to any object oriented variant of that language. However, given that ODP has adopted the object oriented paradigm, there is obvious interest in object oriented variants that can support the required ODP modelling concepts. In particular, Object-Z and ZEST are receiving attention within the ODP community as a specification medium. However, all the object oriented extensions to Z have an unstable definition, or lack a full semantics, or both. Therefore, techniques with a flattening (or approximate flattening) into Z are of considerable interest to our work. By using such a technique we can define unification and consistency checking of viewpoints without compromising the necessity of a standardised formal description technique. Object-Z and ZEST are suitable from this perspective.

Of the Z guidelines the work of Hall is the most general. The style adds no new features to Z, however, there are conventions for writing an object oriented specification. He also provides conventions for modelling classes and their relationships, and, in addition, there is formal support for inheritance through subtyping, \([23]\). In order to support encapsulation, the RM-ODP has
adopted conventions for the use of Z within ODP. Here, encapsulation is achieved by letting each
Z specification denote just one object. This achieves the required encapsulation, but clearly any
specification of an aggregate of objects or interaction between objects cannot then be modelled
within Z. Thus there is clearly a need to extend the framework offered by ODP by considering
further style guidelines for the specification of collections of objects.

The unification techniques described above can be used with ZEST for the specification of
viewpoints. To do so, use ZEST to describe viewpoints consisting of objects or aggregates of
objects. The ZEST can then be flattened to Z, in order to generate the unification of the two
viewpoints and to check for consistency. The unification can then easily be re-assembled into
a ZEST specification if further object oriented development is required. Other object oriented
variants of Z could equally have been used for the basis of this example, in particular, Object-Z
would have provided a similar set of facilities as those we have called upon. Although unification
is applied by first flattening the ZEST, it is important to note that the benefits of using object
orientation are not lost by doing so. The encapsulation can be recovered, and the consistency
checking requirements still lie within the boundaries defined by the object encapsulation.

We have undertaken a number of case studies in order to test the hypothesis that object orien-
ted description is the preferable viewpoint specification medium, and our conclusions so far
support this claim. The studies involving non-object based descriptions were significantly harder
to check for consistency and much harder to specify in an independent fashion in the viewpoints,
the specification in [2] is another indication of the difficulty of non-object based viewpoint spec-
fications.

Conversely, the object based viewpoint descriptions were much more successful. When the
viewpoints contain only references to objects defined in other viewpoints (as opposed to specifying
any of its behaviour) consistency checking is relatively straightforward (although the viewpoints
can still be inconsistent). If two viewpoints both contain (partial) descriptions of the same object,
then there can be a non-trivial consistency checking process, however, due to encapsulation the
boundaries that inconsistency can arise within are well defined.

Examples were undertaken in a number of styles. The style of the object oriented variant chosen
did not significantly affect the success or otherwise of the viewpoint specification or unification.
There are clear merits in using Z without extended syntax, particularly in the use within ISO
initiatives. To that extent, using the work of Hall and Smith have clear advantages. Hall in
particular offers formal and well-defined support for inheritance, which is lacking for some other Z
object oriented variants. Smith’s work at the moment is not sufficiently mature or accepted. The
extended syntax approaches have advantages for the developer, who is then not constrained by
conventions for embedding object orientation in Z, but only if a clear semantics, and preferably a
flattening into Z, can be given.

5.4.1 Relation between Unification and Inheritance

It is important to recognise that unification is a ‘horizontal’ rather than ‘vertical’ development
activity. By that we mean it is used to check development at a particular stage (consistency check-
ing) or possibly to combine development specifications (unification), rather than a development
activity that serves to define the implementation more closely (as in refinement or inheritance).
Given that unification is a horizontal activity, one needs to describe the relationship between it
and vertical development activities. The relationship between unification and refinement is well
known since unification is based upon (least) refinement. We describe here the relation between
inheritance and unification.

To do so we need a formal approach to inheritance and subtyping. Hall, [23], contains a
discussion of known definitions in terms of both extensional and intensional semantics. Given that
we are interested in behaviour of specifications, we shall consider definitions due to intensional
semantics here. His intensional meaning of subclass is in terms of subclass instances being valid
implementations of the superclass, however, the definition is different from a refinement relation
(as one would expect). To exhibit subtyping there must exist a retrieve relation $A \subseteq B$ between the
superclass and subclass such that the following are true.
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S1 \( \forall \text{Superstate}; \text{Substate} • \text{pre Superop} \land \text{Abs} \Rightarrow \text{pre Subop} \)

S2 \( \forall \text{Superstate}; \text{Substate}; \text{Substate'} • \text{pre Superop} \land \text{Abs} \land \text{Subop} \Rightarrow (\exists \text{Superstate'} • \text{Abs'} \land \text{Superop}) \)

S3 \( \forall \text{Substate} • (\exists \text{Superstate} • \text{Abs}) \)

Only the third rule differs from the rules for refinement, see [23] for justification of this. Hall also compares his definition with those of Cusack [14], Lano & Haughton [34] and Liskov & Wing [38]. We are interested here in the relation between unification and the individual viewpoints. In these circumstances the retrieve relations will be partial functions (and only total if one viewpoint is degenerate). In this case Hall’s subtyping imply subtyping in the sense of Cusack, Lano & Haughton and Liskov & Wing (ignoring the history predicates of the latter two). In particular, the rules S1-3 suffice for subtyping in both Hall and ZEST, and we will thus work with this definition.

It is easy to construct examples to show that the unification of two viewpoints is not in general a subtype of each viewpoint. However, this is unsurprising because one viewpoint is only a partial description of an objects behaviour. Instead the natural result to seek is the following:

**Theorem 2** Let \( P_i, O_i \) be objects in viewpoint \( i \). Let \( P_i \) be a subtype of \( O_i \). Then \( U(P_1, P_2) \) is a subtype of \( U(O_1, O_2) \), where \( U \) is the unification operator between viewpoints.

**Proof**
The full proof involves construction of appropriate retrieve relations between \( U(P_1, P_2) \) and \( U(O_1, O_2) \) in a manner similar to the proof that unification is the least refinement, see [16]. The outline of the proof is as follows:

The subtyping rules S1 and S2 between \( U(P_1, P_2) \) and \( U(O_1, O_2) \) are satisfied because unification is the least refinement.

For S3, note that every state in the \( U(O_1, O_2) \) unification appears in either \( O_1 \) or \( O_2 \) or both. Thus every state in \( U(P_1, P_2) \) is related to some state in \( U(O_1, O_2) \) via the retrieve relation defined for the least refinement. \( \square \)

It is straightforward to construct examples to show the converse is not true, that is \( U(P_1, P_2) \) being a subtype of \( U(O_1, O_2) \) does not imply that \( P_i \) is a subtype of \( O_i \).

The theorem then provides a sound footing for the use of object oriented techniques in viewpoint descriptions. The relationship between unification and multiple inheritance is clearly of importance (especially w.r.t. method consistency), and is currently under investigation.
Chapter 6

Conclusion

In conclusion of this deliverable we summarise the results so far and describe the key open problems that remain.

6.1 Summary of results

We have reported on three major areas of investigation:

1. Defining consistency;
2. Consistency checking in LOTOS; and
3. Consistency checking in Z.

6.1.1 Defining consistency

The RM-ODP contains three different definitions of consistency. These definitions have been formalised and related to one another (see chapter 3). The different definitions seem incomparable unless they had been instantiated with a particular FDT. To resolve this deficiency a more general definition of consistency was formulated that can incorporate all three RM-ODP definitions.

In general terms, a number of viewpoint specifications are consistent if and only if a new specification can be found that is a development of all viewpoint specifications. In addition, the new specification is required to be internally valid (i.e. can be implemented). We distinguish between balanced consistency, where one development relation is used in all viewpoints, and unbalanced consistency, where different development relations can be used in different viewpoints.

If a number of viewpoint specifications are consistent with respect to certain development relations, then a common development can be found, which we call unification. Conversely, the existence of an internally valid unification implies consistency of the viewpoint specifications.

Although the definition of consistency is general enough to cope with inter language consistency checking, we have initially considered intra language consistency checking only.

6.1.2 Consistency checking in LOTOS

In chapter 4, we describe techniques to check for binary, balanced consistency of LOTOS specifications with respect to four different development relations: trace preorder, reduction, extension and testing equivalence. Techniques to construct unifications of two LOTOS specifications with respect to these development relations are considered.

In LOTOS it is possible to find necessary and sufficient conditions to show that two specifications are consistent. Once consistency has been established, it is possible to construct a unification. However, the generated unification is usually not the least common development.
6.1.3 Consistency checking in Z

In chapter 5, techniques are given to unify Z viewpoint specifications with respect to the Z refinement relation. It also describes how the internal validity of the derived unification can be verified. Since the derived unification is always the least unification, internal validity of the unification implies the consistency of the original specifications.

The use of object oriented specification techniques was also considered. Initial evidence indicates such techniques can reduce the complexity of the consistency checking process.

6.2 Open problems

This deliverable has outlined a number of issues relating to consistency checking mechanisms for ODP. There clearly remains much work to be done, we discuss briefly here some of the future directions for research in this area.

6.2.1 Inter language consistency checking

It is commonly recognised that different FDTs will be applicable to different viewpoints. In this respect, to be of practical use, inter language consistency checking mechanisms have to be built upon intralanguage mechanisms. An approach could be envisaged where to check the consistency of two different FDT viewpoint specifications, one specification is translated into the language of the other, after which an intralanguage mechanism is applied. If this is possible then checking across language boundaries becomes feasible.

Translation is discussed briefly below, languages to be considered include LOTOS, Z and IDLs. The translation between IDLs and Z clearly has practical significance, and we will investigate this as resources allow. The translation between LOTOS and Z would be a fundamentally important result, given the different semantic bases for the two languages, and we are considering to what extent this can be achieved.

6.2.2 Translation

There has been some success in relating formal languages that have similar underlying semantics, e.g. [43, 5]. However, the common semantics used in these approaches is typically very ugly. ODP consistency checking requires translation across FDT families. Some directions that could be pursued to make such translations possible are discussed below.

**Syntactic translation** Translation based upon a direct relation of syntactic terms in one FDT to terms in another FDT is one possible approach. However, it is difficult to envisage how such an approach could offer a general solution. In particular, a lot of semantic meaning will certainly be lost in such a crude translation of FDTs. Partial syntactic translations may, however, be feasible.

**Common underlying semantics** Another, more promising, approach is to define a common underlying semantic model for the required FDTs. Specifications could then be translated into the common semantic domain, in which a consistency check can be performed. Such translation could either use the semantics of one of the FDTs as the intermediate semantics or use a third semantics. The former of these is not fully general; for example, Z and LOTOS are so fundamentally different that relating one to the others semantic model is very difficult to envisage. Relating FDTs using a third intermediate form is a more likely approach.

A sufficiently general semantic model has been developed in [57]. One-sorted first-order predicate logic is proposed to capture the semantics of specifications in different formalisms. The proposed semantic model is very general, but it does not necessarily have the same properties as the standard semantics of the FDTs. Also, there are certain features of specification languages which cannot be captured in first-order logic. Nevertheless, the proposed method presents a promising step towards a general technique for consistency checking.
There have been several other attempts to relate different formal languages:

- A link between model based action systems (and thereby Z) and CSP has been made by showing that refinements (forwards and backwards simulation) in an action system are sound and jointly complete with respect to the notion of refinement in CSP \cite{55}.

- The requirement for highly expressive intermediate semantics suggests that logical notations may be appropriate. \cite{20} and \cite{7} consider logical characterisations of LOTOS in temporal logic. However, relating temporal logic to the Z first order logic remains an open issue. Categorical approaches and the theory of institutions offer a possible solution \cite{7}.

- A final alternative which has the benefit of being ODP specific is suggested by the work of \cite{13}. This work offers a direct denotational semantics for the computational viewpoint language. This semantics could, theoretically, be used to relate different FDT interpretations of the computational viewpoint language. Clearly, this work does not give a complete solution to consistency as the semantics are restricted to a single viewpoint. However, it may be possible to extrapolate this approach to a general solution.

It is clear, though, that a usable translation mechanism is likely to represent a pragmatic, compromise solution. In particular, complete preservation of semantic meaning during translation will not be possible. In addition, different viewpoints describe different sets of features and thus may not be directly translatable between each other.

### 6.2.3 ODP specific concepts

The work presented in this deliverable has concentrated on a general framework for consistency checking mechanisms and, in particular, we have not considered specific ODP viewpoints. Consideration of the viewpoints will include discussion of the role of the ODP architectural semantics. Specifically, part 4 should provide a basis for relating FDTs. ODP concepts, in particular viewpoint languages, are defined in different FDTs in the architectural semantics. Thus, when relating complete viewpoint specifications in different FDTs these definitions can be used as components of a consistency check.

However, it is important to note that the architectural semantics will only provide a framework for consistency checking. Actual viewpoint language specifications will extend the ODP architectural semantics, which are non-prescriptive by nature, with FDT specific behaviour. There is then a need to combine the framework provided by the architectural semantics with actual consistency checking relationships arising from FDTs. Such cross viewpoint consistency checks will clearly involve correspondence rules, and further work is required on identifying appropriate correspondences between the viewpoints in order to provide formal support for cross viewpoint mappings.

### 6.2.4 Tool Development

Work on tool development based upon consistency checking mechanisms is clearly important if this work is to move into application areas. We envisage roles for unification tools (in particular for Z), coupled to semi-automatic consistency checkers which will aim to provide support for consistency checks. The extent of such tool development will be considered in the light of available resources.

### 6.2.5 Object orientation

The ODP viewpoint languages are object based. Current consistency checking and unification algorithms do not take the object oriented nature of the specifications into consideration. Further research is necessary to assess in which way object orientation will effect the consistency checking process. Some experiments with object oriented Z specifications suggest that the encapsulation property of objects will simplify the consistency checking mechanisms.
6.3 Future plans

In this section, we list the activities planned for the future. In the short term, we plan to consider the following tasks:

- Identification and characterisation of ODP correspondence rules;
- Building a unification and consistency checking tool for $Z$;
- Extending the LOTOS consistency checking techniques to unbalanced and global consistency;
- Starting work on a LOTOS tool for unification and consistency checking;
- Investigation of inter language consistency checking between LOTOS and $Z$ specifications;
- Identification of case studies for intra language consistency checking in LOTOS and $Z$; and
- Provision of input to RM-ODP Part 1 on consistency checking.

For the long term we envisage the following activities:

- Completion of unification and consistency checking tools for both $Z$ and LOTOS;
- Building a tool for consistency checking and/or translation between $Z$ and LOTOS;
- Performing several case studies on both intra language and inter language consistency checking; and
- Dissemination of final results to the ODP community.
Bibliography


